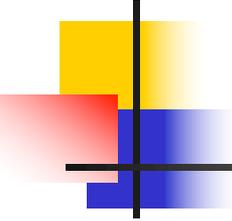


Knowledge Compilation:

Principles and Applications

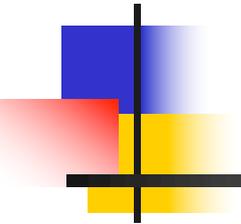
Adnan Darwiche

Computer Science Department, UCLA



Agenda

- Languages and Operations
- Knowledge Compilers
- Applications:
 - Explaining and Verifying AI Systems
 - Probabilistic Reasoning
 - Machine Learning



Languages & Operations

Knowledge Compilation

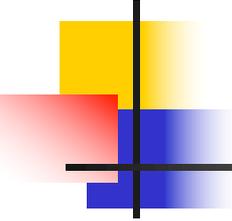
$A \ \& \ okX \Rightarrow \neg B$
 $\neg A \ \& \ okX \Rightarrow B$
 $B \ \& \ okY \Rightarrow \neg C$
 $\neg B \ \& \ okY \Rightarrow C$

Compiler

**Compiled
Structure**

Queries

**Evaluator
(Polytime)**



Knowledge Compilation

$A \ \& \ okX \Rightarrow \neg B$
 $\neg A \ \& \ okX \Rightarrow B$

 $B \ \& \ okY \Rightarrow \neg C$
 $\neg B \ \& \ okY \Rightarrow C$

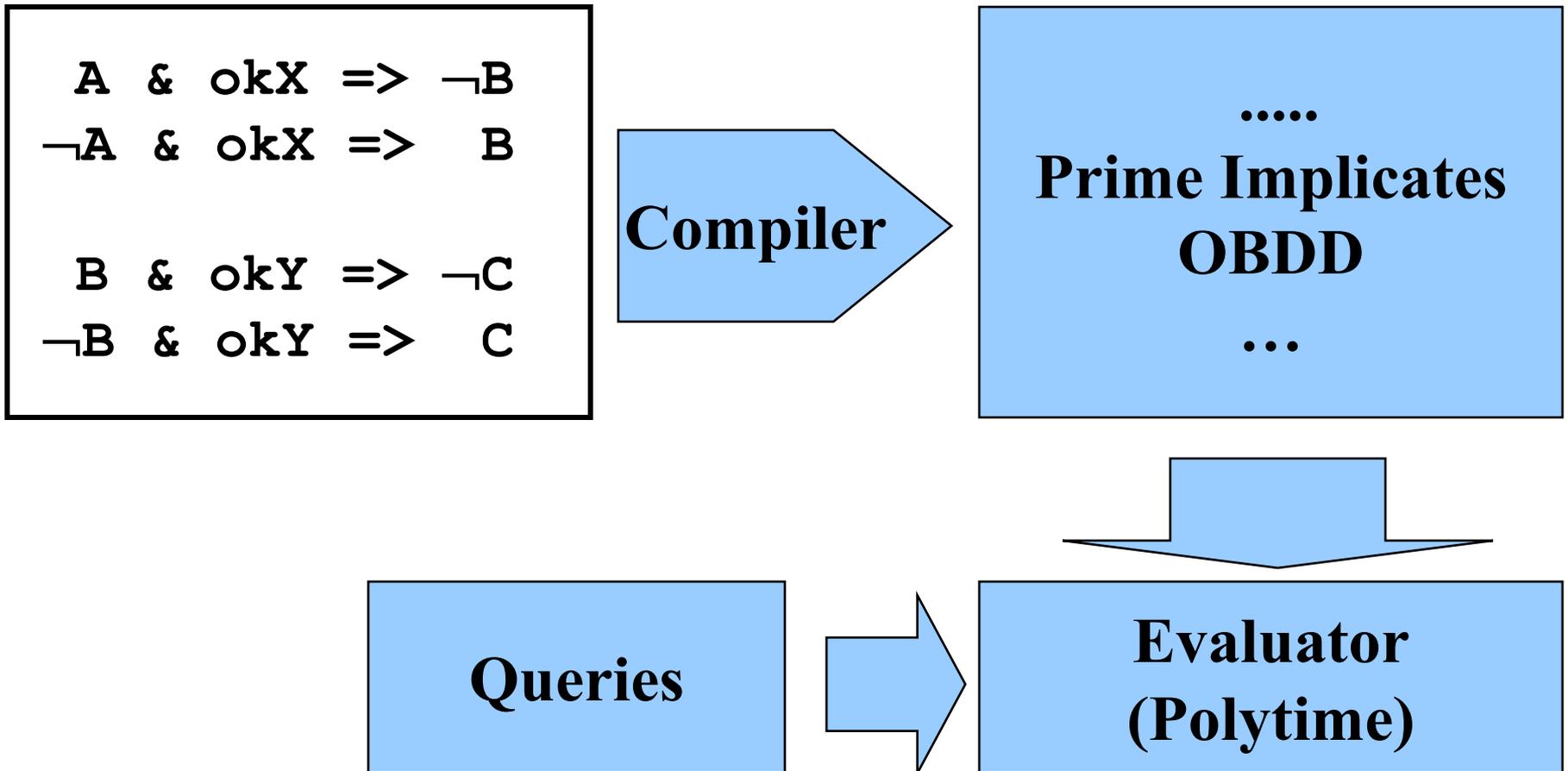
Compiler

?

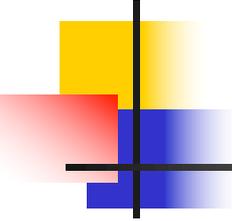
Queries

**Evaluator
(Polytime)**

Knowledge Compilation

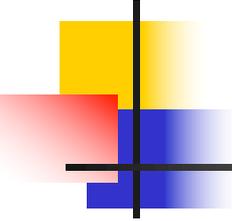


Knowledge Compilation Map



- What's the space of possible target compilation languages?
 - Can it be synthesized in a semantically systematic way?
- How do the languages compare?
 - Succinctness (relative size)
 - Operations they support in polytime

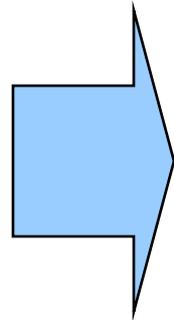
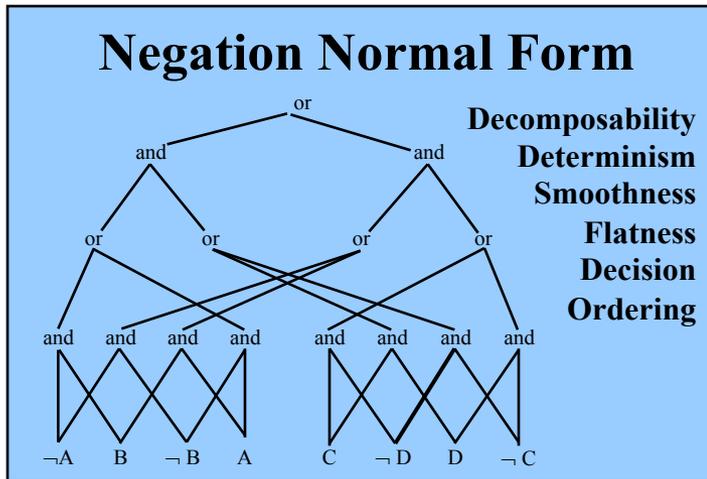
Knowledge Compilation



Map

- For a given application: identify needed operations
- Choose most succinct language that supports desired operations
- Compile knowledge base into chosen language

A Knowledge Compilation MAP

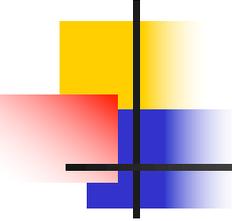


Polytime Operations

Consistency (CO)
Validity (VA)
Clausal entailment (CE)
Sentential entailment (SE)
Implicant testing (IP)
Equivalence testing (EQ)
Model Counting (CT)
Model enumeration (ME)

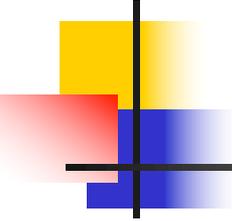
Projection (exist. quantification)
Conditioning
Conjoin, Disjoin, Negate

Succinctness



Propositional Logic

- **Literal** $X, \neg X$
- **Clause** $(X \vee \neg Y \vee \neg Z)$
- **Term** $(\neg X \wedge Y \wedge Z)$
- **CNF**: Conjunctive Normal Form
 $(X \vee \neg Y \vee \neg Z) \wedge \dots \wedge (Y \vee \neg W)$
- **DNF**: Disjunctive Normal Form
 $(\neg X \wedge Y \wedge Z) \vee \dots \vee (X \wedge \neg Z \wedge W)$



Propositional Logic

- **Truth assignment (TA)**

$X : true, Y : false, Z : true, W : false$

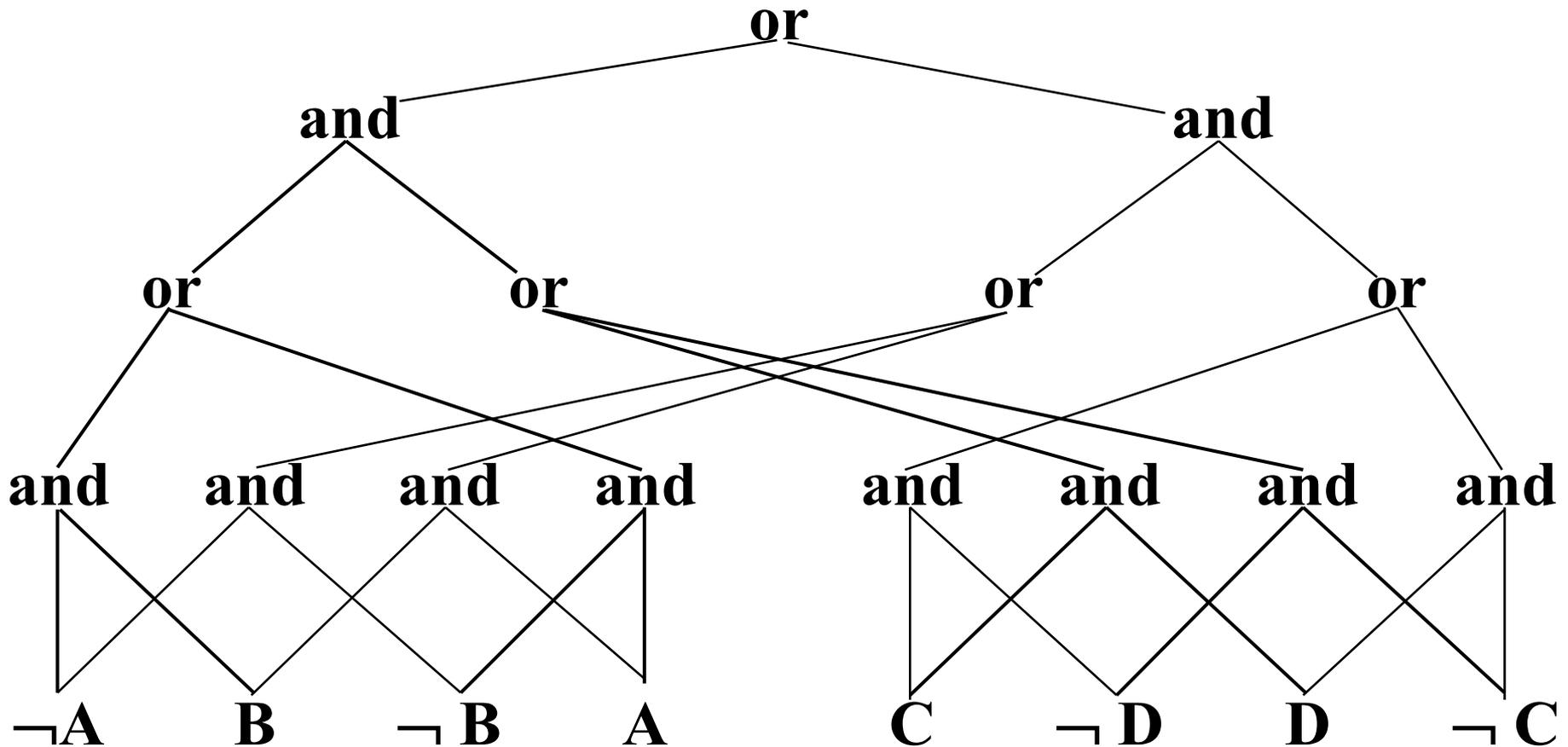
- **TA satisfies sentence (model)**

$(X \vee \neg Y \vee \neg Z) \wedge \dots \wedge (Y \vee \neg W)$

- **Following TA is not a model**

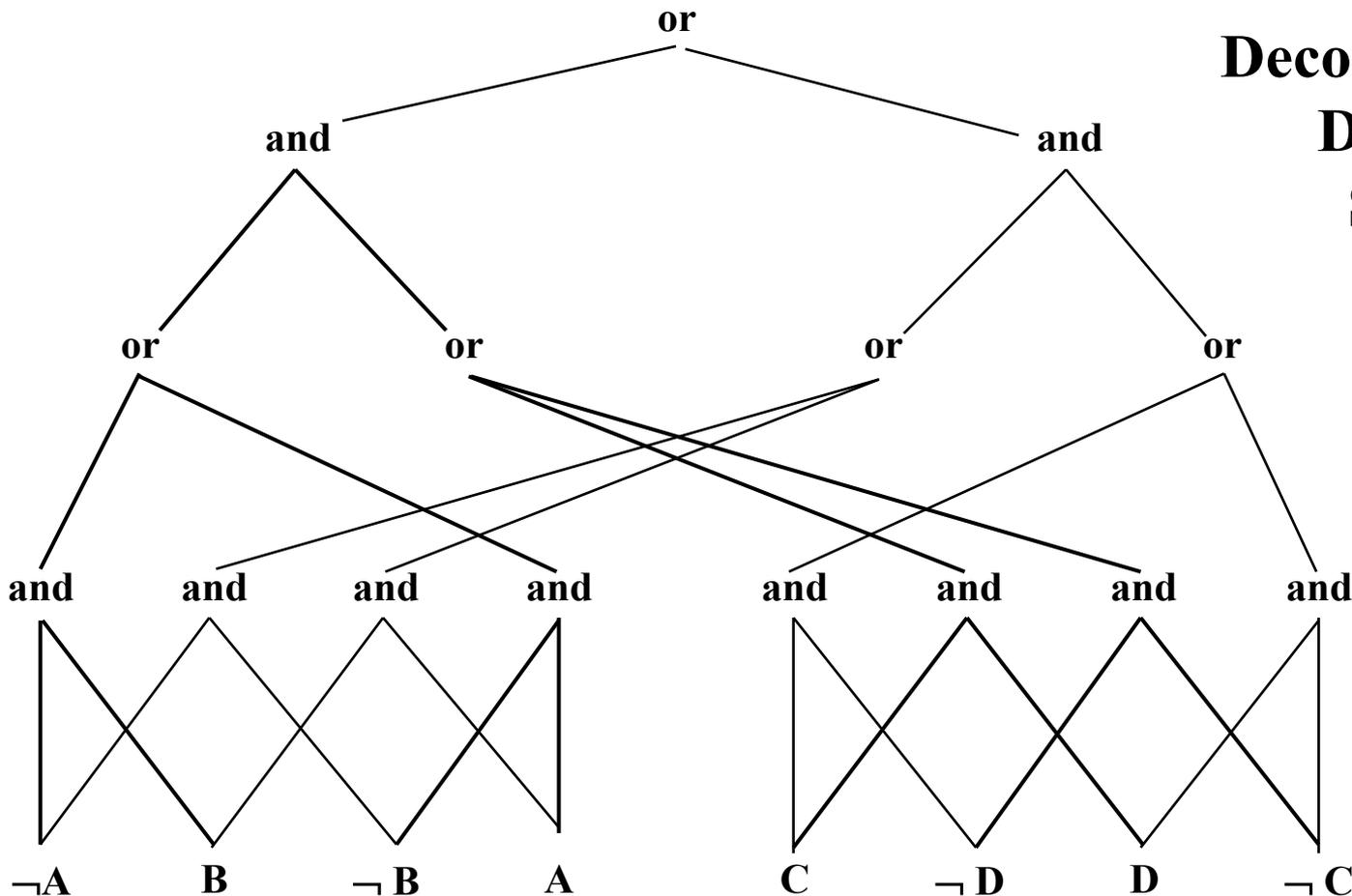
$X : true, Y : false, Z : true, W : true$

Negation Normal Form



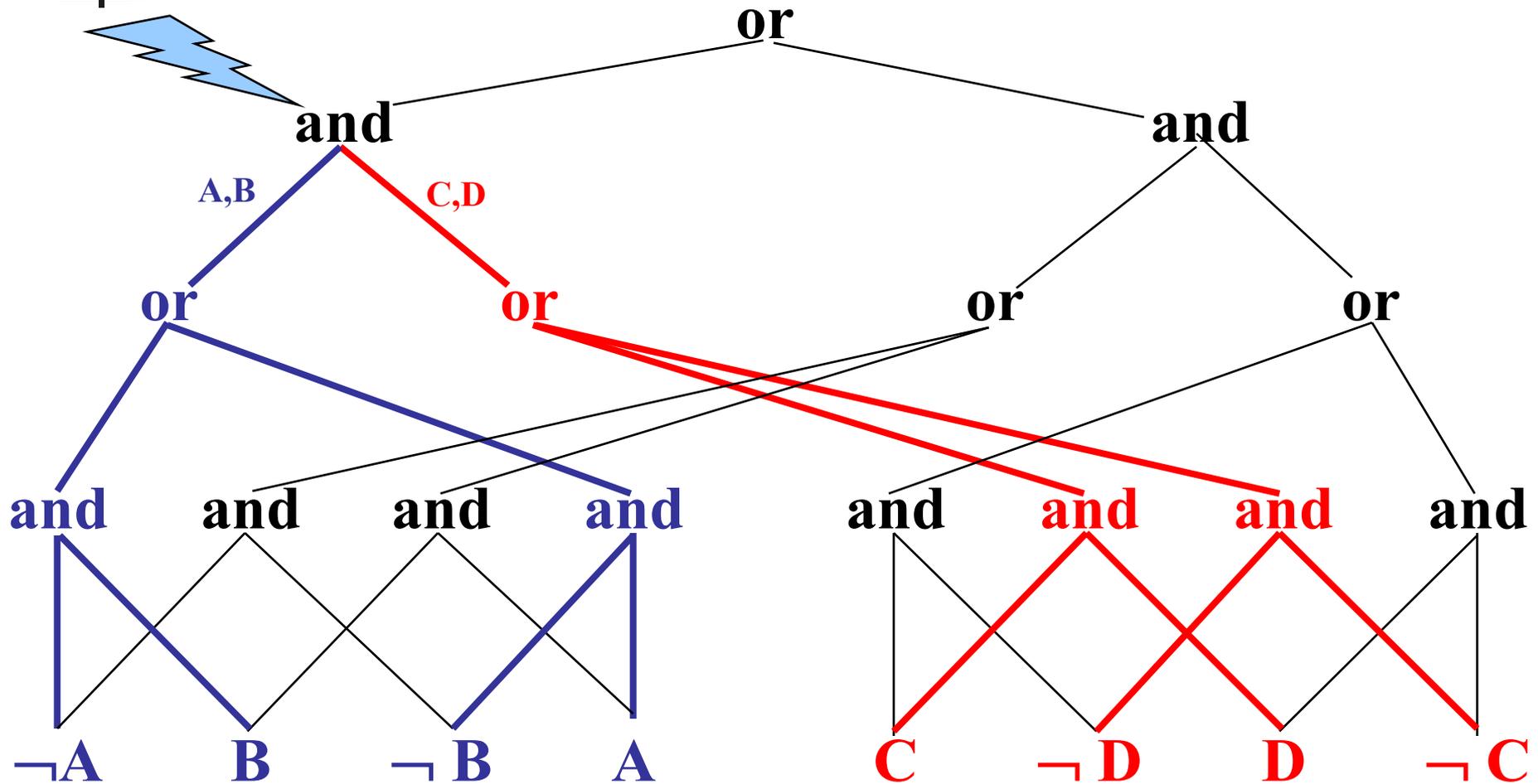
rooted DAG (Circuit)

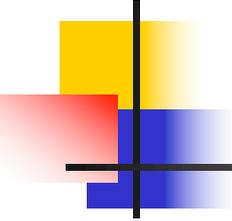
Negation Normal Form



Decomposability
Determinism
Smoothness
Flatness
Decision
Ordering

Decomposability



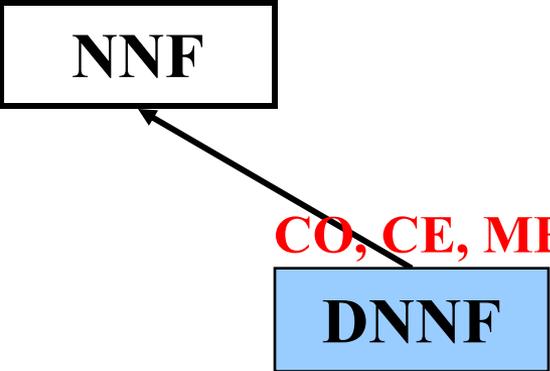


NNF Subsets

NNF

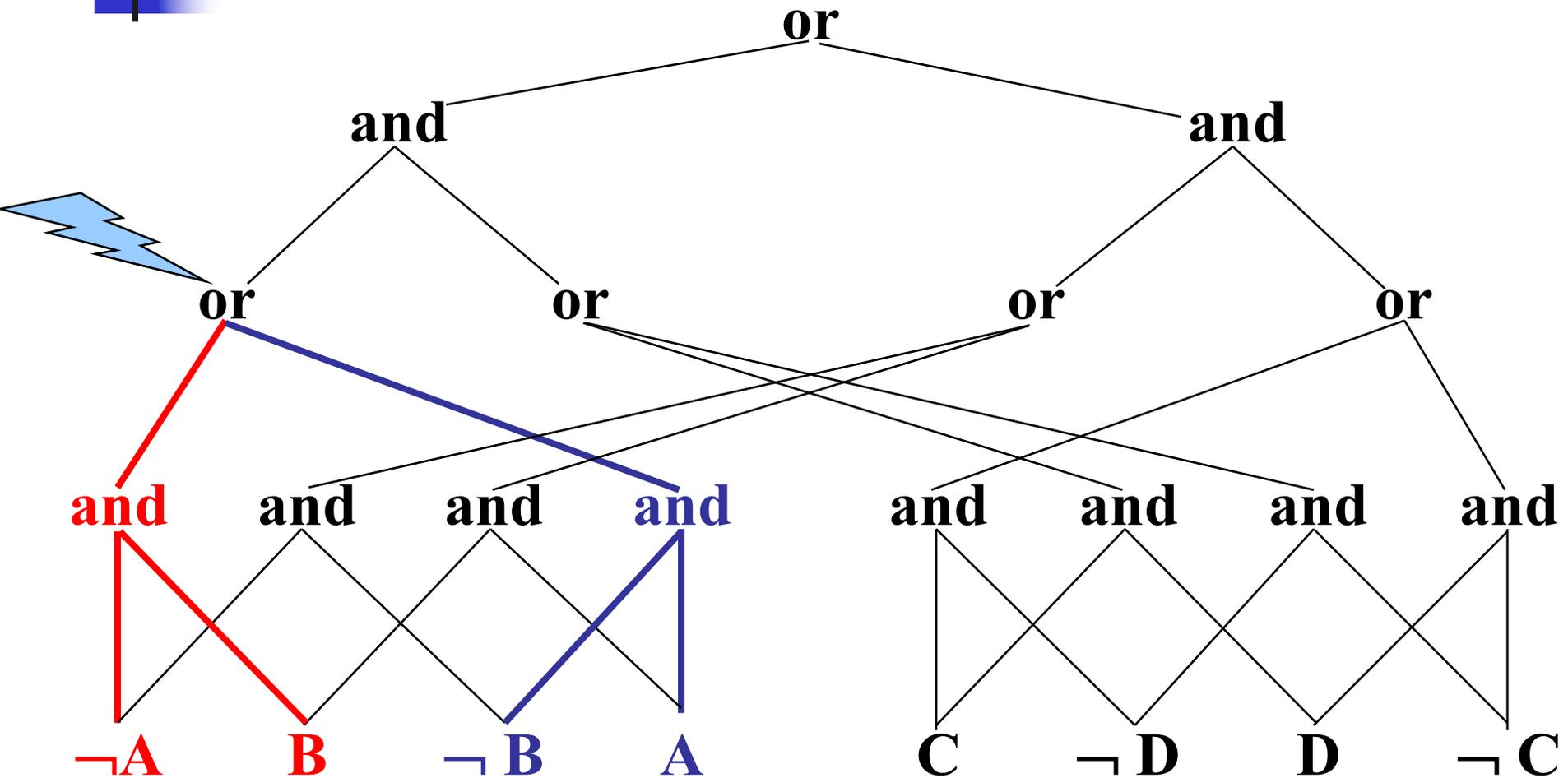
CO, CE, ME

DNNF

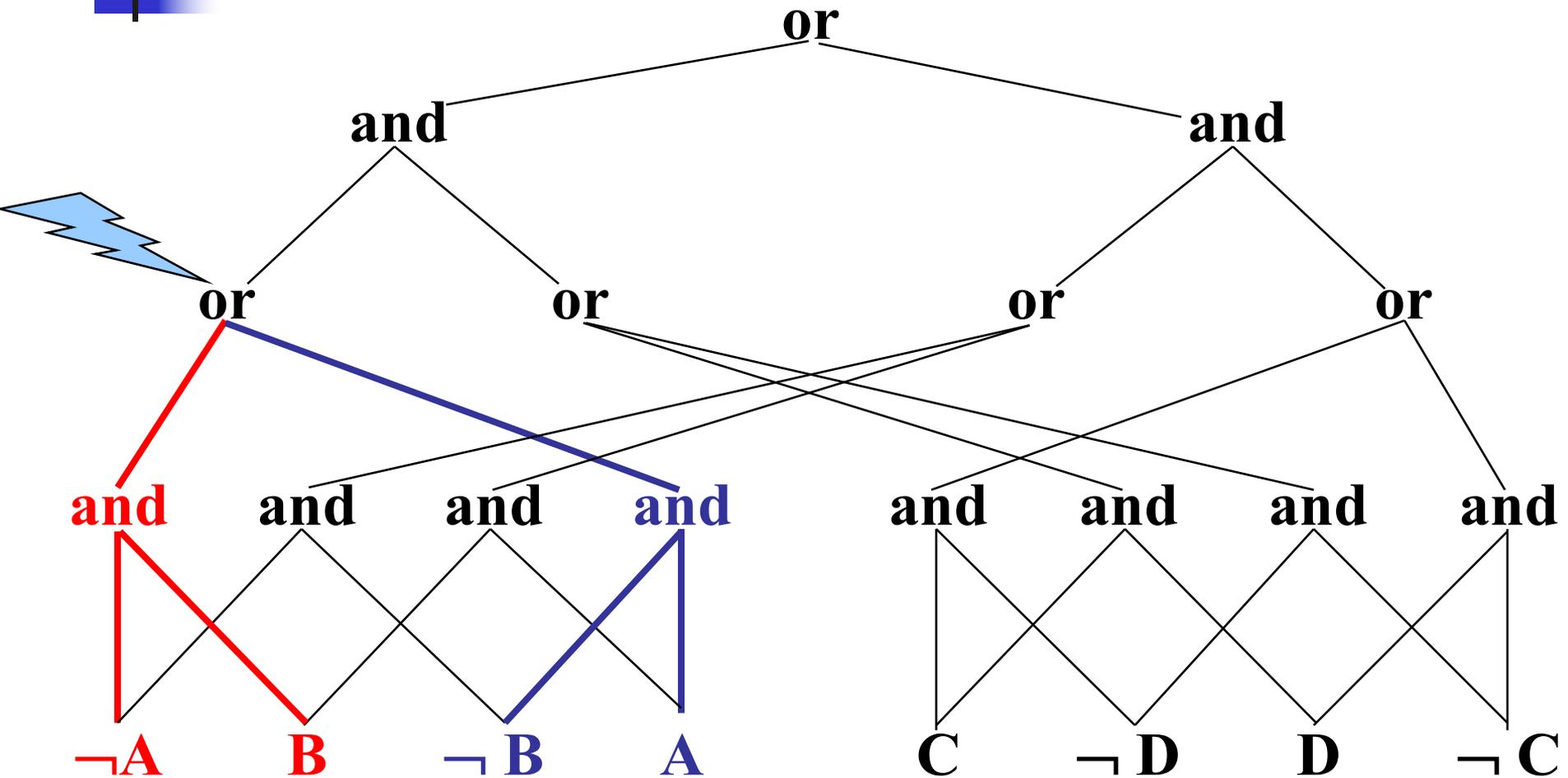


```
graph BT; DNNF[DNNF] -- "CO, CE, ME" --> NNF[NNF];
```

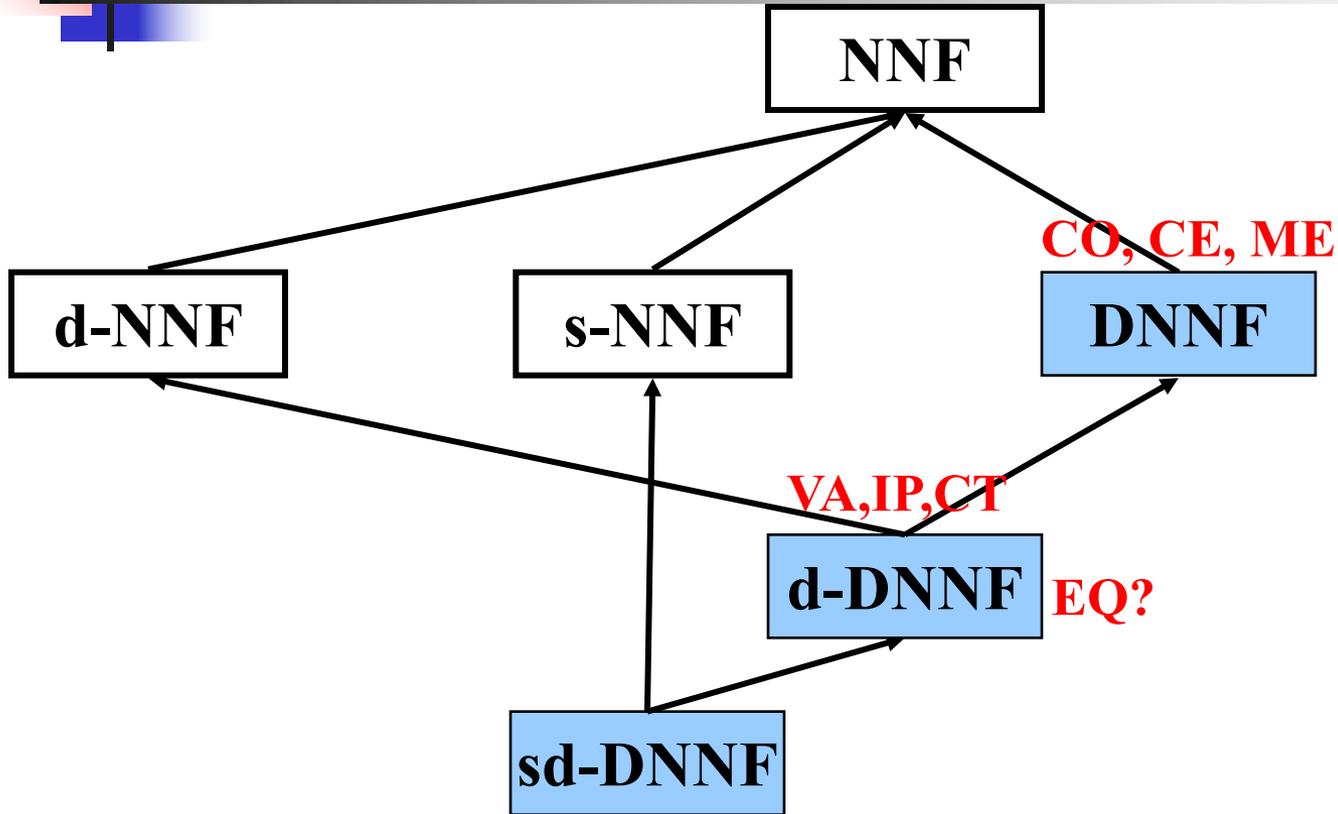
Determinism



Smoothness

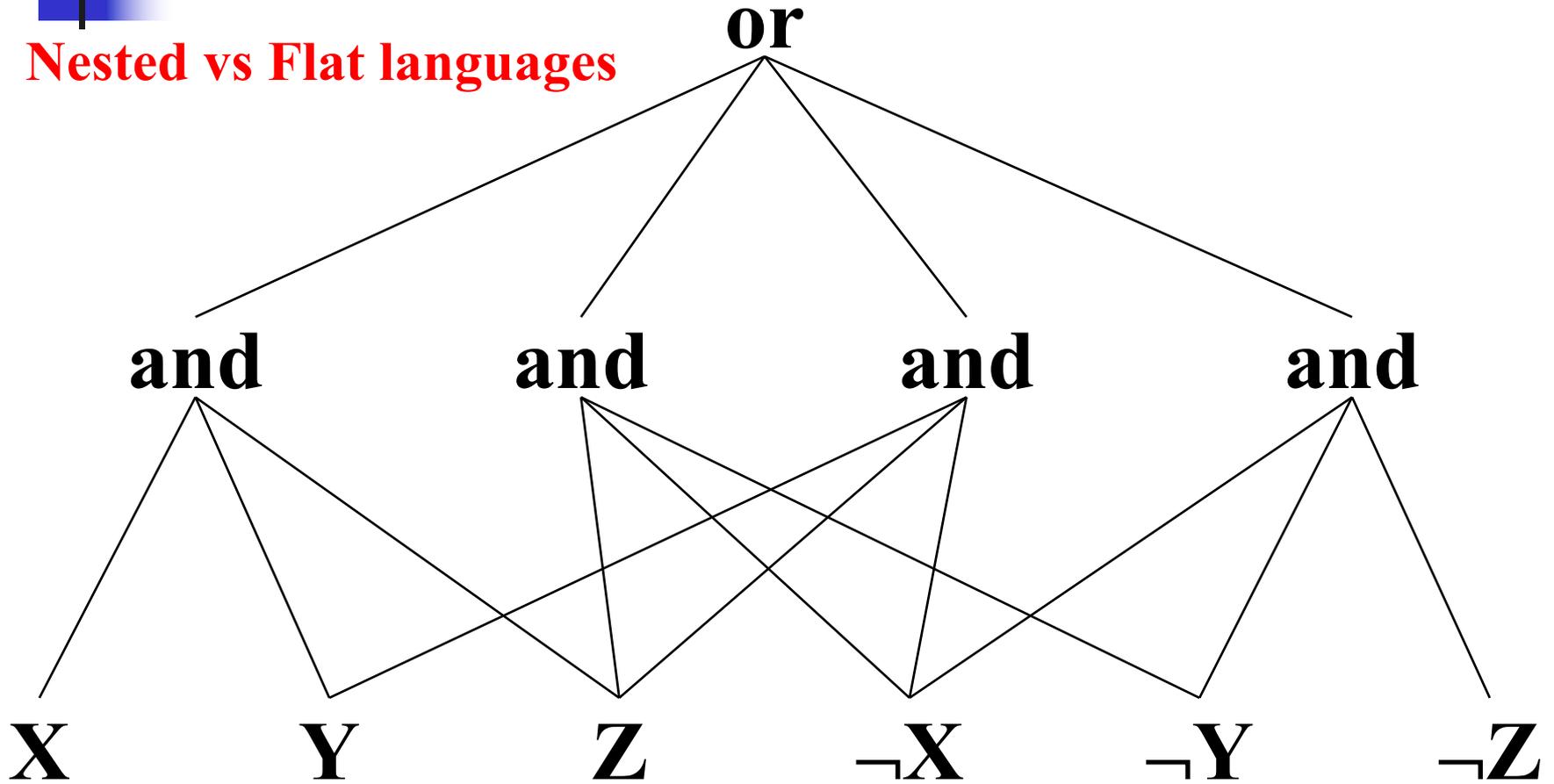


NNF Subsets

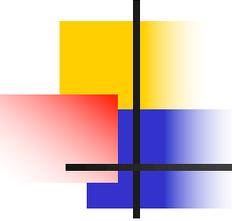


Flatness

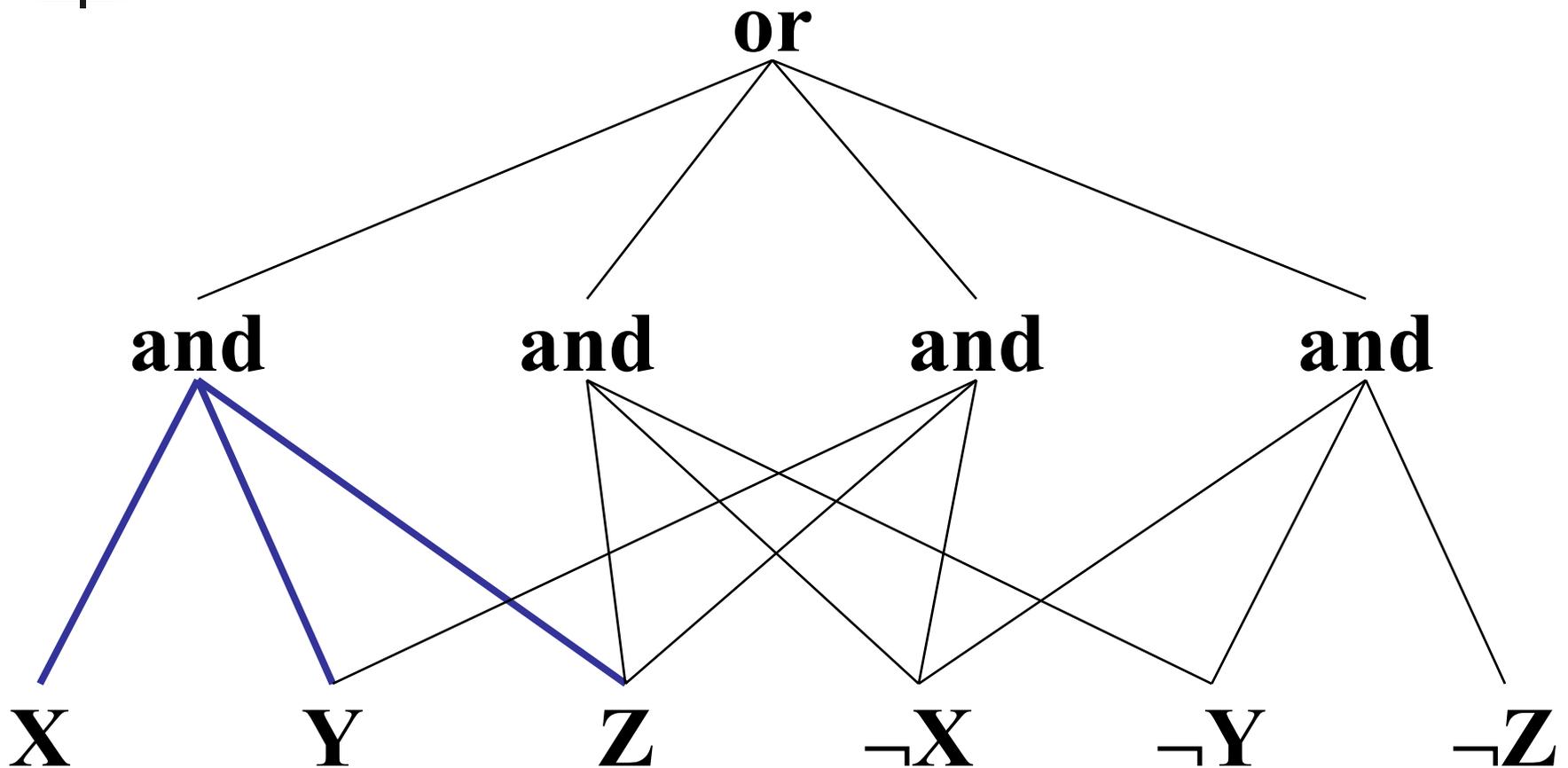
Nested vs Flat languages



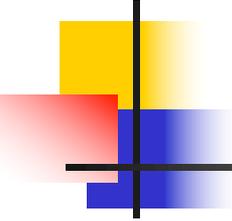
$$(X \wedge Y \wedge Z) \vee (Z \wedge \neg X \wedge \neg Y) \vee (Y \wedge Z \wedge \neg X) \vee (\neg X \wedge \neg Y \wedge \neg Z)$$



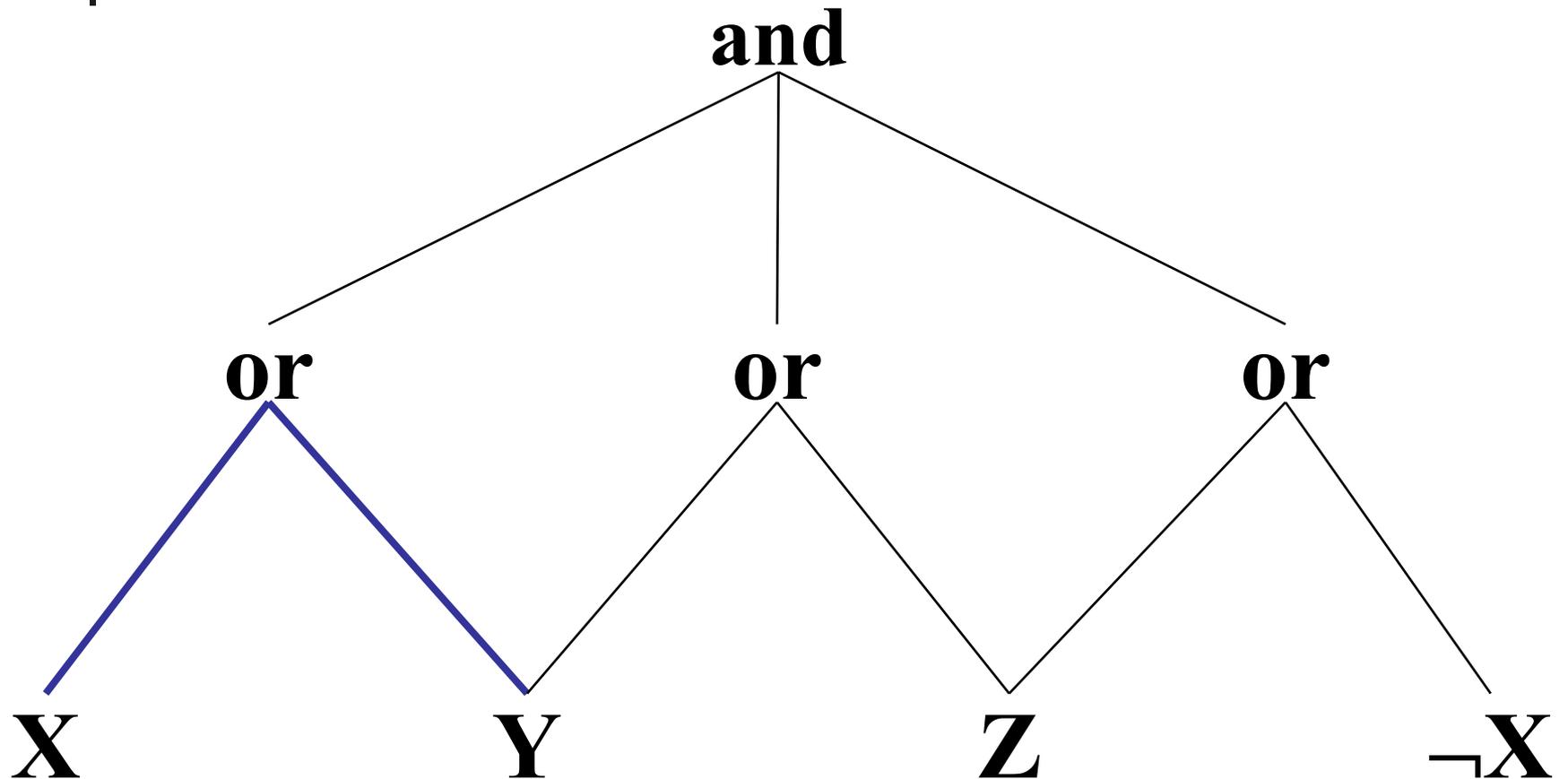
Simple Conjunction



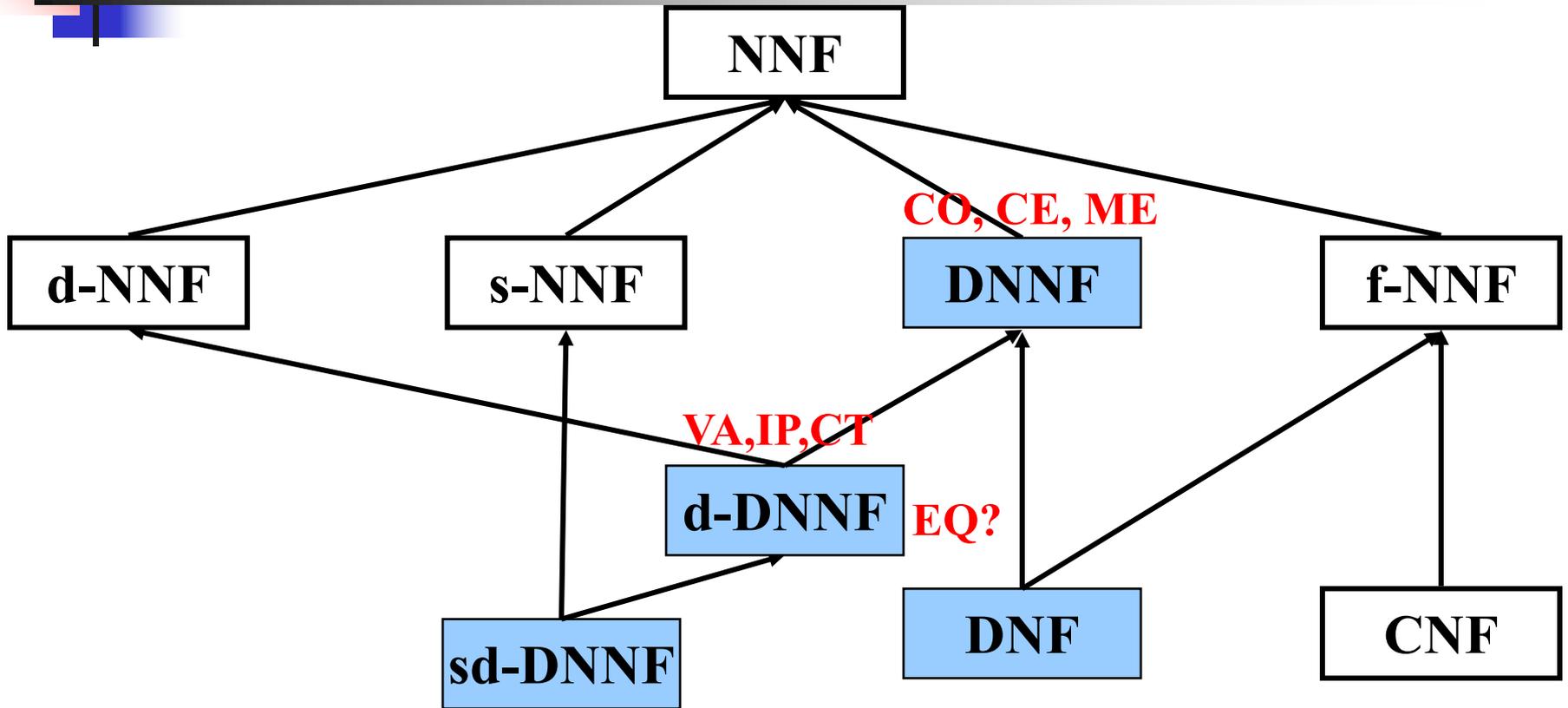
Simple conjunction implies decomposability



Simple Disjunction



NNF Subsets



Prime Implicates (PI)

Resolution that:

$$\frac{(\alpha \vee X), (\beta \vee \neg X)}{(\alpha \vee \beta)}$$

er

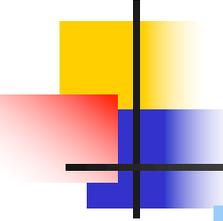
e CNF, it must
e CNF

- CNF:

$$(\neg A \vee B) \wedge (\neg B \vee C) \wedge (\neg C \vee D)$$

- PI:

$$\begin{aligned} &(\neg A \vee B) \wedge (\neg B \vee C) \wedge (\neg C \vee D) \wedge \\ &(\neg A \vee C) \wedge (\neg A \vee D) \wedge (\neg B \vee D) \end{aligned}$$



Prime Implicants (IP)

Consensus h that:

$$\frac{(\alpha \wedge X), (\beta \wedge \neg X)}{(\alpha \wedge \beta)}$$

it must imply a

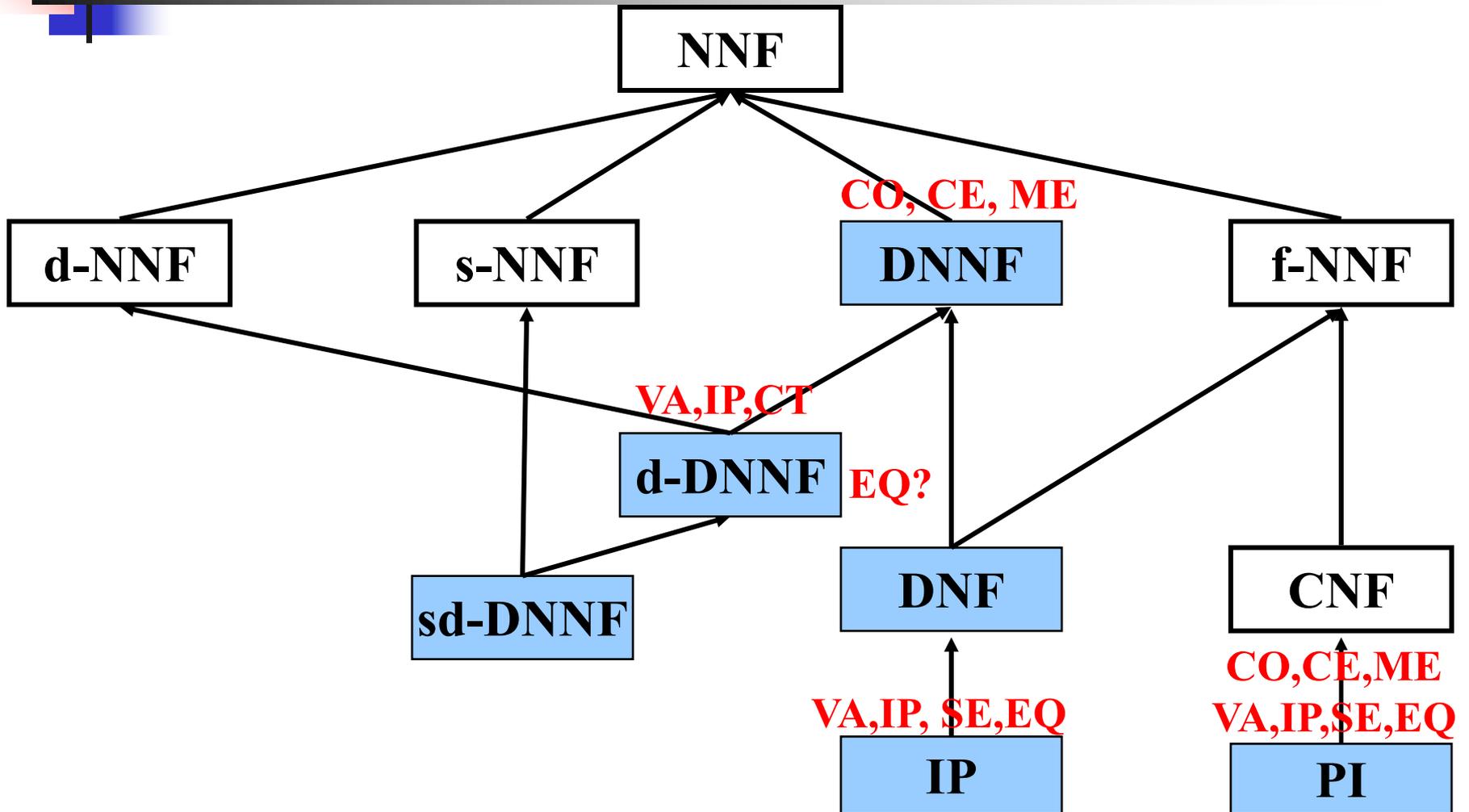
- DNF:

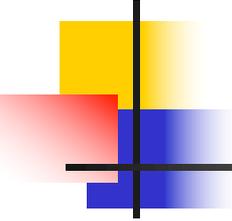
$$(A \wedge B) \vee (\neg B \wedge C)$$

- IP:

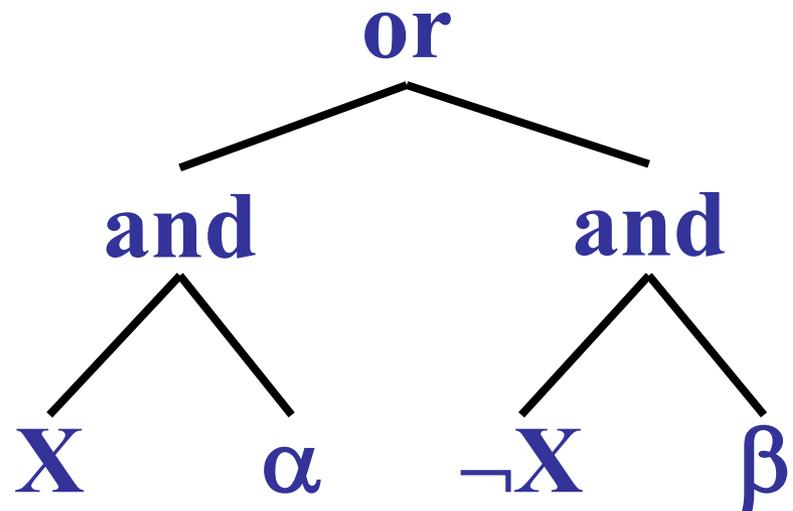
$$(A \wedge B) \vee (\neg B \wedge C) \vee (A \wedge C)$$

NNF Subsets



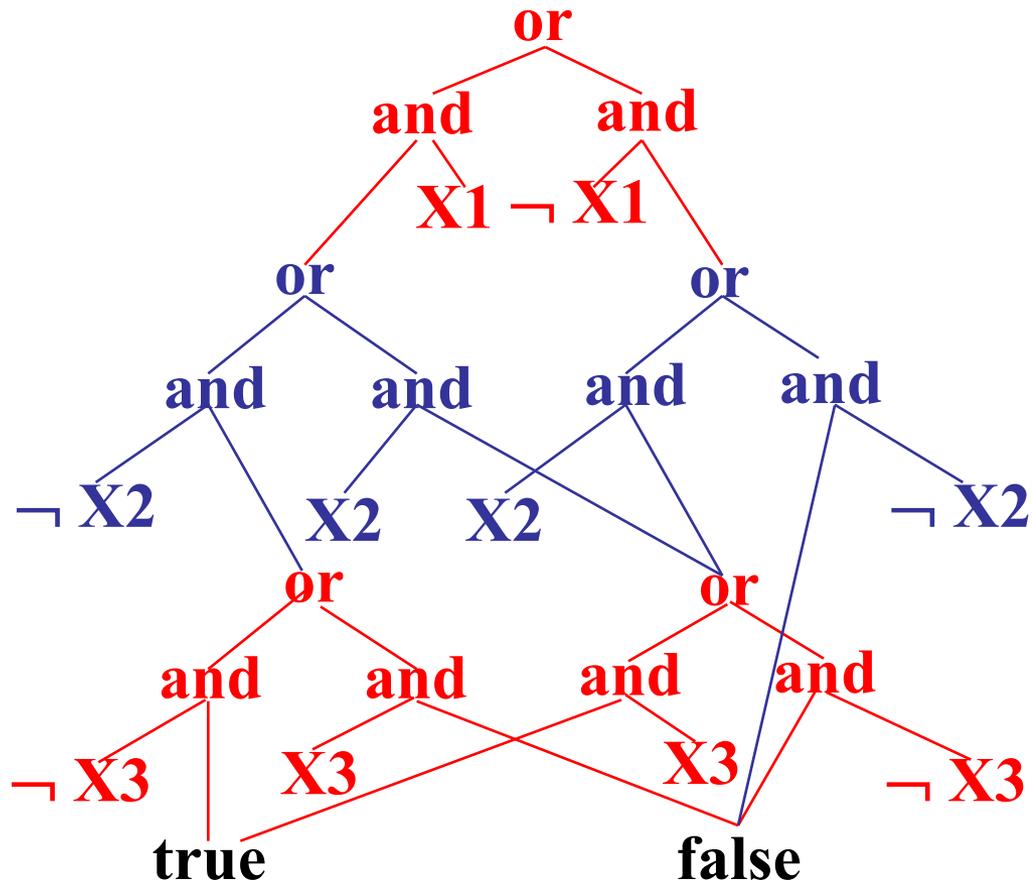


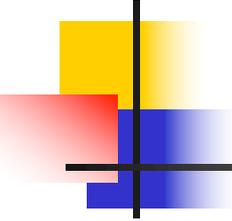
Decision



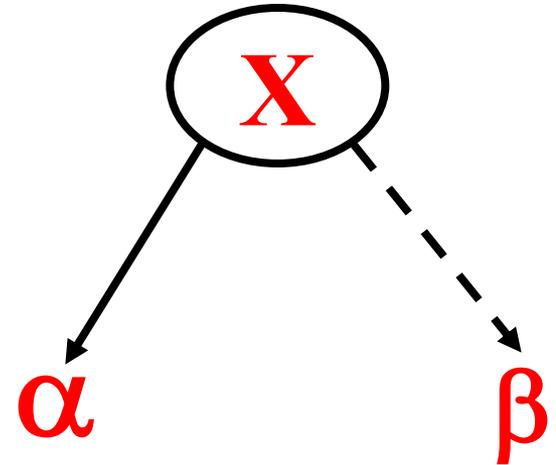
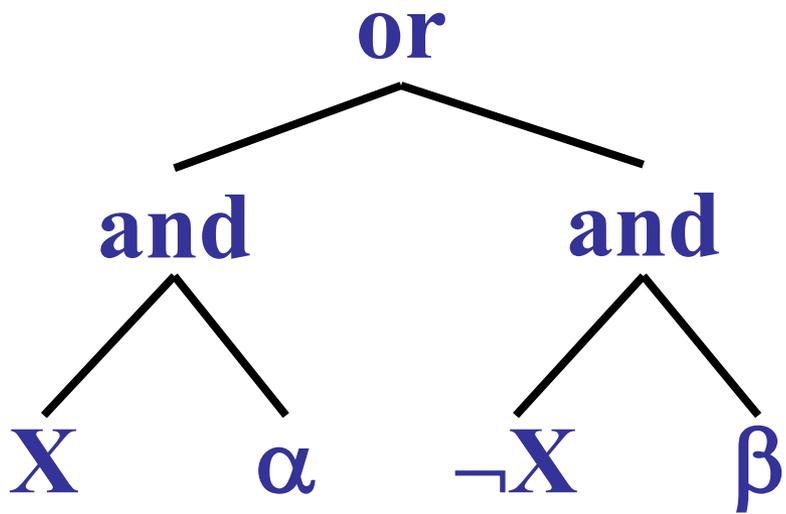
α, β : Are decision nodes

Decision

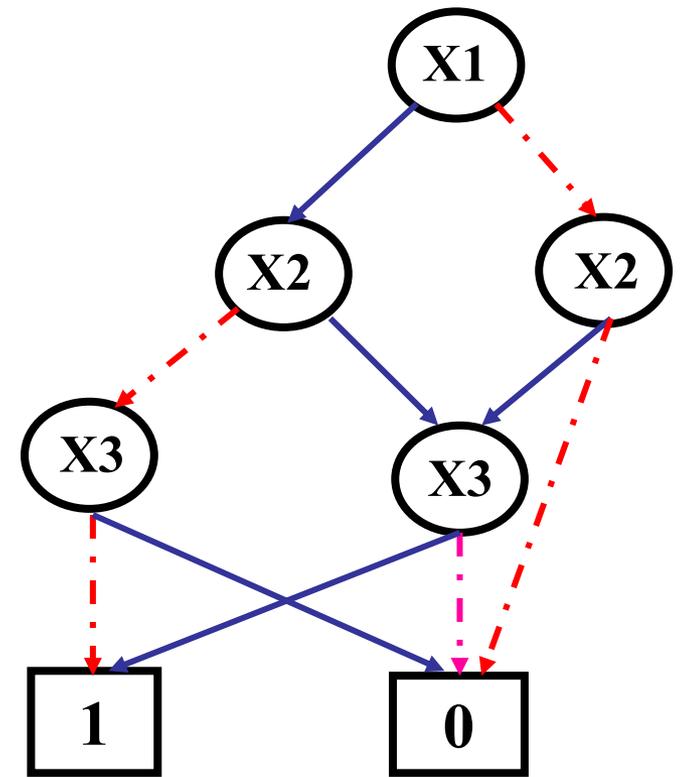
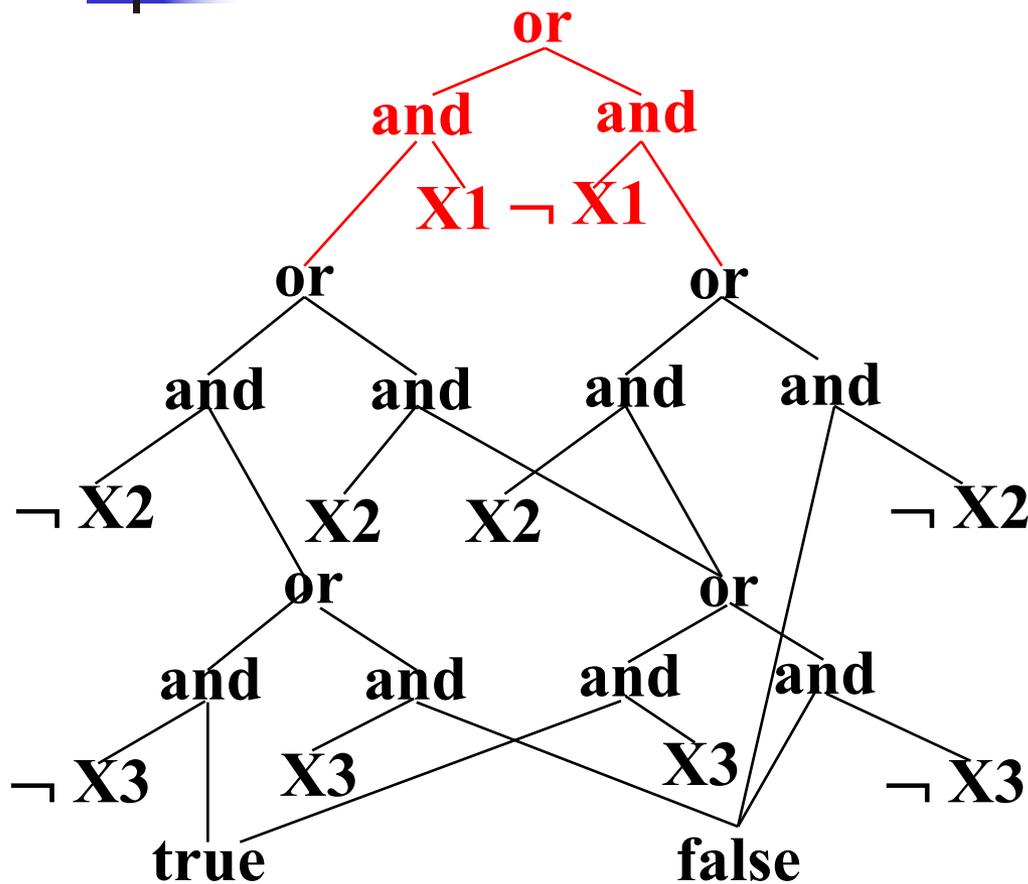




Decision

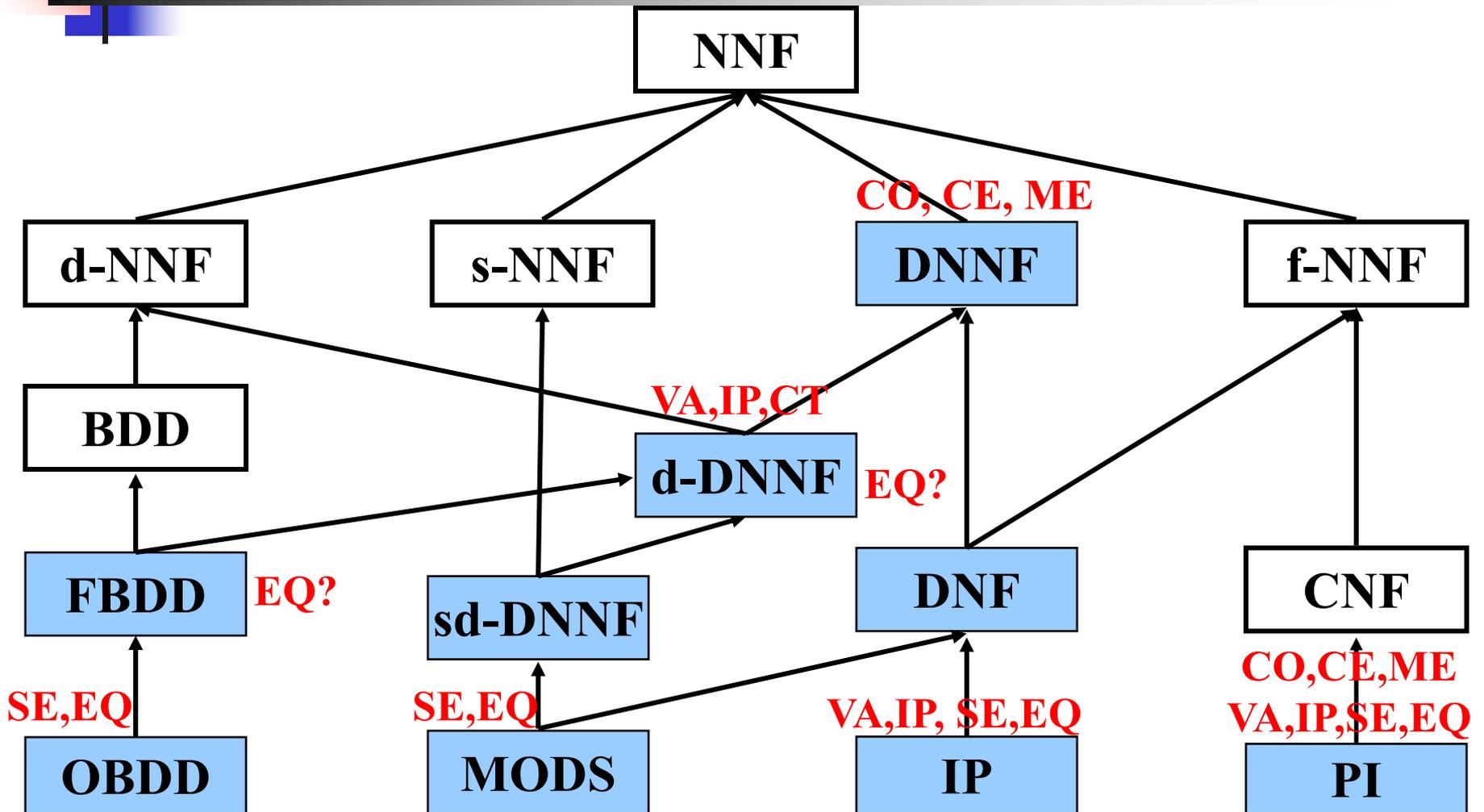


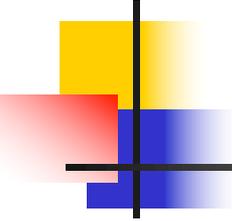
Binary Decision Diagrams (BDDs)



Decision implies determinism

NNF Subsets





Language Succinctness

L1 at least as succinct as L2

$$\mathbf{L1} \leq \mathbf{L2}$$

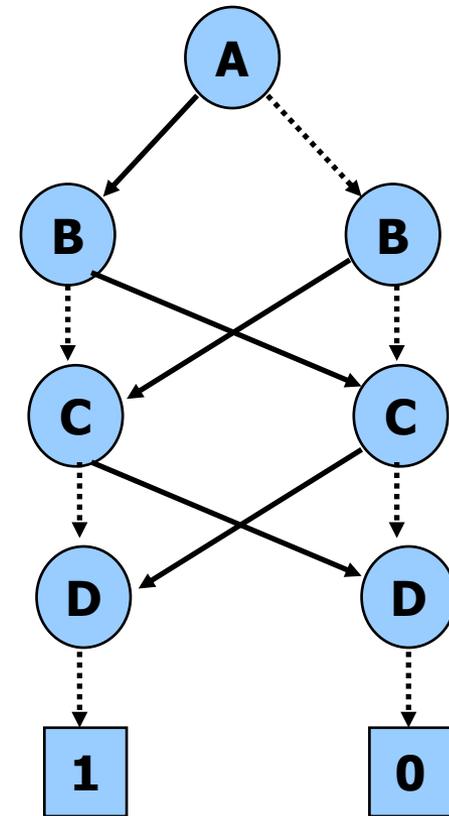
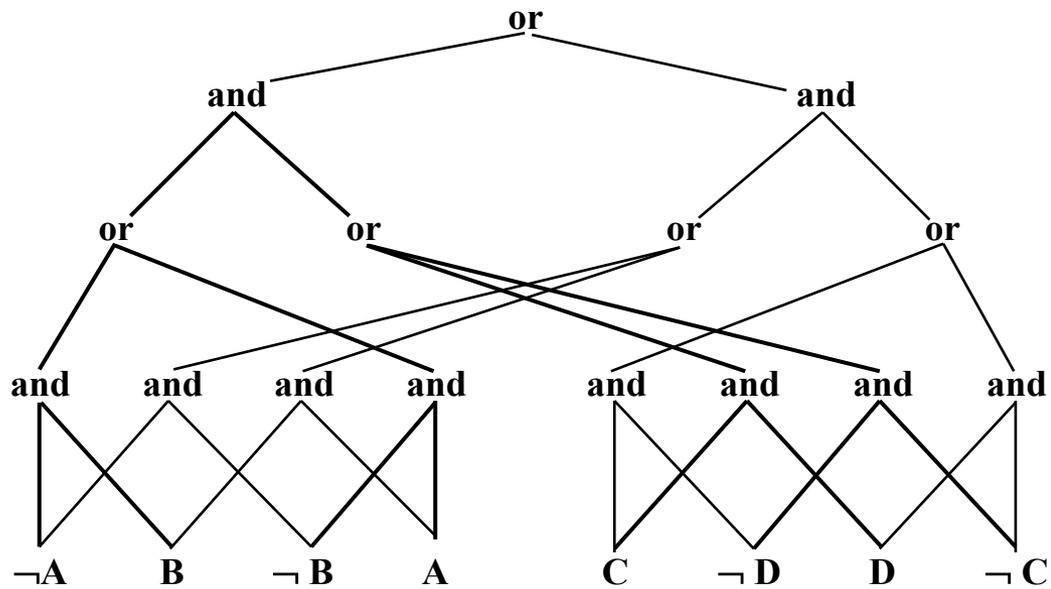
Size $p(n)$

Size n

L1 is more succinct than L2

$$\mathbf{L1} < \mathbf{L2}$$

Odd Parity Function



Tractability & Succinctness

NNF

Tractable Operations

DNNF

decomposability

**Diagnosis,
Non-mon**

d-DNNF

determinism

**Probabilistic
reasoning**

FBDD

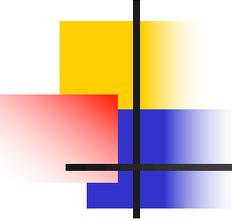
decision

OBDD

ordering

Space Efficiency (succinctness)

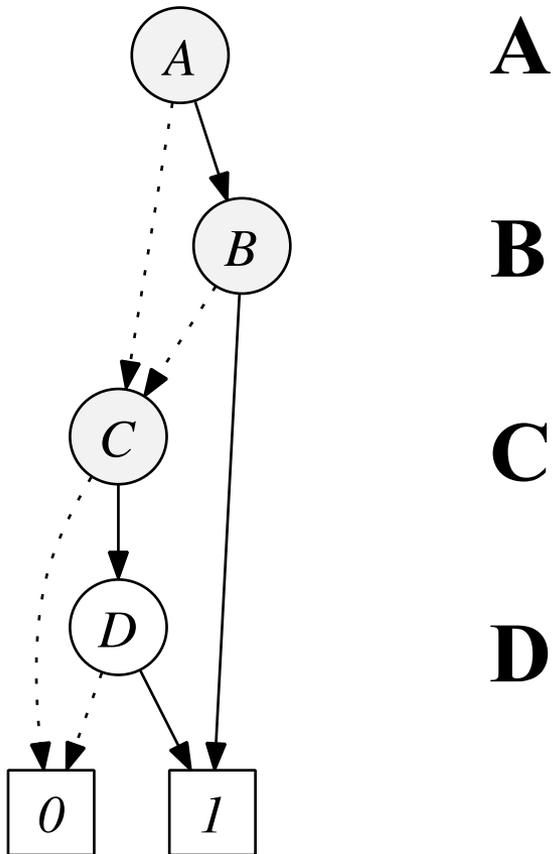




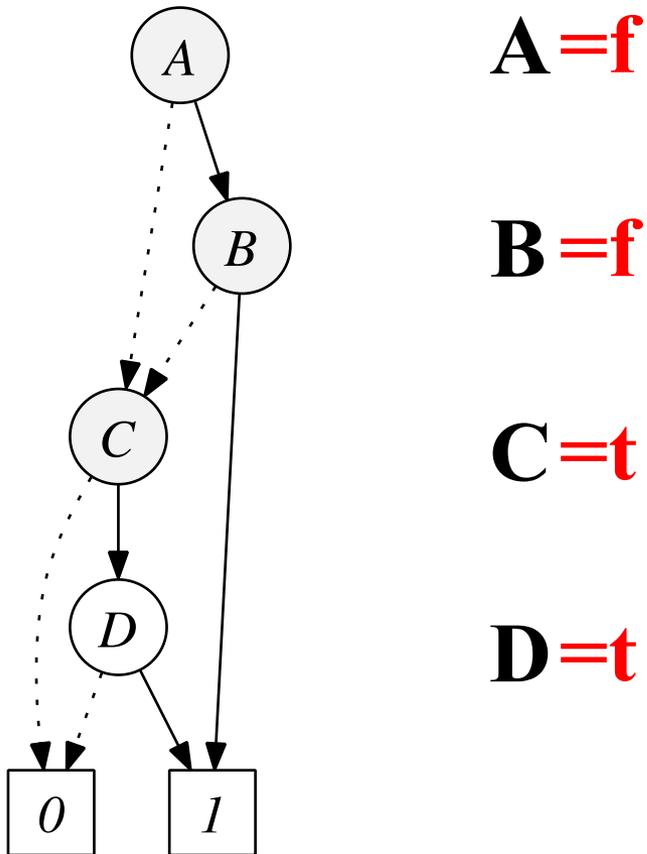
Separating Functions

- OBDD/FBDD:
 - Hidden weighted bit function $\text{hwb}(x_1, \dots, x_n)$
- DNNF/DNF:
 - odd parity function $\text{parity}(x_1, \dots, x_n)$
- DNNF/OBDD:
 - Distinct integers function $\text{distinct}(x_1, \dots, x_n)$

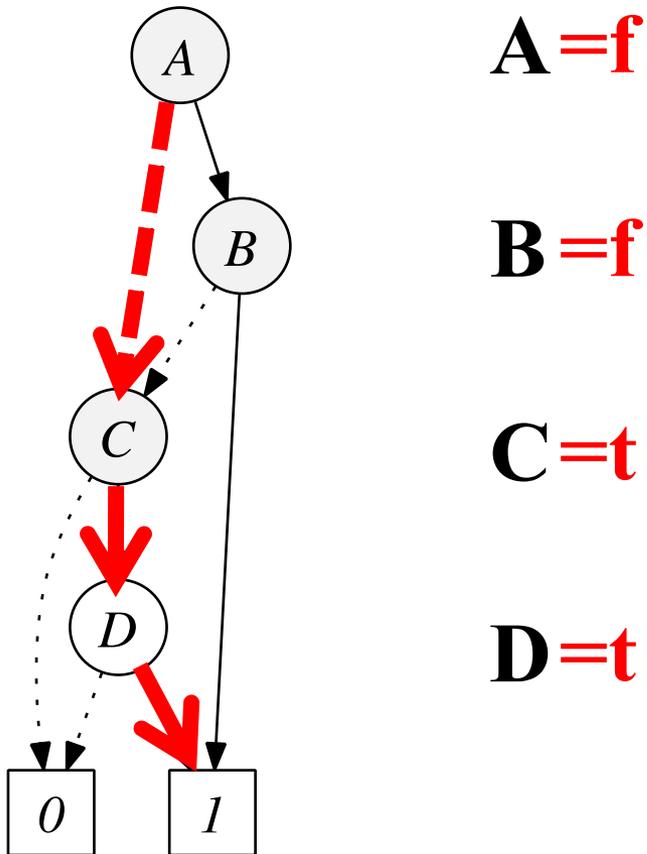
From OBDD to SDD



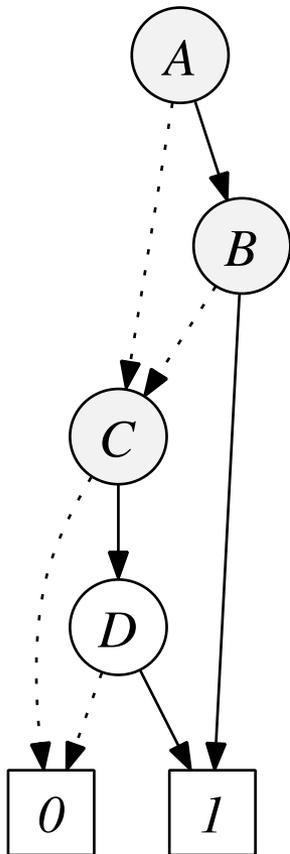
From OBDD to SDD



From OBDD to SDD



From OBDD to SDD

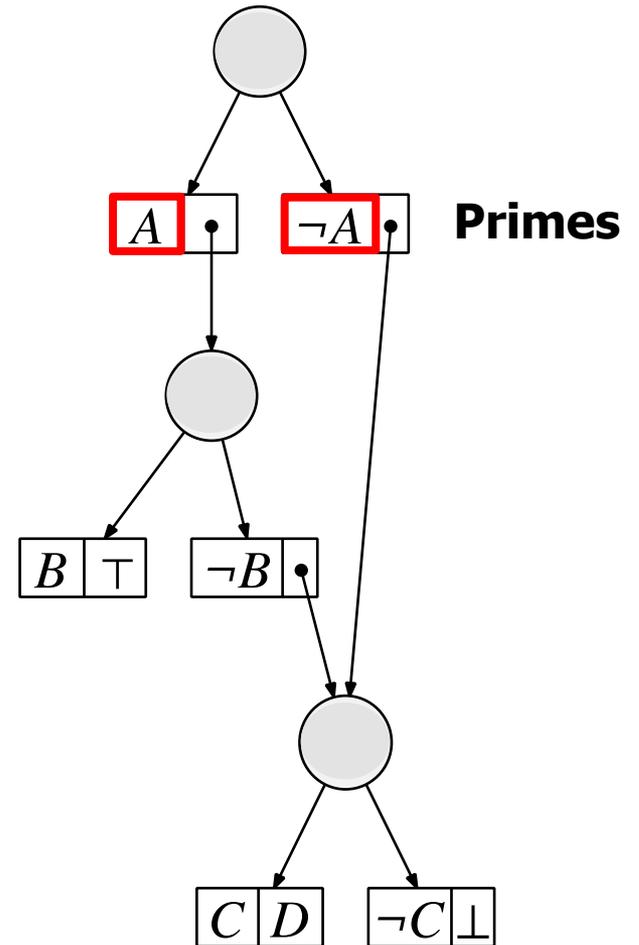


A

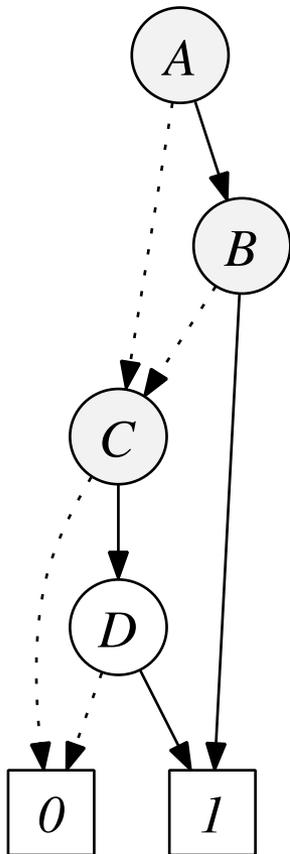
B

C

D



From OBDD to SDD

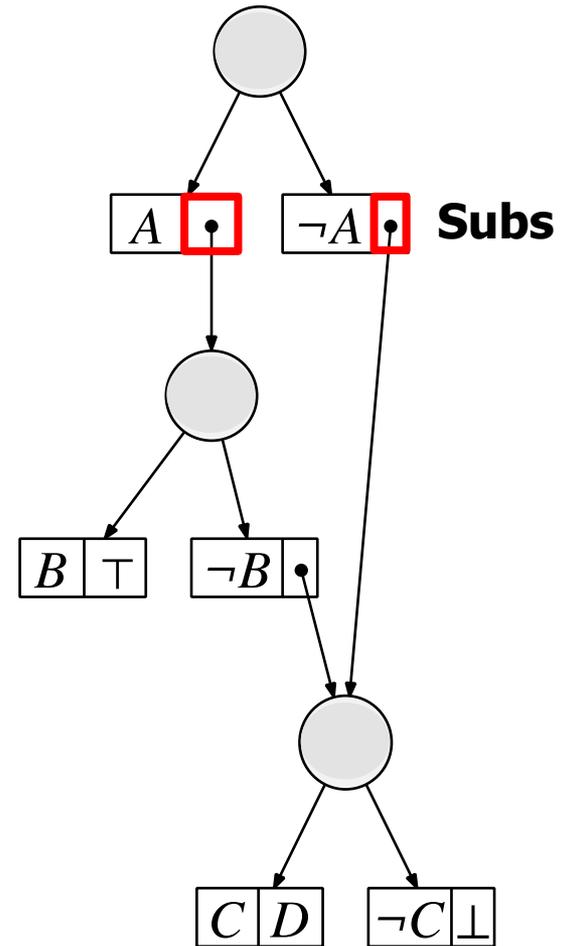


A

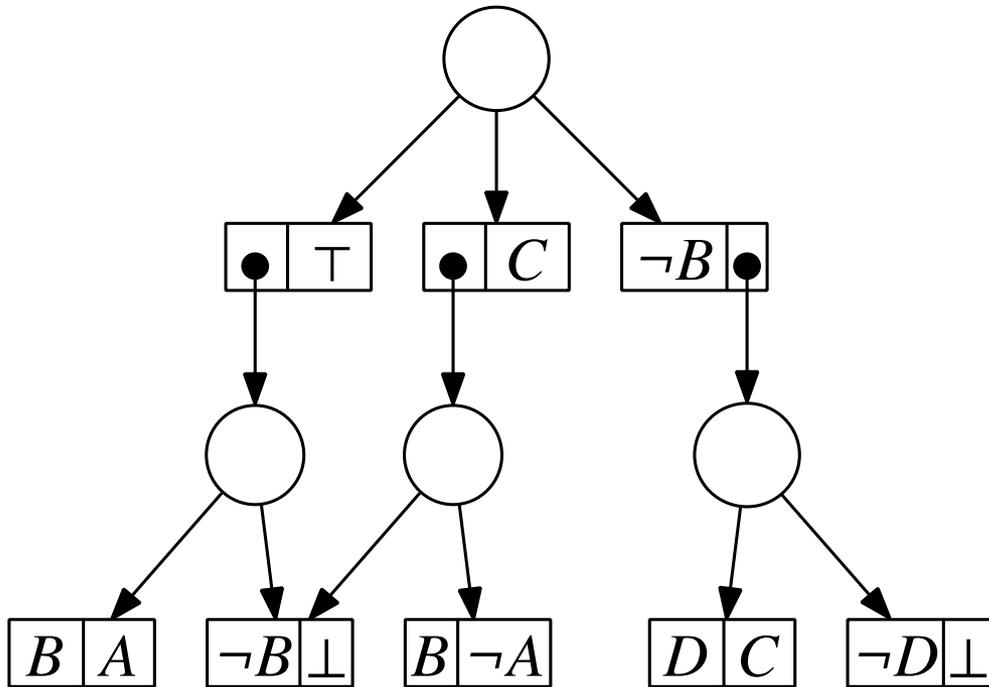
B

C

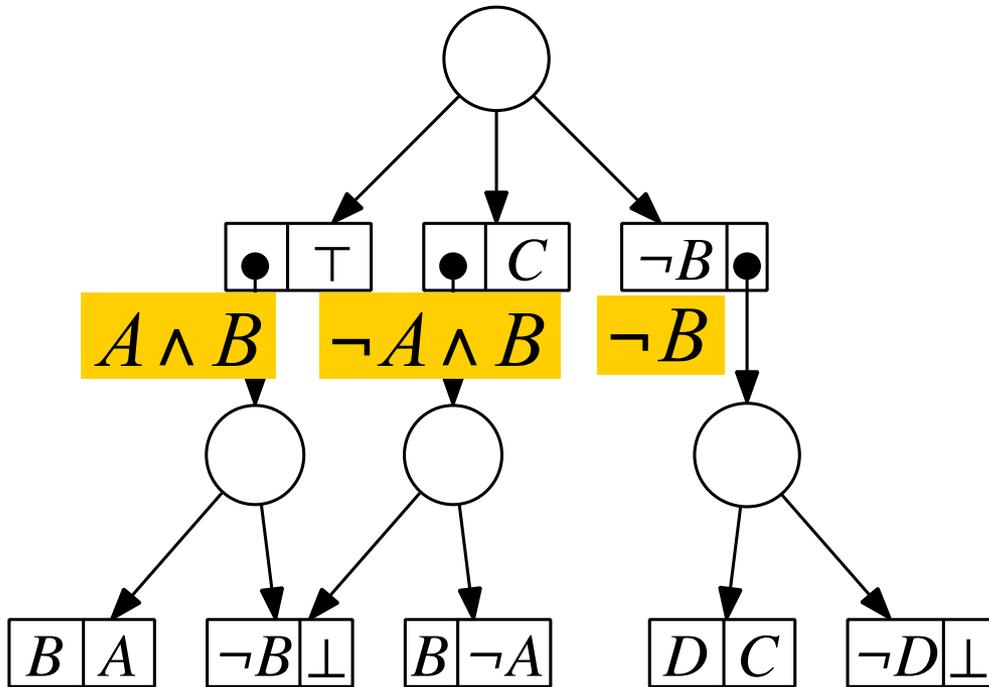
D



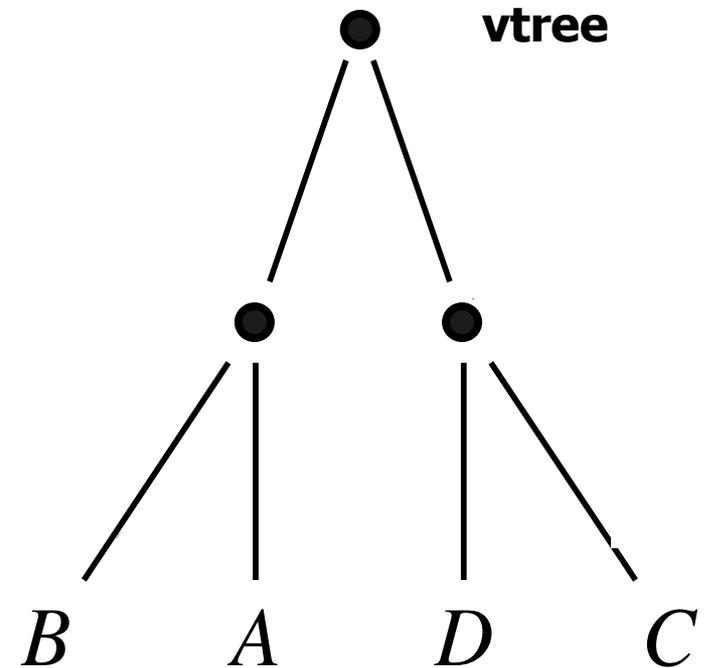
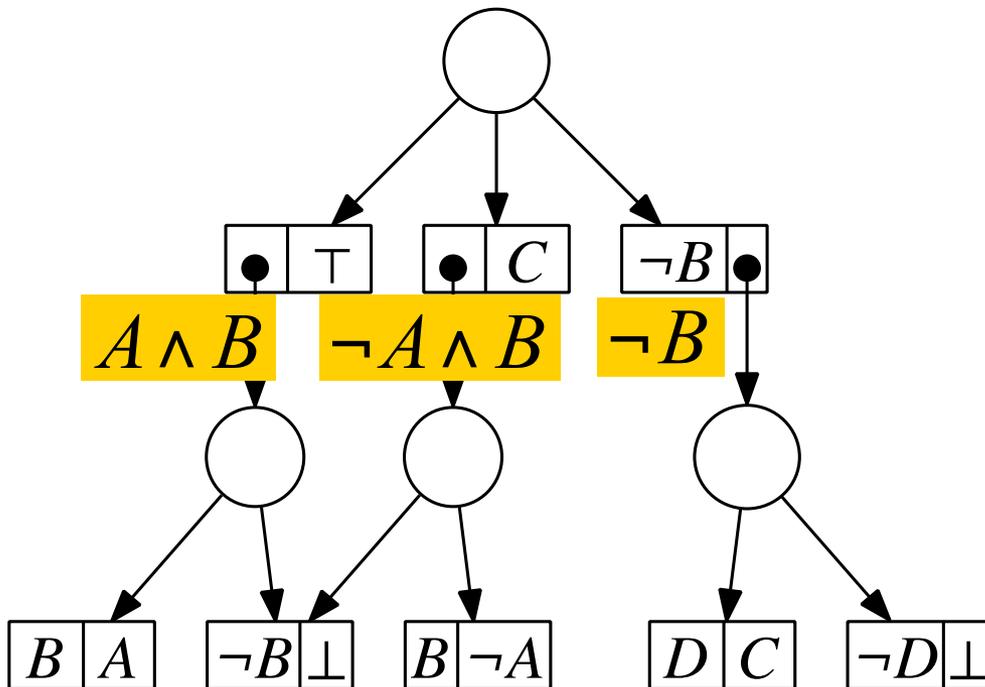
SDD: Sentential Decision Diagram



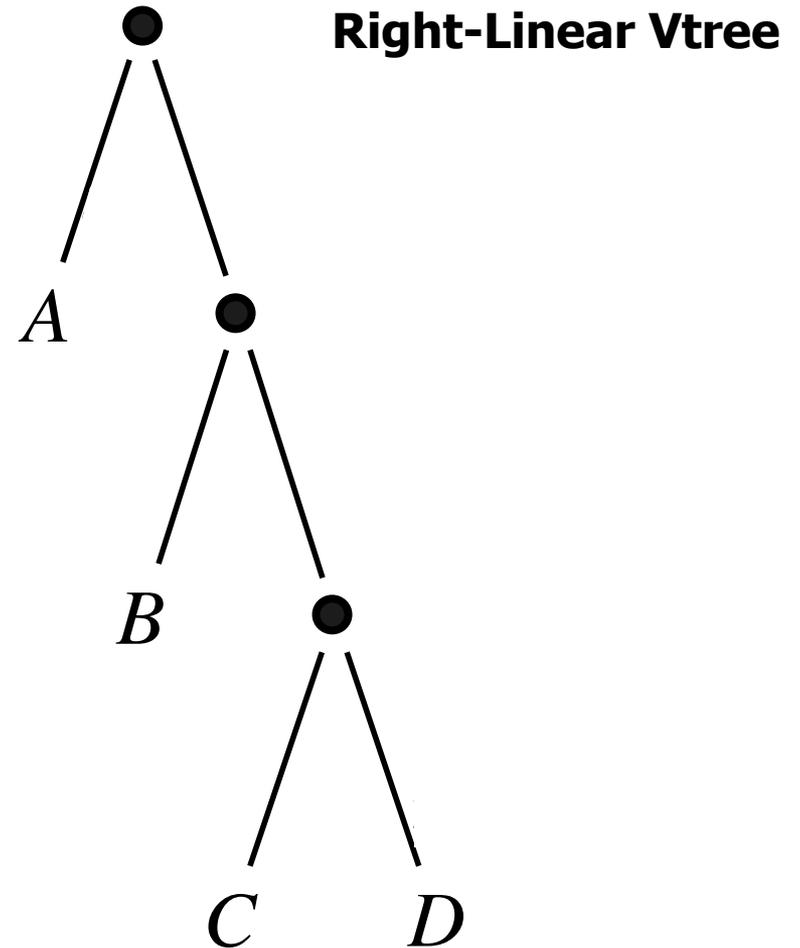
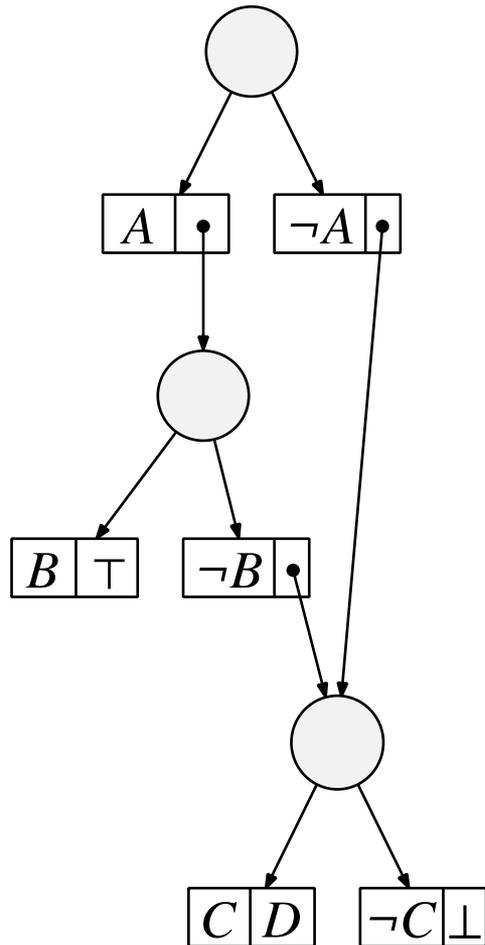
SDD: Sentential Decision Diagram

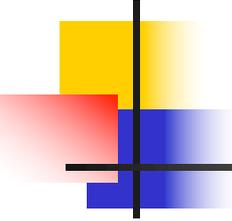


SDD: Sentential Decision Diagram



OBDD as SDD





X-Partitions: The insight!

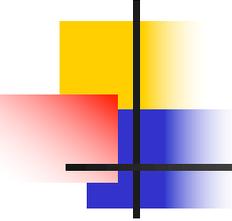
- Write function $f(\mathbf{X}, \mathbf{Y})$ as

$$h_1(\mathbf{X}) g_1(\mathbf{Y}) + \dots + h_n(\mathbf{X}) g_n(\mathbf{Y})$$

such that:

$h_1(\mathbf{X}), \dots, h_n(\mathbf{X})$ is a partition

(mutually exclusive and exhaustive; $h_i \neq \text{false}$)



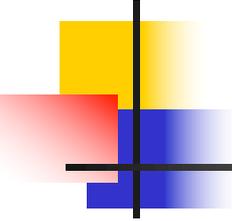
X-Partitions

- **X**-partition

$$h_1(\mathbf{X}) g_1(\mathbf{Y}) + \dots + h_n(\mathbf{X}) g_n(\mathbf{Y})$$

written as $\{ (h_1, g_1), \dots, (h_n, g_n) \}$

- We call h_i primes, g_i subs

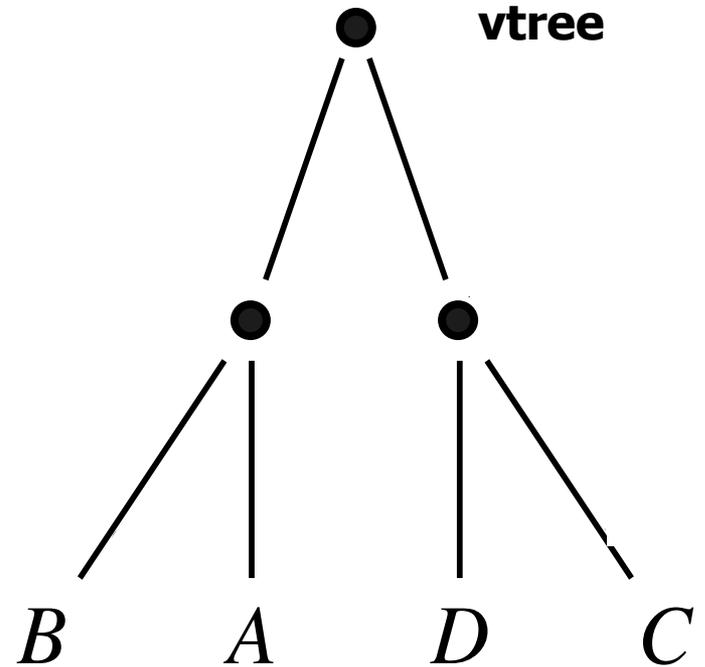
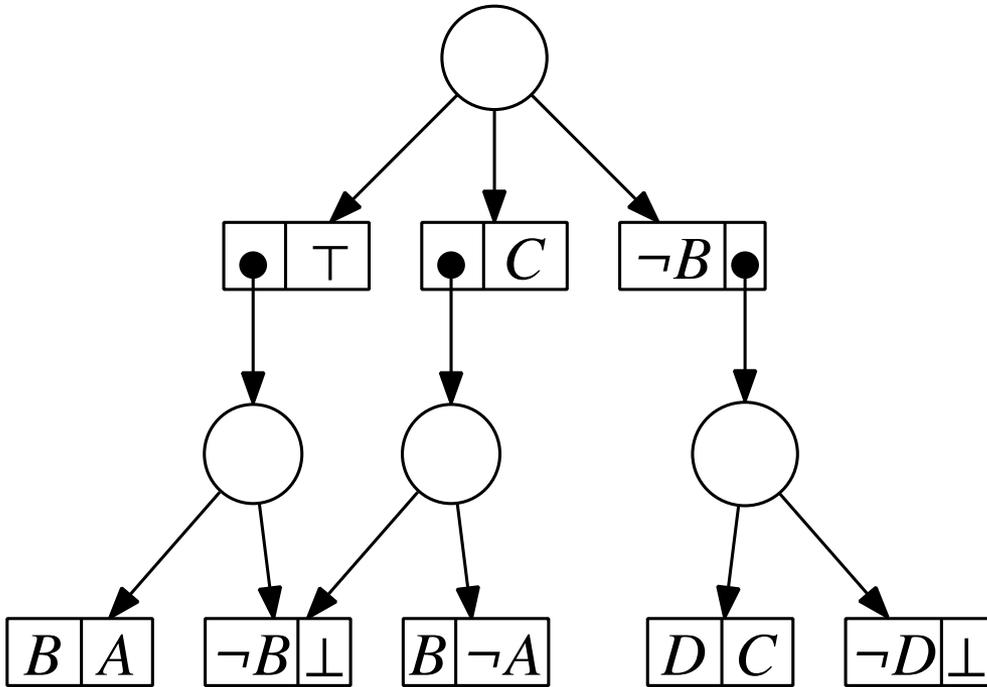


Compression & Canonicity

- **X**-partition is compressed if no equal subs
 $\{ (h_1, g_1), (h_2, g_2) \dots, (h_n, g_n) \}$
- If equal subs $g_1 = g_2$, compress to:
 $\{ (h_1 + h_2, g_1), \dots, (h_n, g_n) \}$
- Every function $f(\mathbf{X}, \mathbf{Y})$ has a unique
compressed **X**-partition

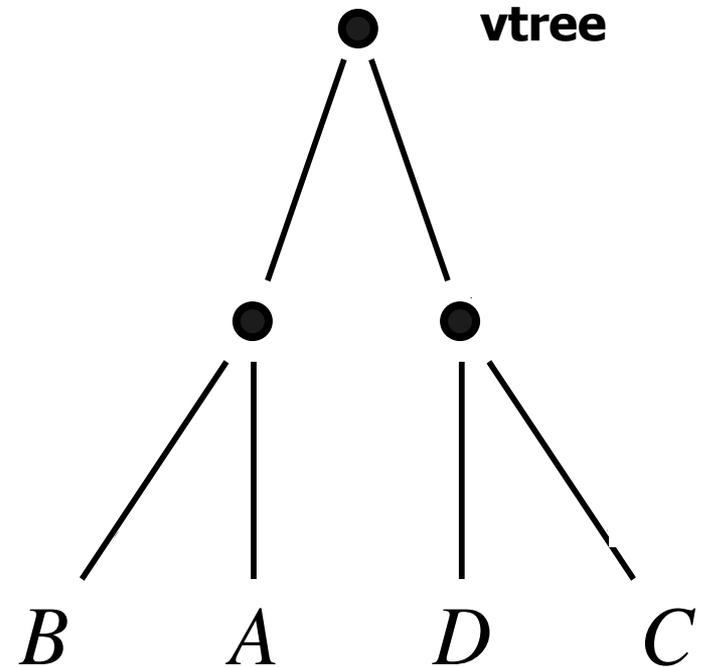
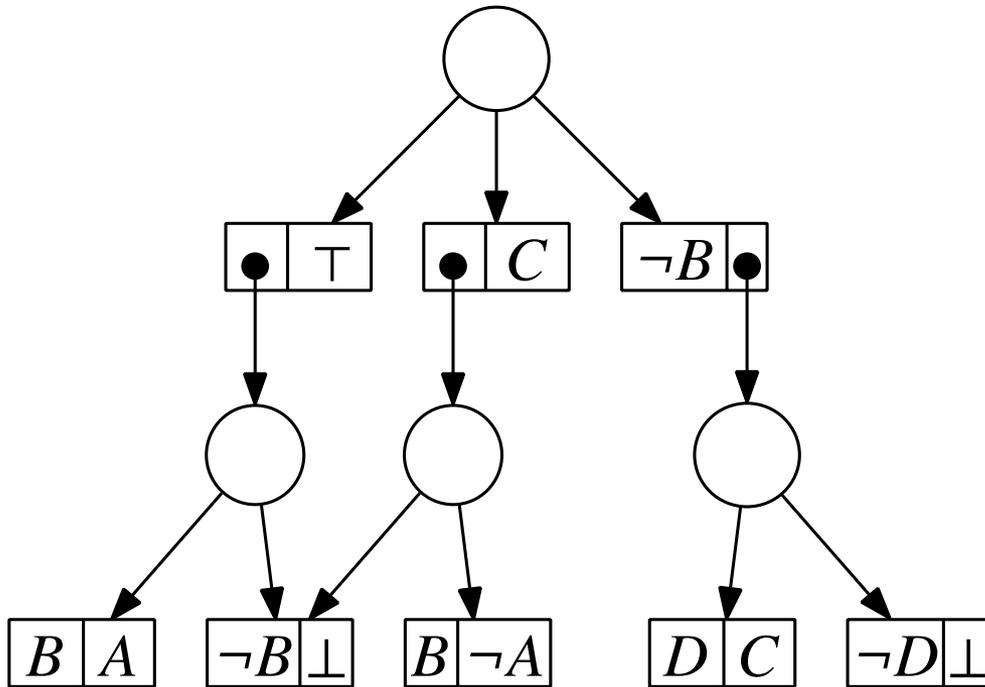
SDD

$$f = (A \wedge B) \vee (B \wedge C) \vee (C \wedge D)$$



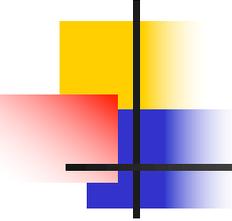
SDD

$$f = (A \wedge B) \vee (B \wedge C) \vee (C \wedge D)$$



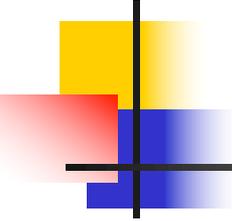
Compressed X-Partition of function f with $X = \{A, B\}$

$$\{(A \wedge B, true), (\neg A \wedge B, C), (\neg B, C \wedge D)\}$$



OBDDs are SDDs

- When $\mathbf{X}=\{X\}$ (single variable), an \mathbf{X} -partition corresponds to a Shannon decomposition
- Primes: $X, \neg X$
- Subs: $f \mid X, f \mid \neg X$
- \mathbf{X} -partition: $\{(X, f \mid X), (\neg X, f \mid \neg X)\}$



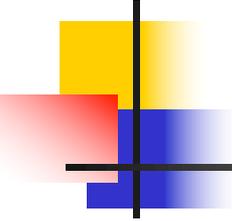
Polytime Apply Operation

- **X**-partition of $f(\mathbf{X}, \mathbf{Y})$: $\{(p_1, q_1), \dots, (p_n, q_n)\}$
- **X**-partition of $g(\mathbf{X}, \mathbf{Y})$: $\{(r_1, s_1), \dots, (r_m, s_m)\}$

- **X**-partition of

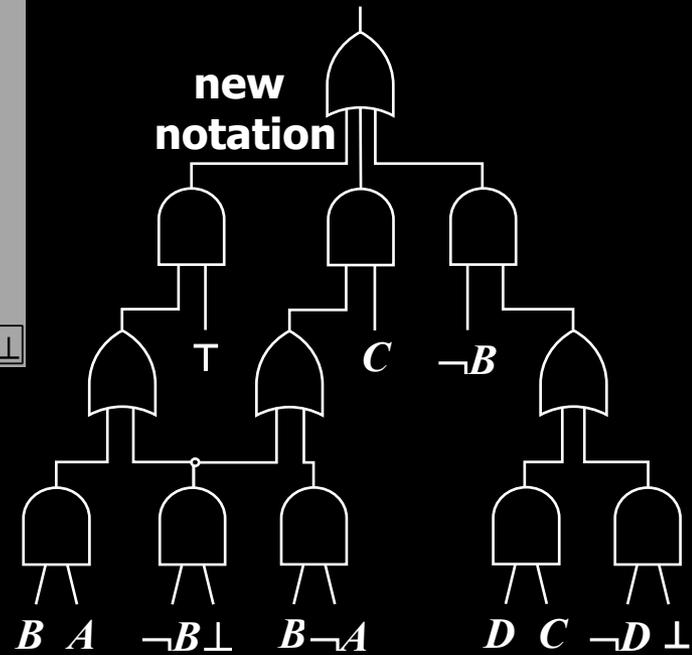
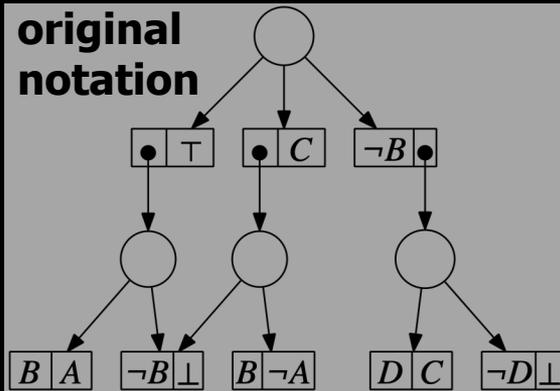
$$f \circ g = \{(p_i \wedge r_j, q_i \circ s_j) \mid p_i \wedge r_j \neq \text{false}\}$$

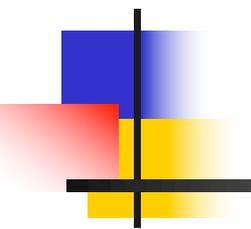
- Result may not be compressed



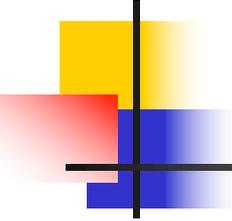
SDD vs OBDD

- SDD a strict superset of OBDD:
 - Characterized by trees, which include orders
 - Branch over sentences, which include literals
- SDDs maintain key properties of OBDDs:
 - Canonical when compressed
 - Polytime Apply Operation (no compression)
- SDDs: treewidth, OBDD: pathwidth
- SDDs more succinct than OBDDs

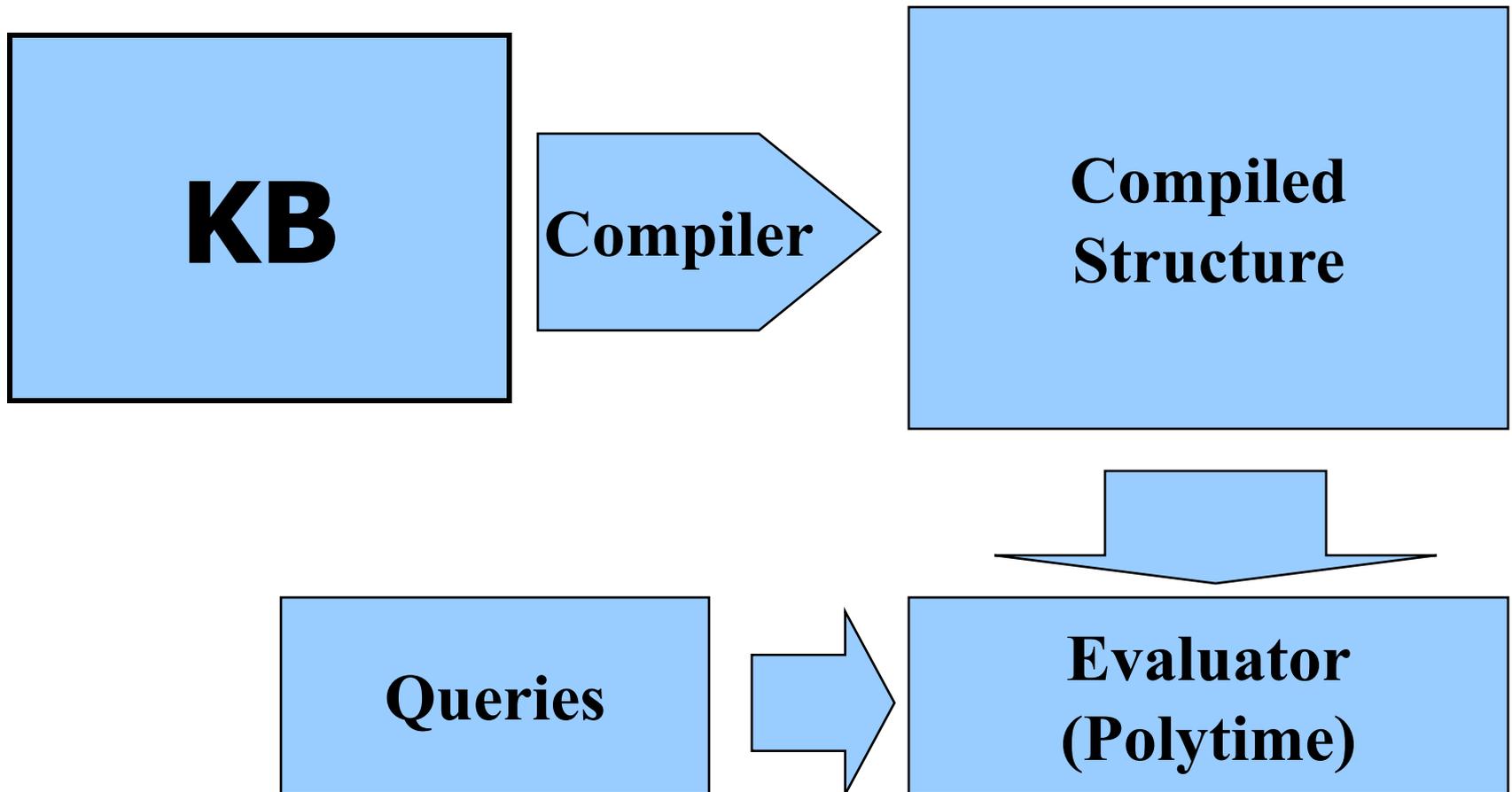


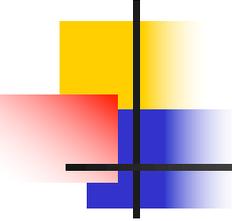


Operations



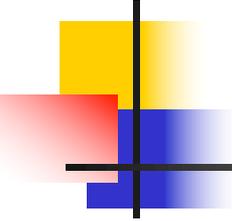
Knowledge Compilation





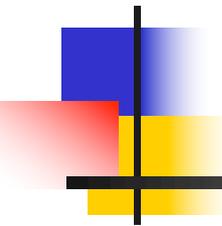
Queries

- **Consistency (CO)**
- **Validity (VA)**
- **Sentential entailment (SE)**
- **Clausal entailment (CE): KB implies clause**
- **Implicant testing (IP): term implies KB**
- **Equivalence testing (EQ)**
- **Model Counting (CT)**
- **Model enumeration (ME)**



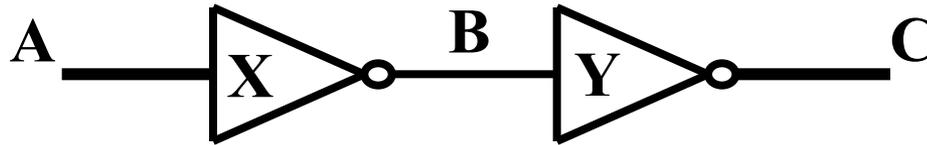
Transformations

- **Projection (existential quantification)**
- **Conditioning**
- **Conjoin**
- **Disjoin**
- **Negate**



Decomposability

Example Knowledge Base

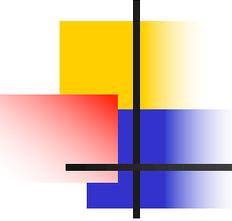


$A \ \& \ okX \Rightarrow \sim B$

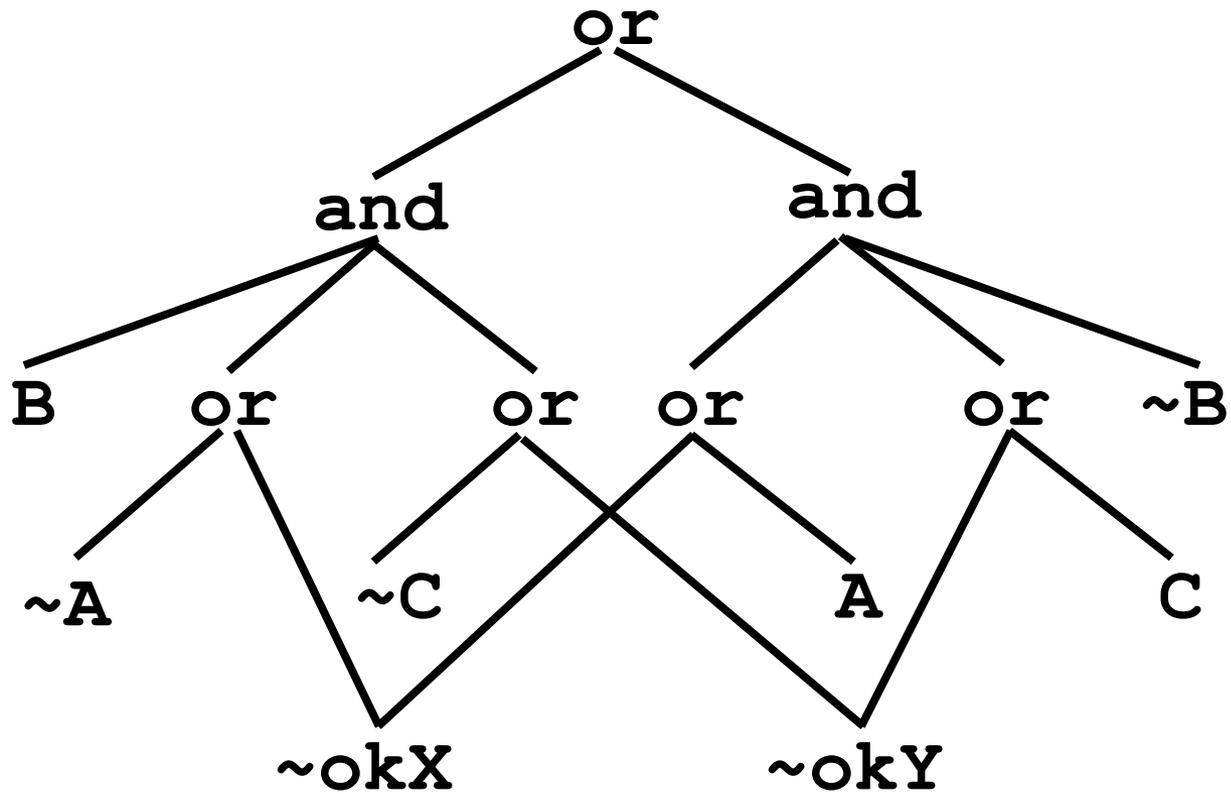
$\sim A \ \& \ okX \Rightarrow B$

$B \ \& \ okY \Rightarrow \sim C$

$\sim B \ \& \ okY \Rightarrow C$



Decomposable

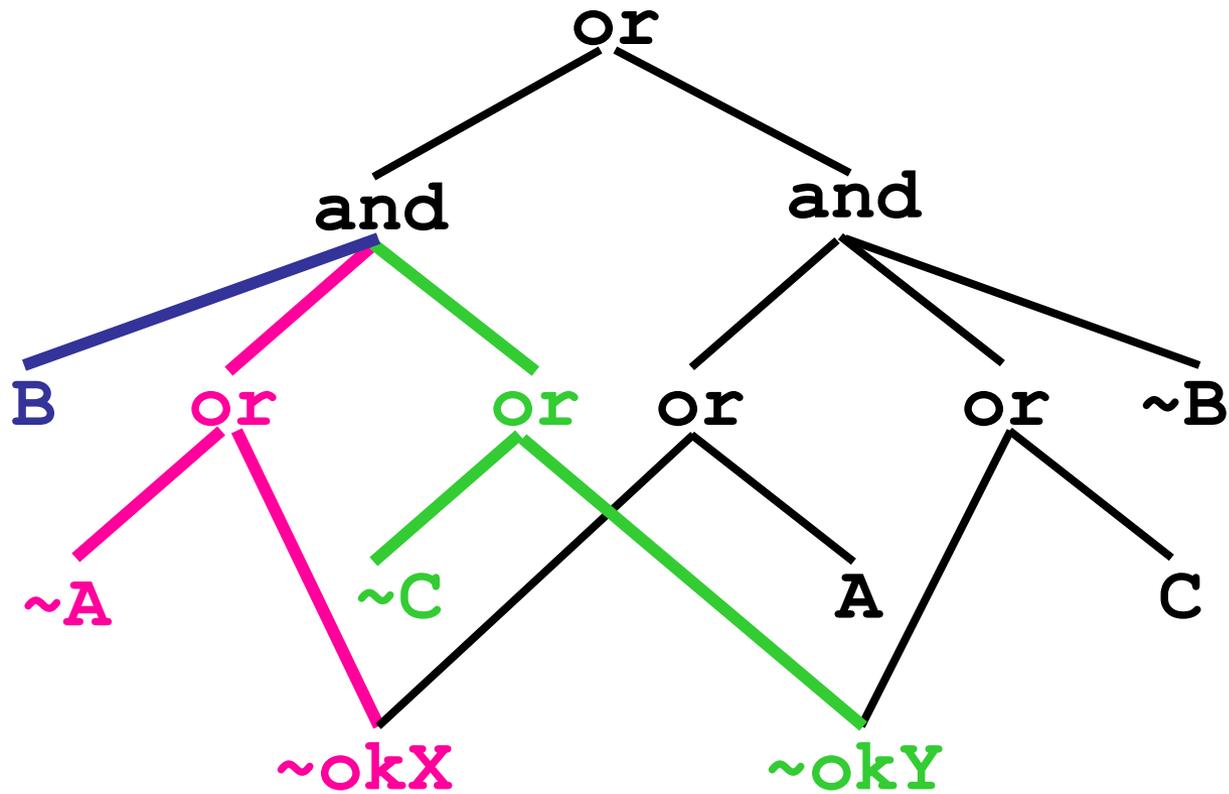


Decomposable

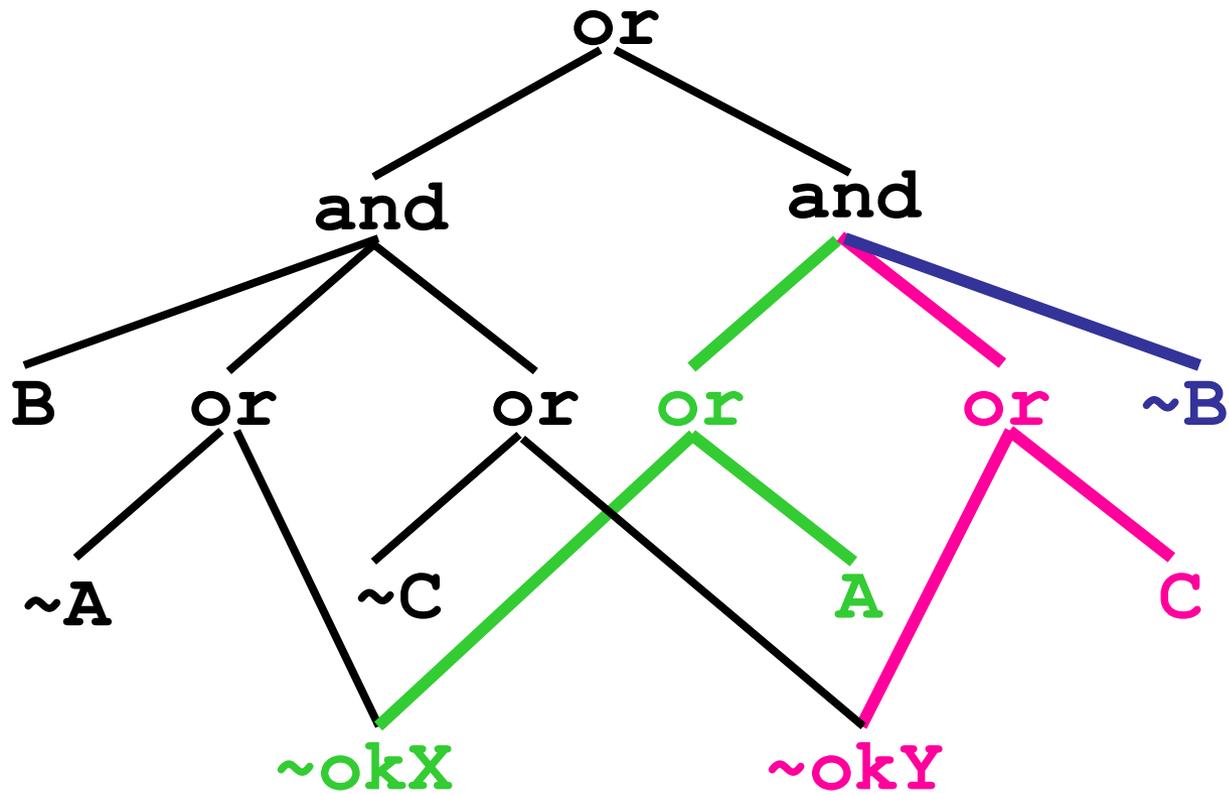
B

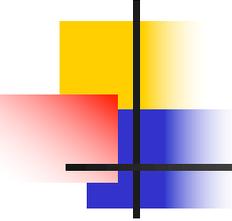
A, okX

C, okY



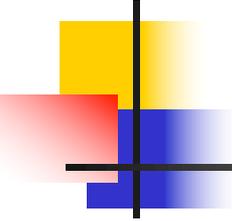
Decomposable



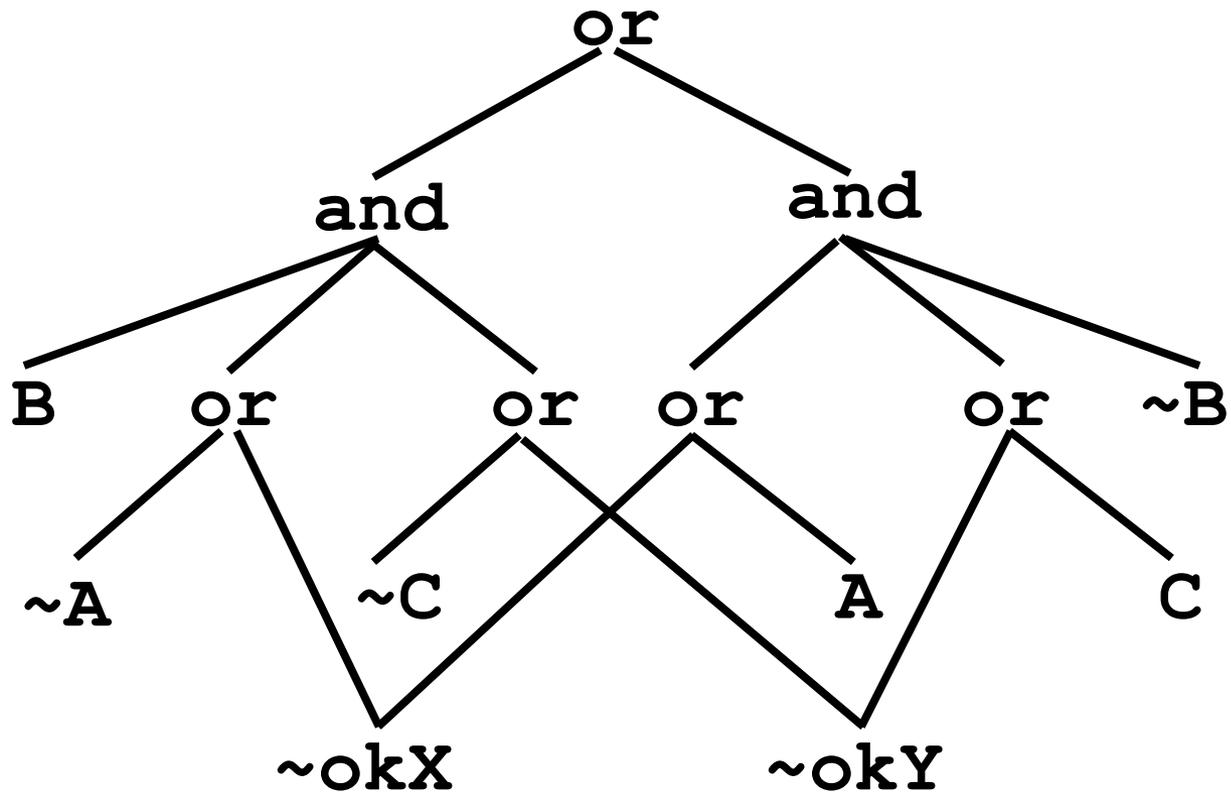


Satisfiability

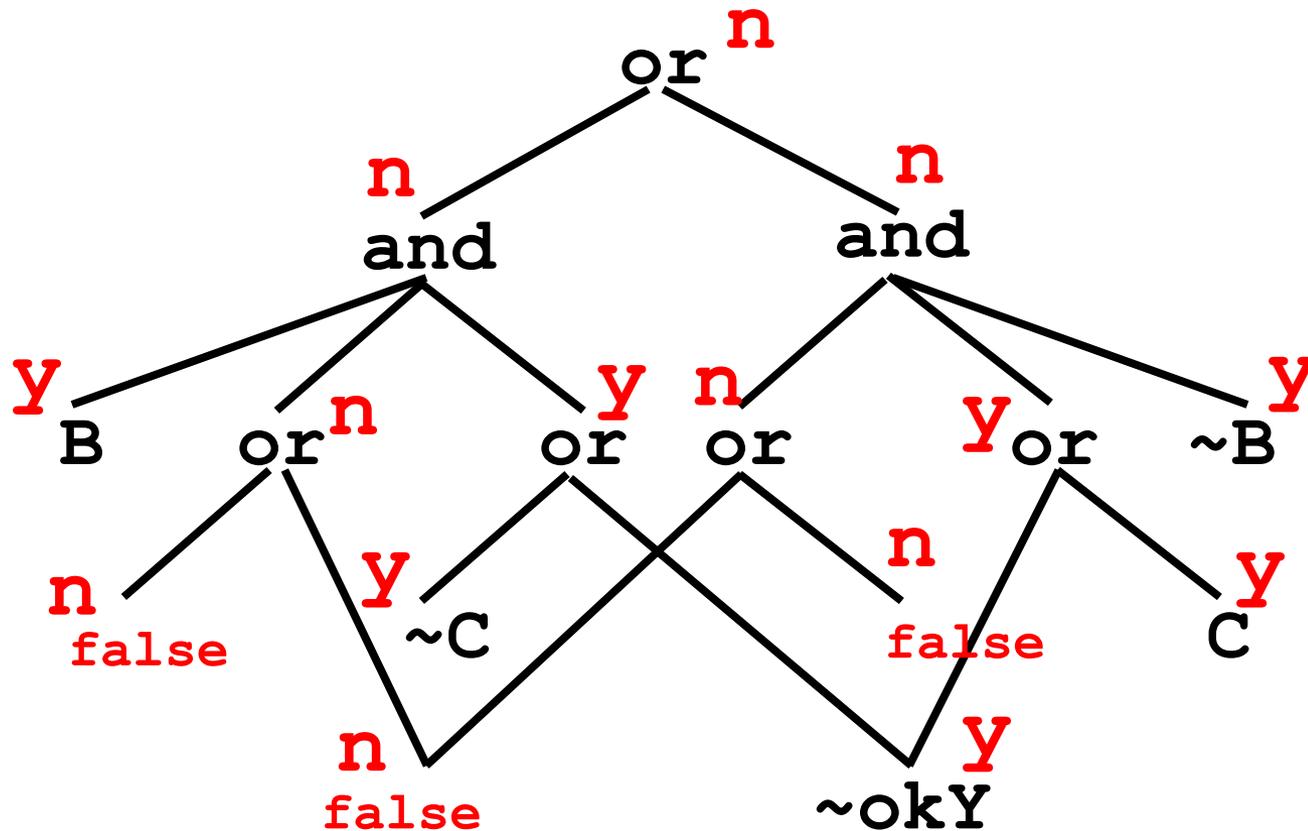
- $\text{SAT}(A \text{ or } B)$ iff $\text{SAT}(A)$ or $\text{SAT}(B)$
- $\text{SAT}(A \text{ and } B)$ iff $\text{SAT}(A)$ and $\text{SAT}(B)$
- $\text{SAT}(X)$ is true
- $\text{SAT}(\sim X)$ is true
- $\text{SAT}(\text{True})$ is true
- $\text{SAT}(\text{False})$ is false

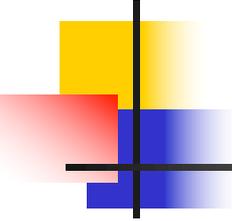


Satisfiability



Satisfiability



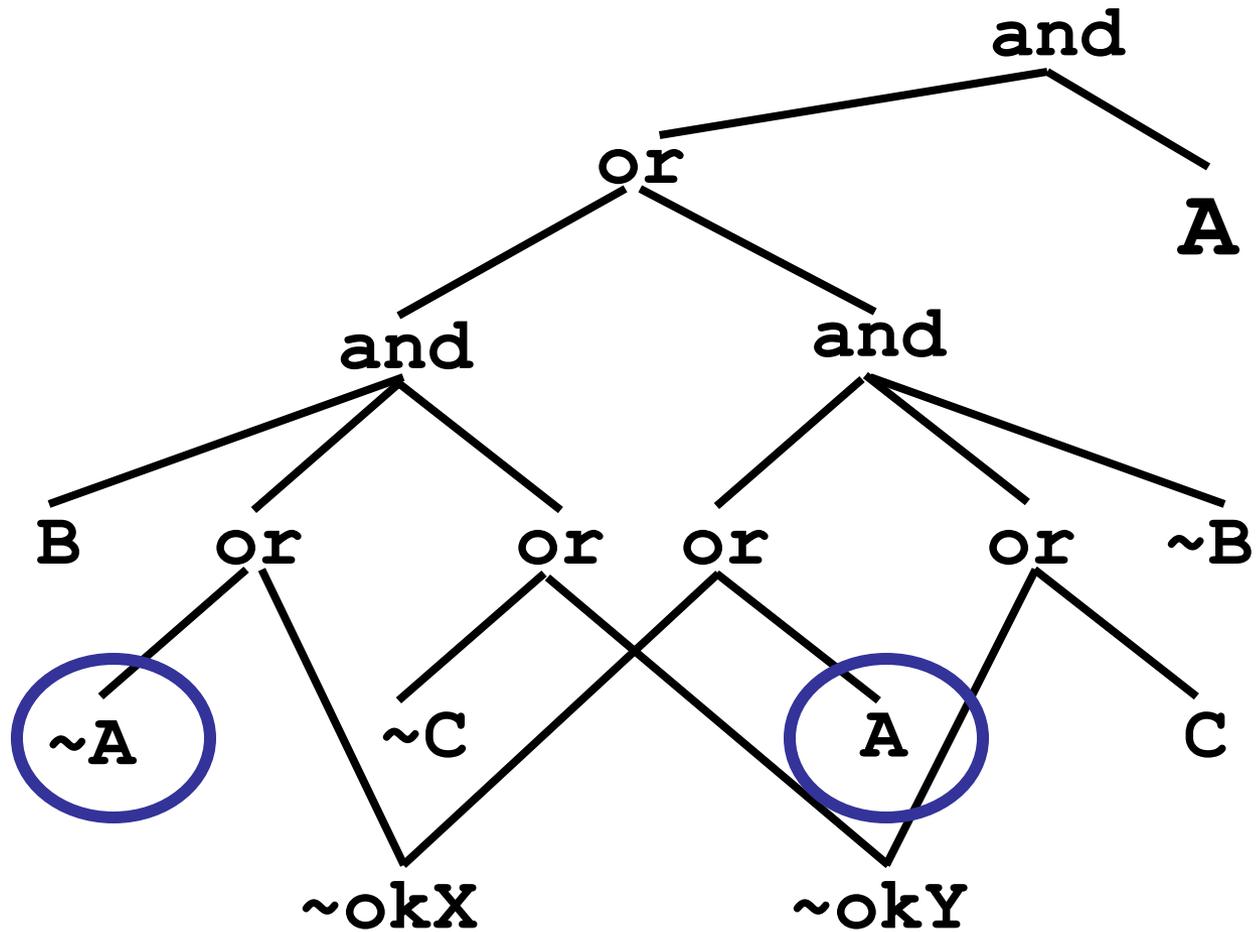


Clausal Entailment

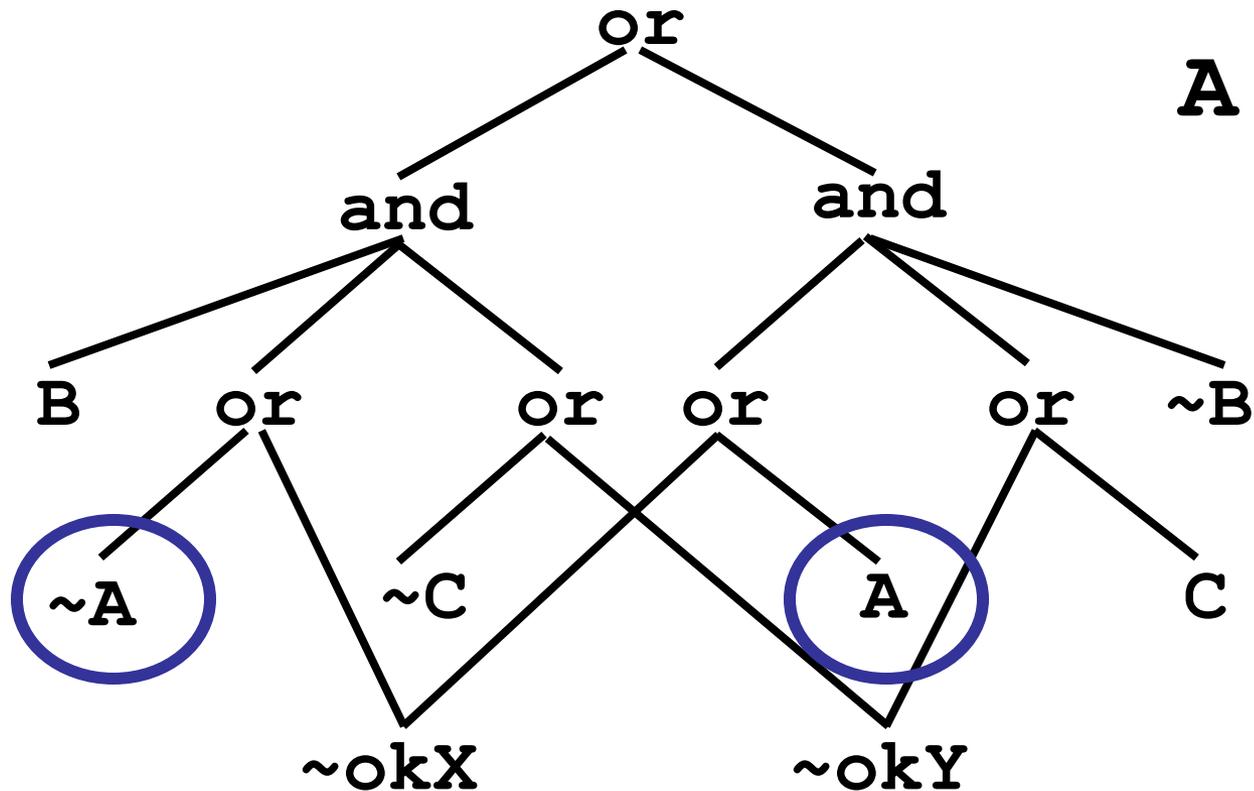
KB *entails* $L_1 \vee L_2 \vee \dots \vee L_n$?

KB $\& \sim L_1 \& \sim L_2 \& \dots \& \sim L_n$ *SAT?*

Literal Conjoin

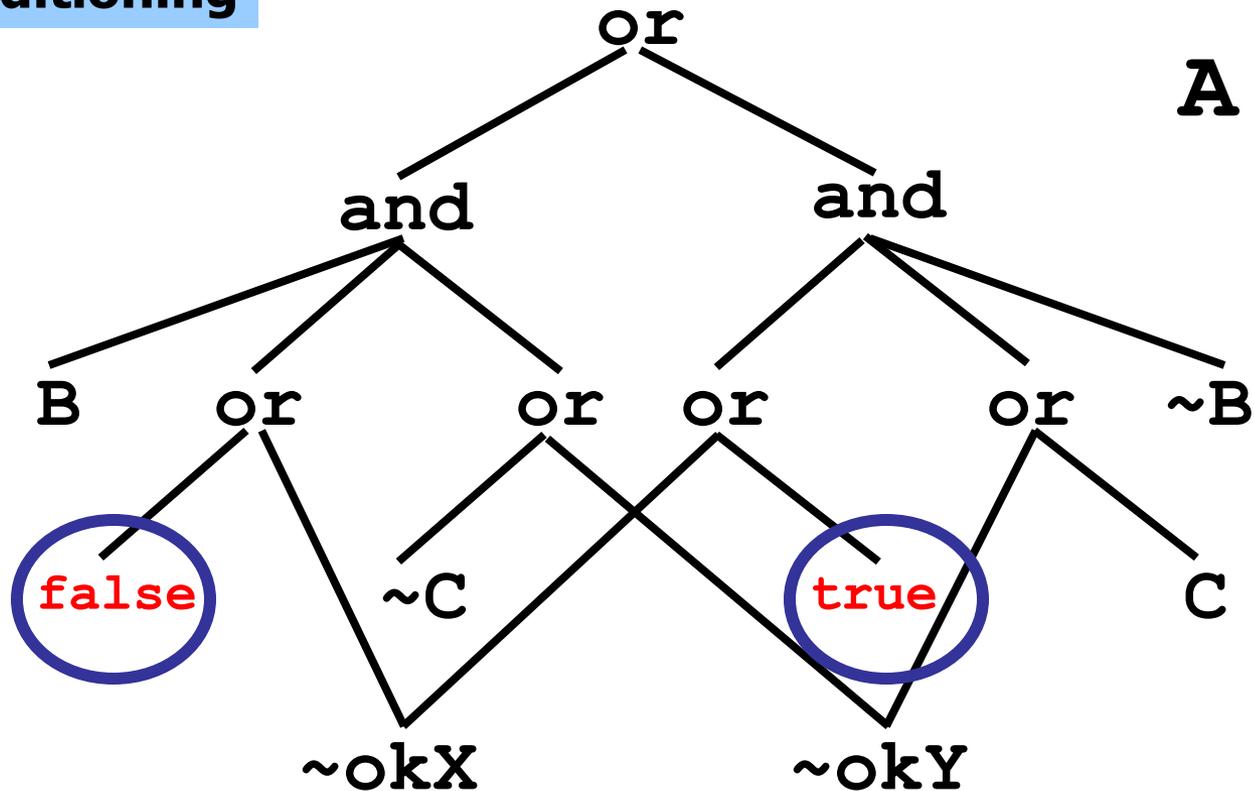


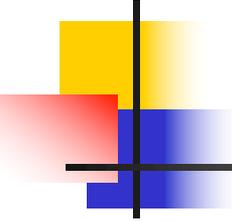
Literal Conjoin



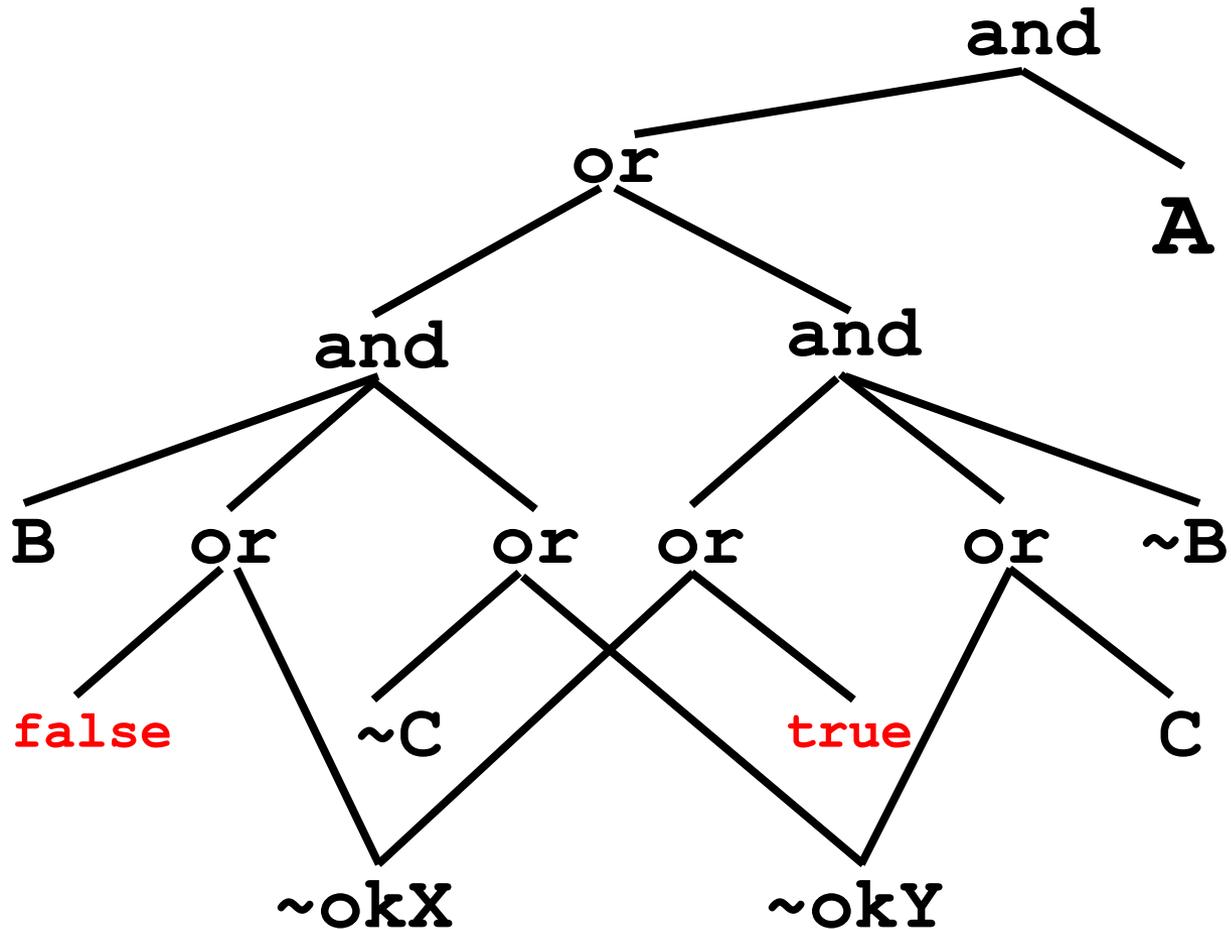
Literal Conjoin

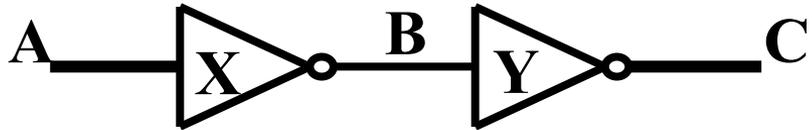
Conditioning



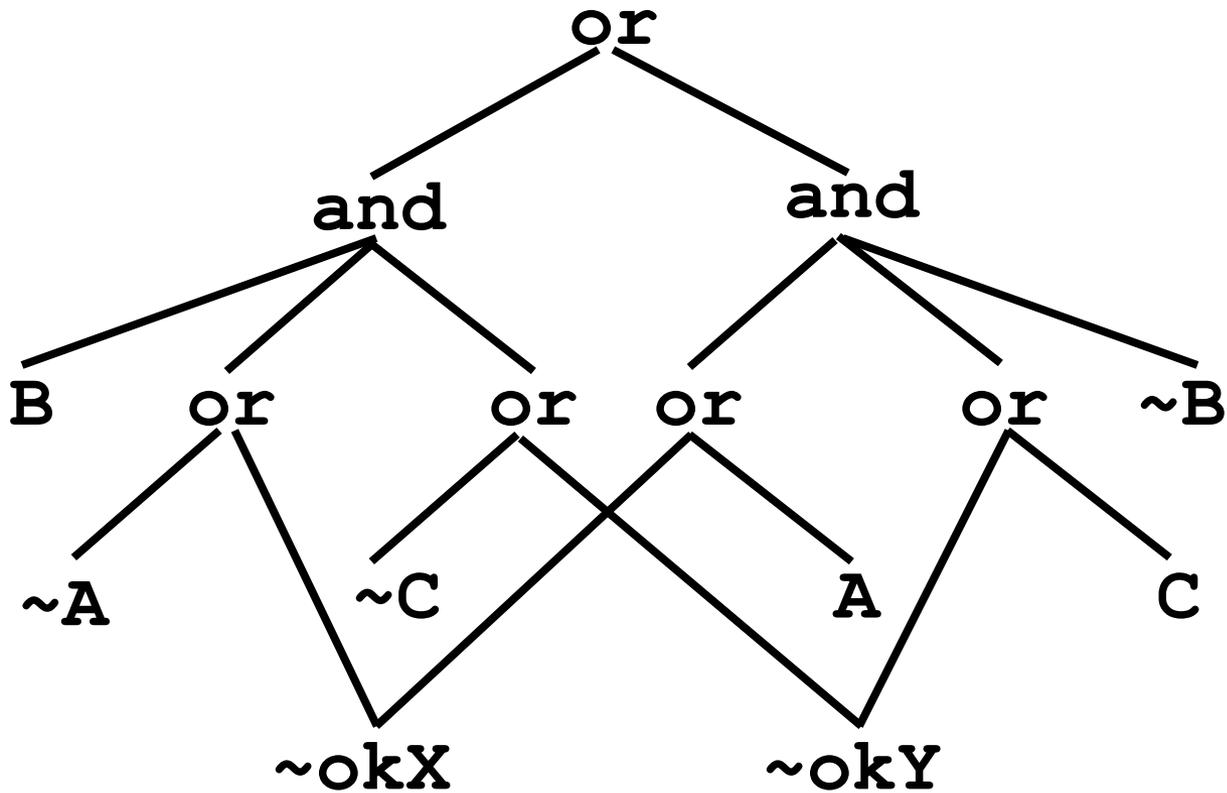


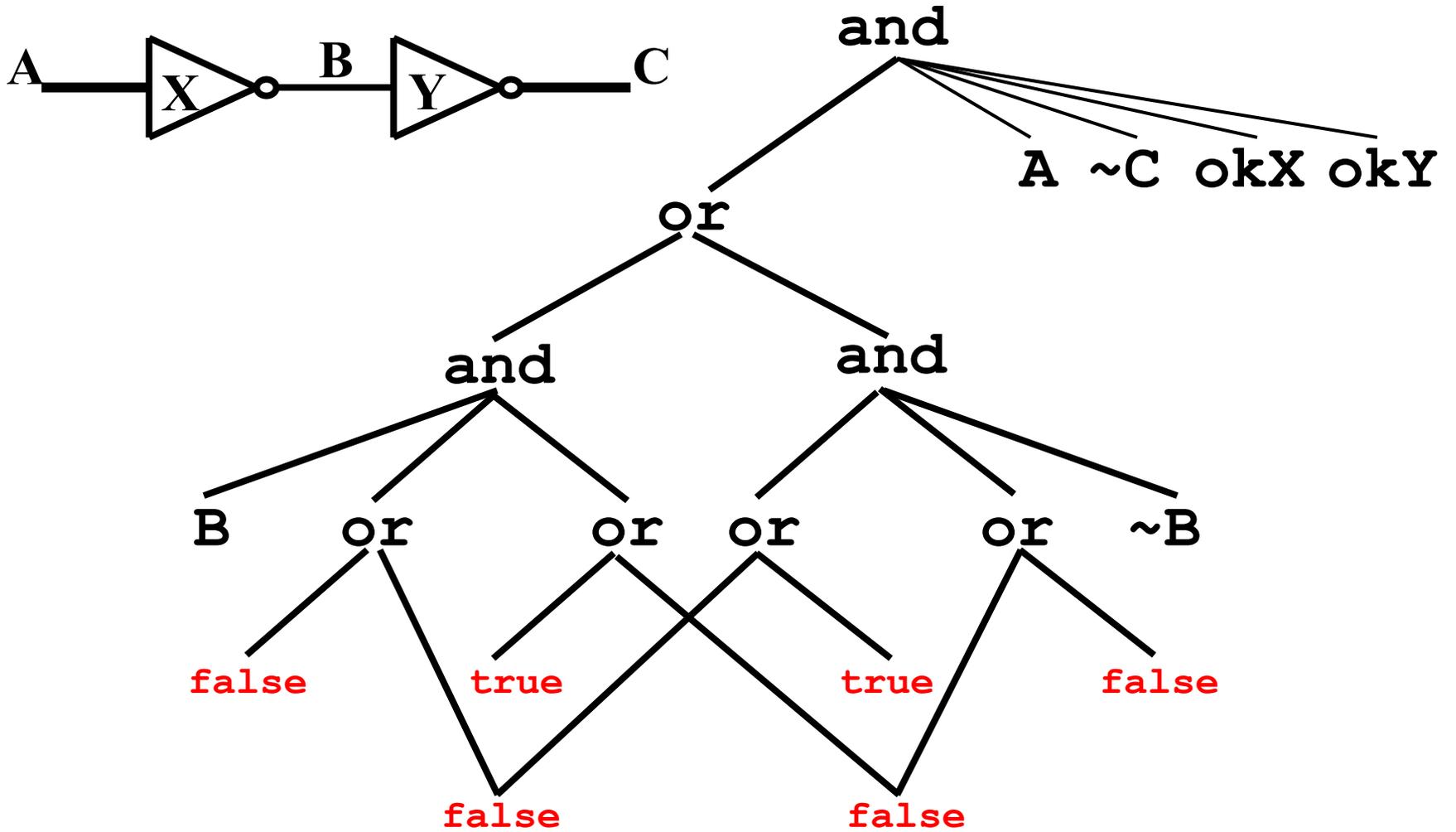
Literal Conjoin

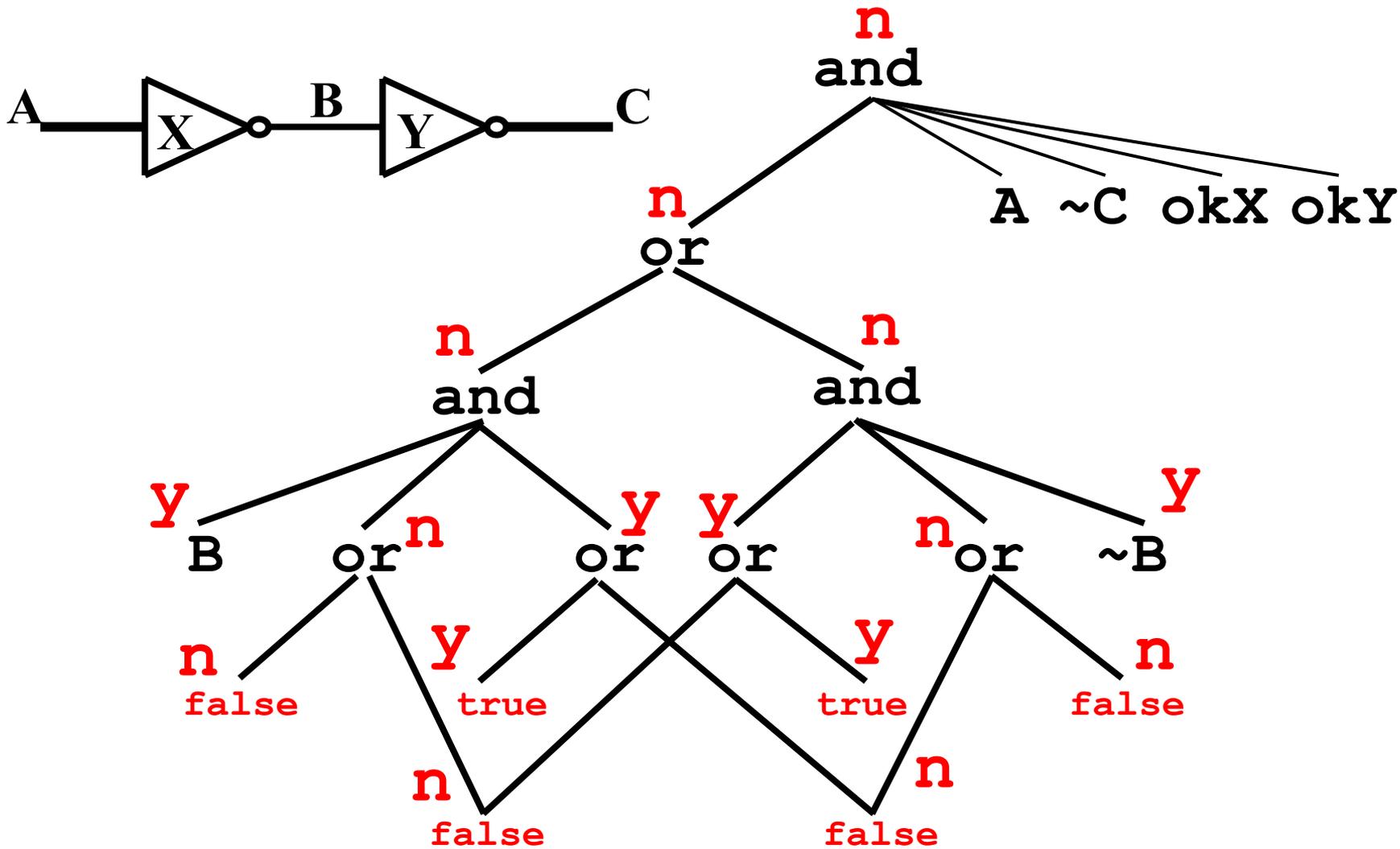


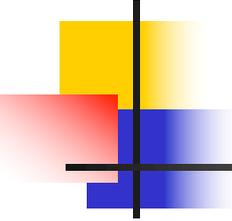


A ~C okX okY









Projection: Existential Quantification

Knowledge Base

$$\Delta = A \Rightarrow B, B \Rightarrow C, C \Rightarrow D$$

Existentially quantifying B,C

Forgetting B,C

Projecting on A,D

$$(\exists B \exists C \Delta) = A \Rightarrow D$$

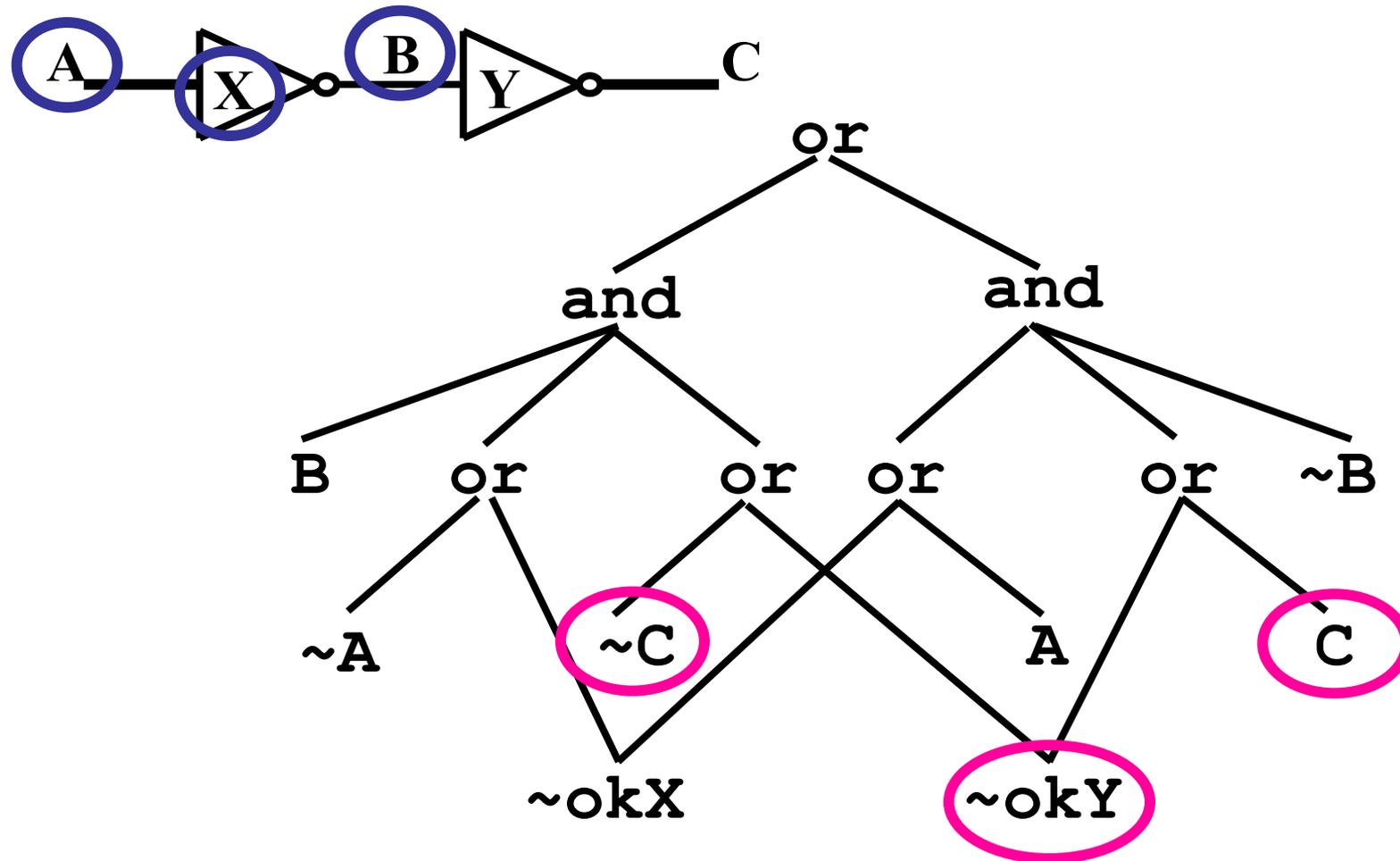
Projection: Existential Quantification

Formal Definition

$$\exists X \Delta = (\Delta \mid X) \vee (\Delta \mid \neg X)$$

- **If Knowledge base is a CNF:**
 - **Close under resolution**
 - **Remove all clauses that mention X**

Projection: Existential Quantification



Projection: Existential Quantification

A

$$\exists X (\Delta \wedge \Gamma)$$

$$= ((\Delta \wedge \Gamma) | X) \vee ((\Delta \wedge \Gamma) | \neg X)$$

$$= ((\Delta | X) \wedge (\Gamma | X)) \vee ((\Delta | \neg X) \wedge (\Gamma | \neg X))$$

$$= (\Delta \wedge (\Gamma | X)) \vee (\Delta \wedge (\Gamma | \neg X))$$

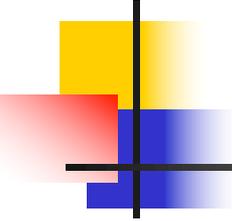
$$= \Delta \wedge ((\Gamma | X) \vee (\Gamma | \neg X))$$

$$= \Delta \wedge (\exists X \Gamma)$$

$\sim \text{ok} X$

true

$\sim B$
B



Minimum Cardinality

$$A \ \& \ okX \Rightarrow \sim B$$

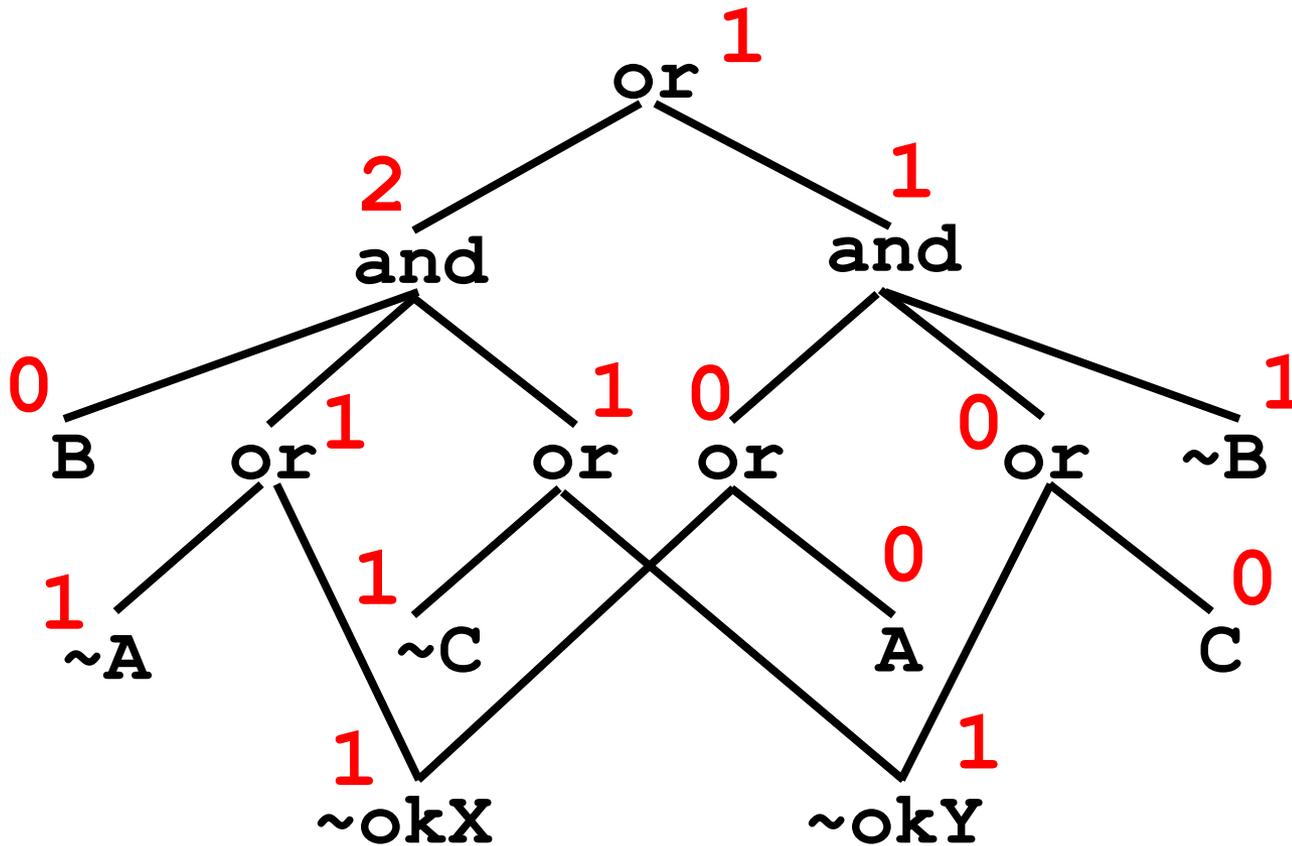
$$\sim A \ \& \ okX \Rightarrow B$$

$$B \ \& \ okY \Rightarrow \sim C$$

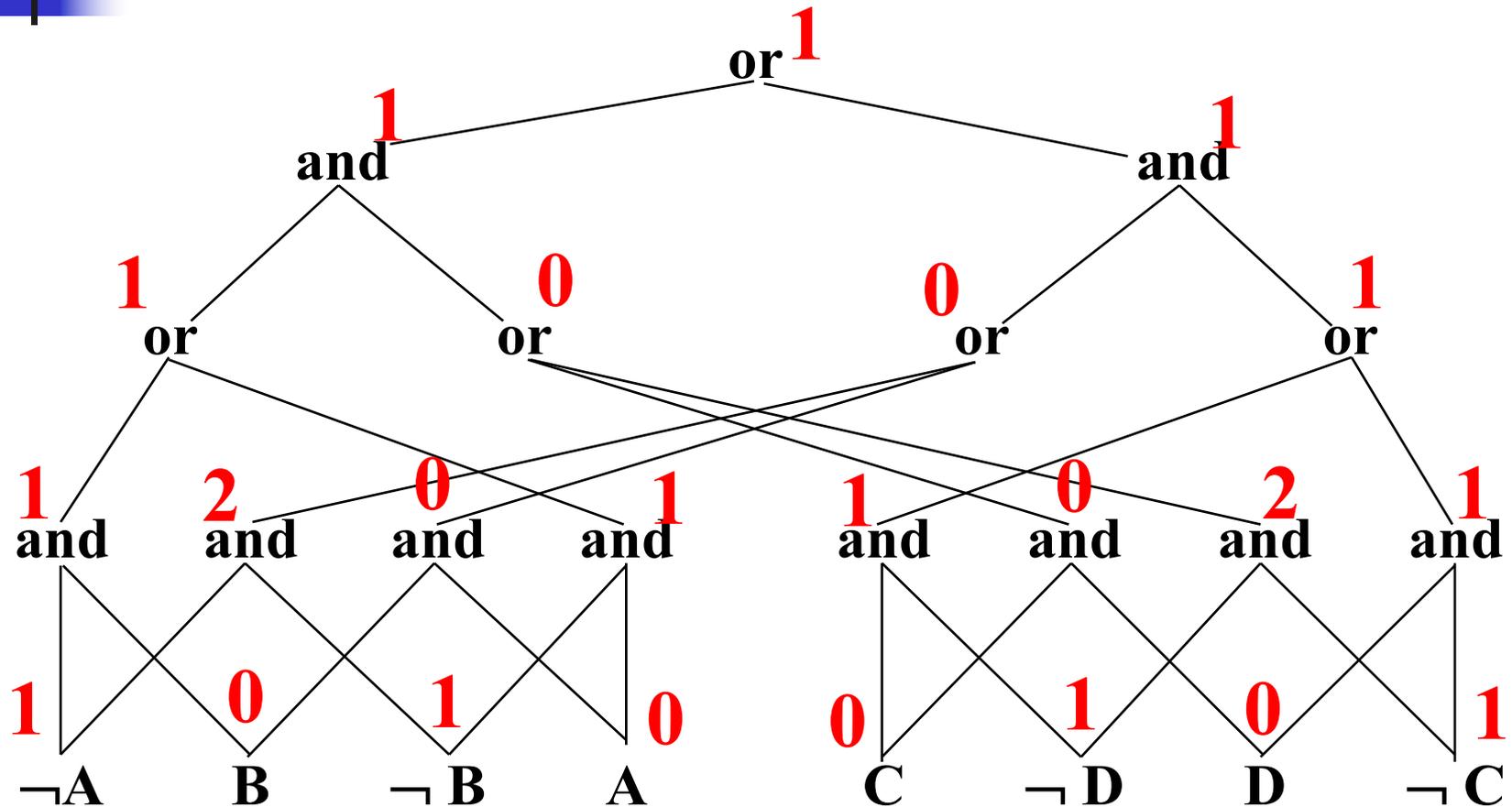
$$\sim B \ \& \ okY \Rightarrow C$$

okX	okY	A	B	C	
true	true	true	false	true	1
true	false	true	false	false	3
.
.
					<hr/>
					1

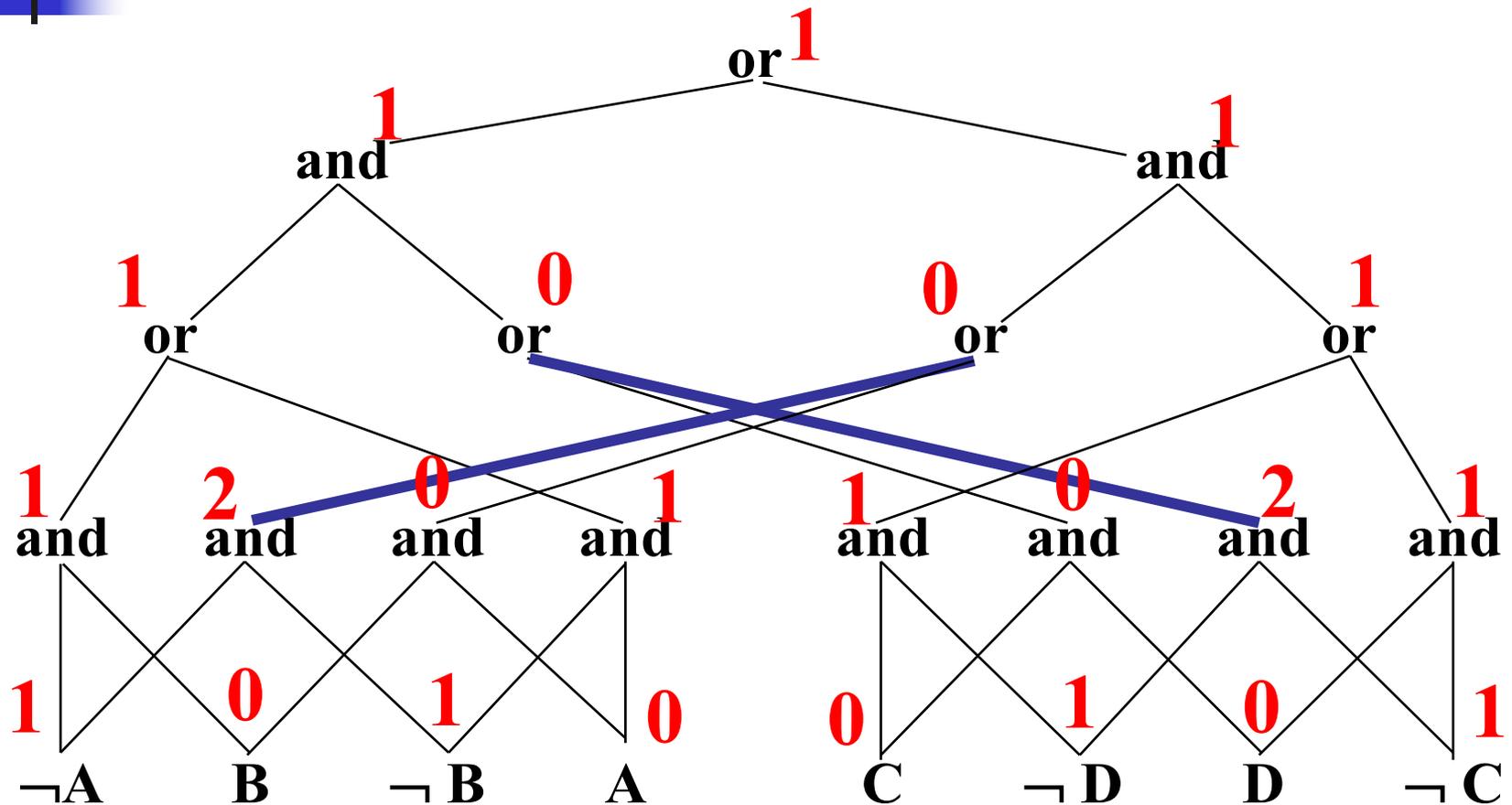
Minimum Cardinality



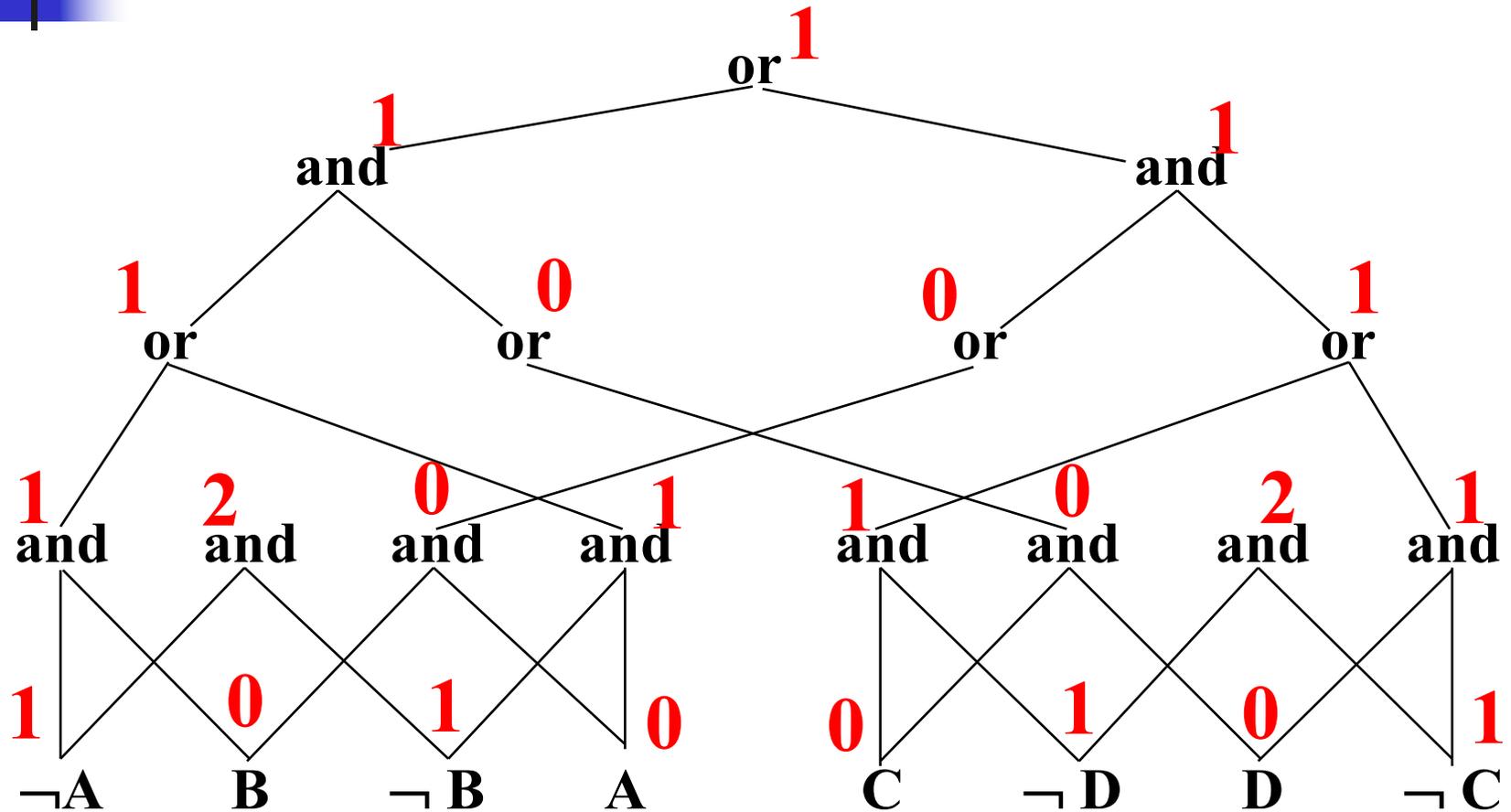
Minimizing: Requires Smoothness



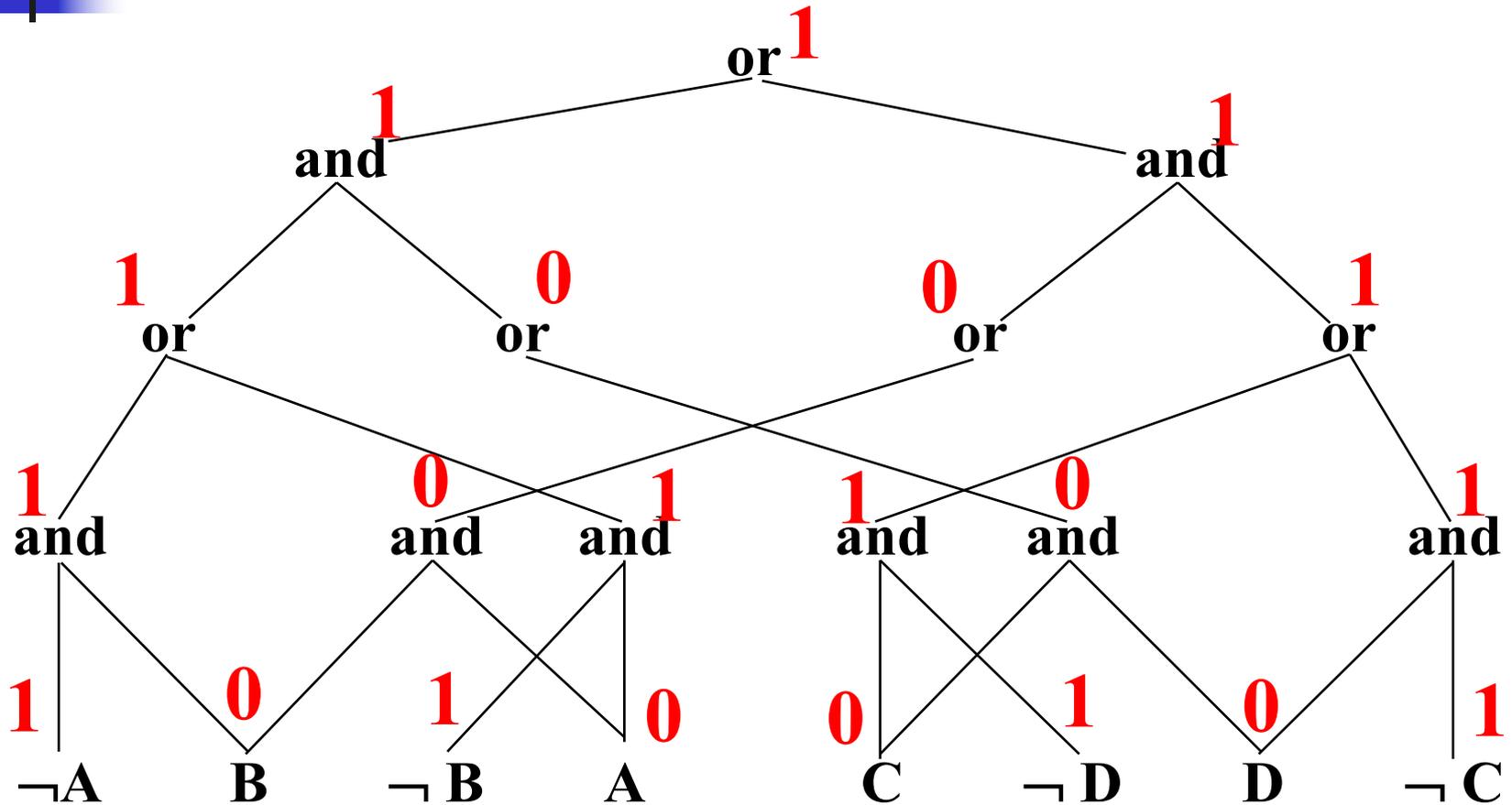
Minimizing



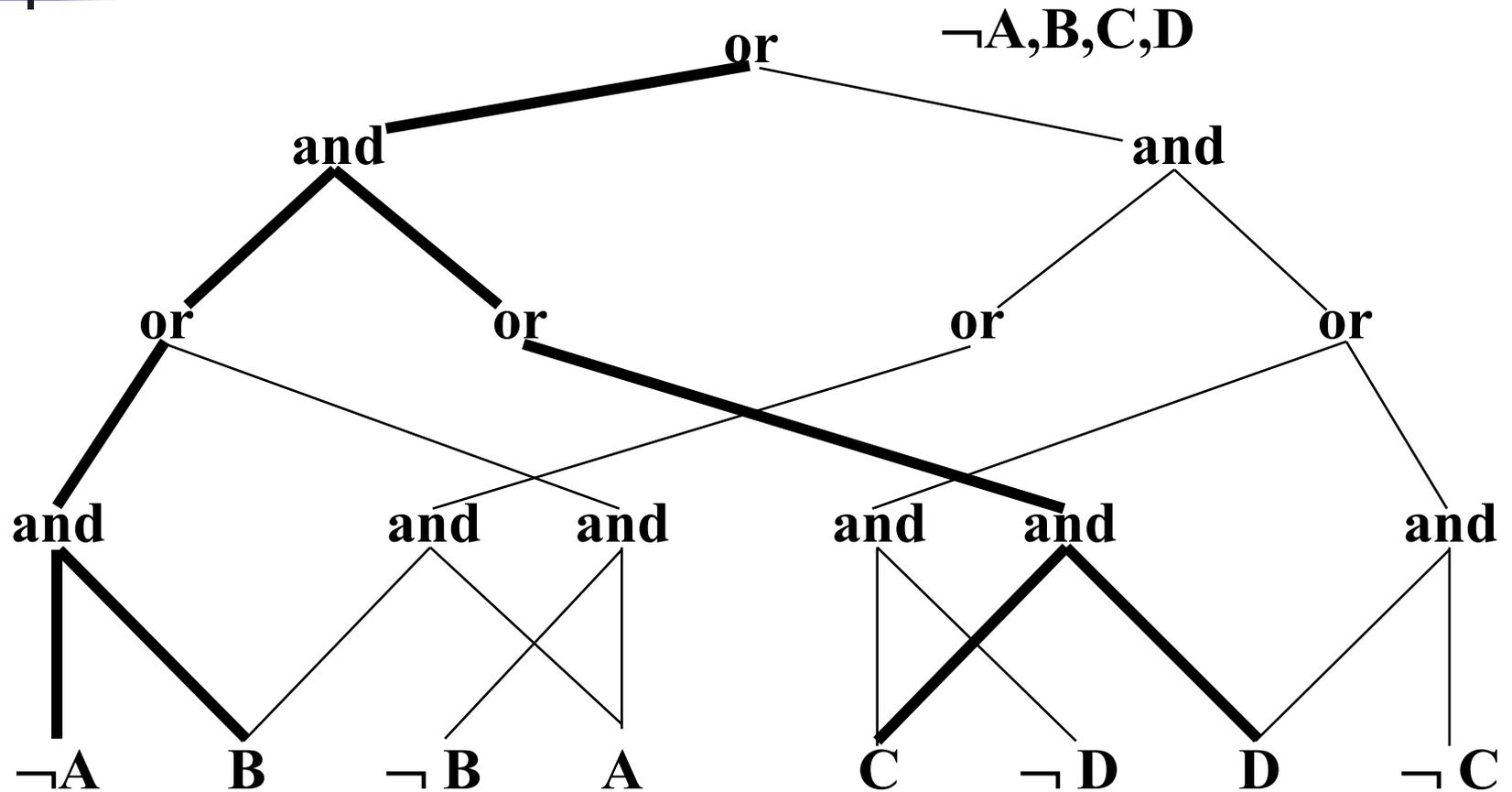
Minimizing



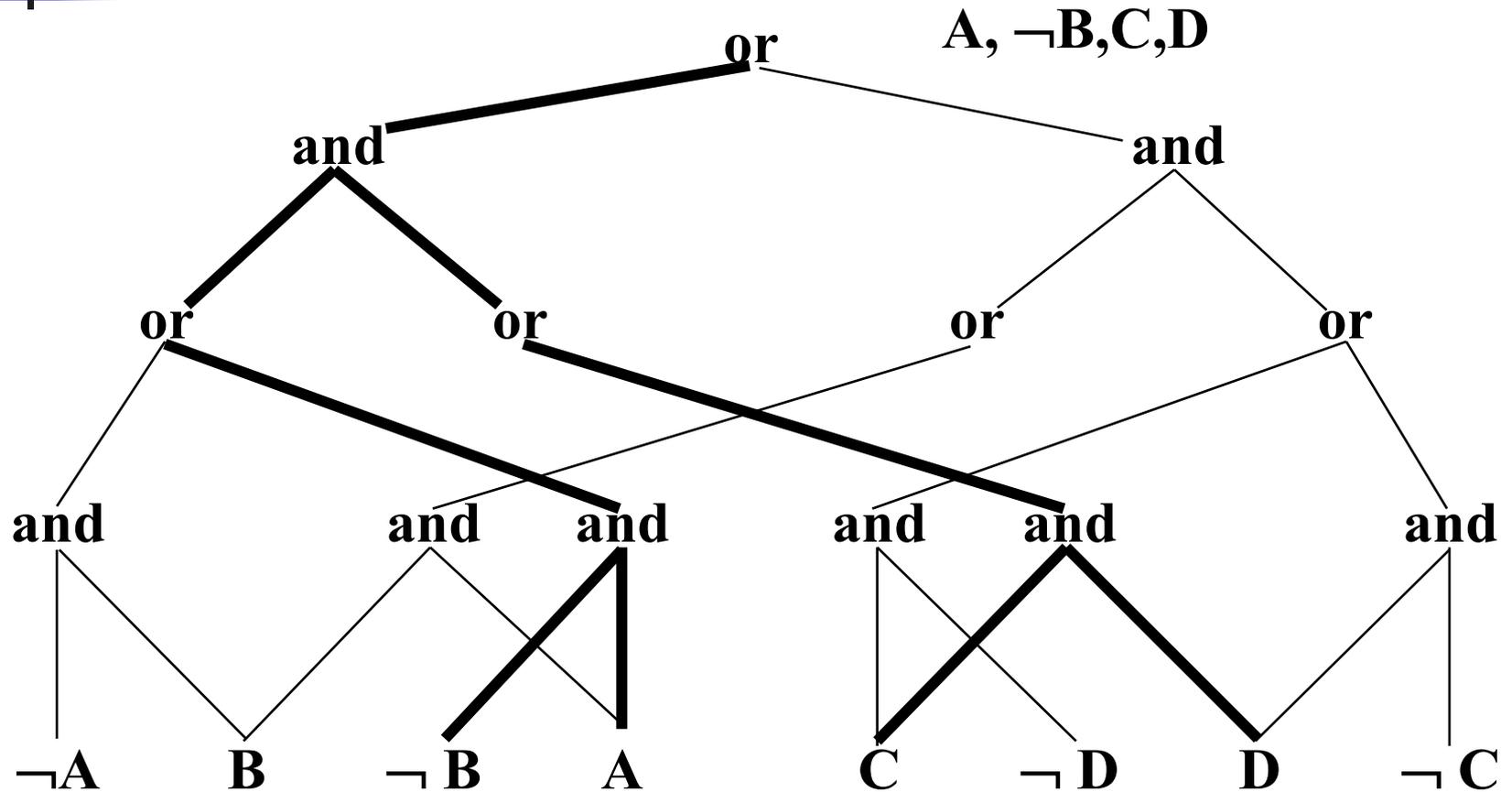
Minimizing



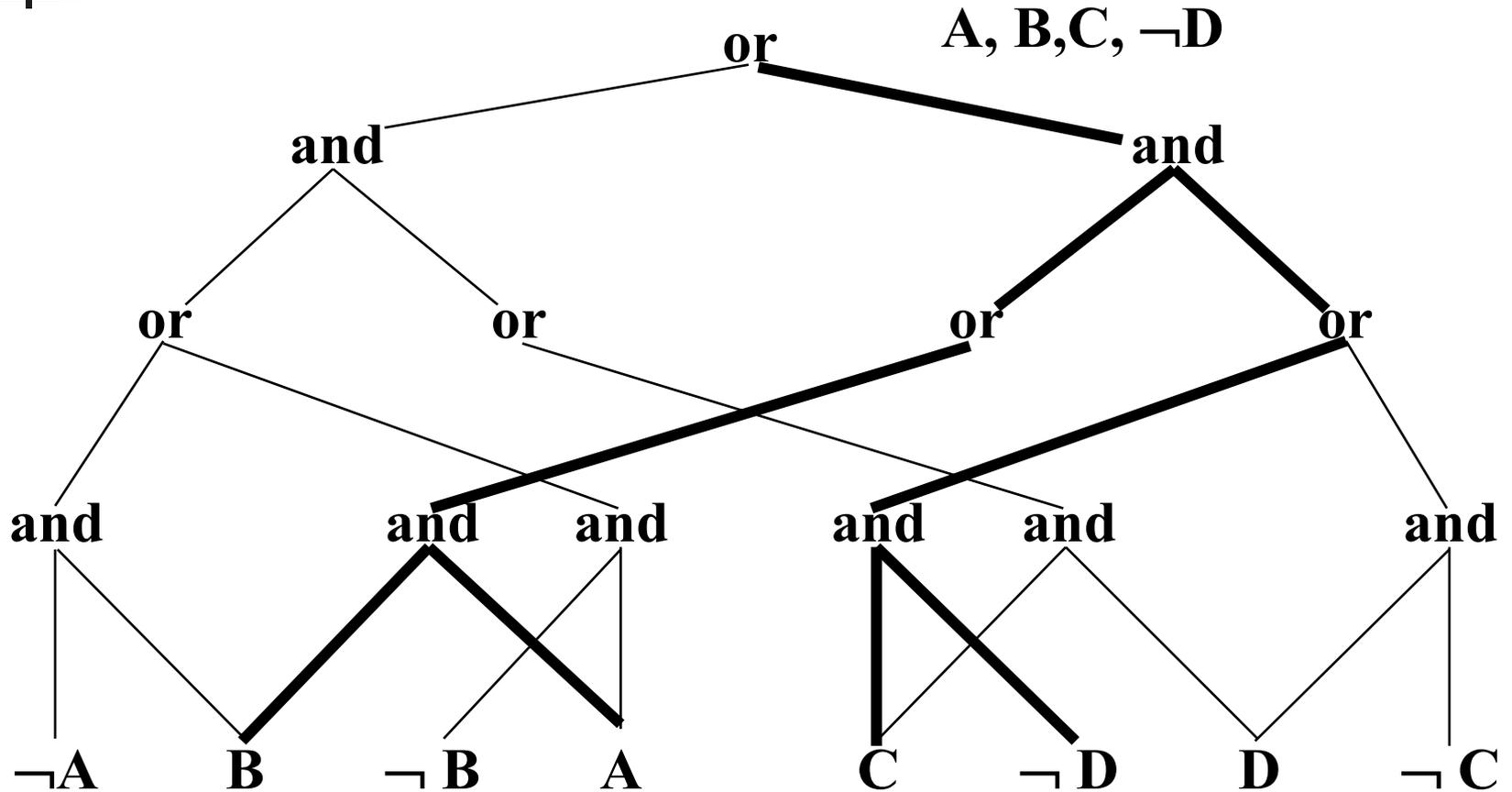
Minimizing



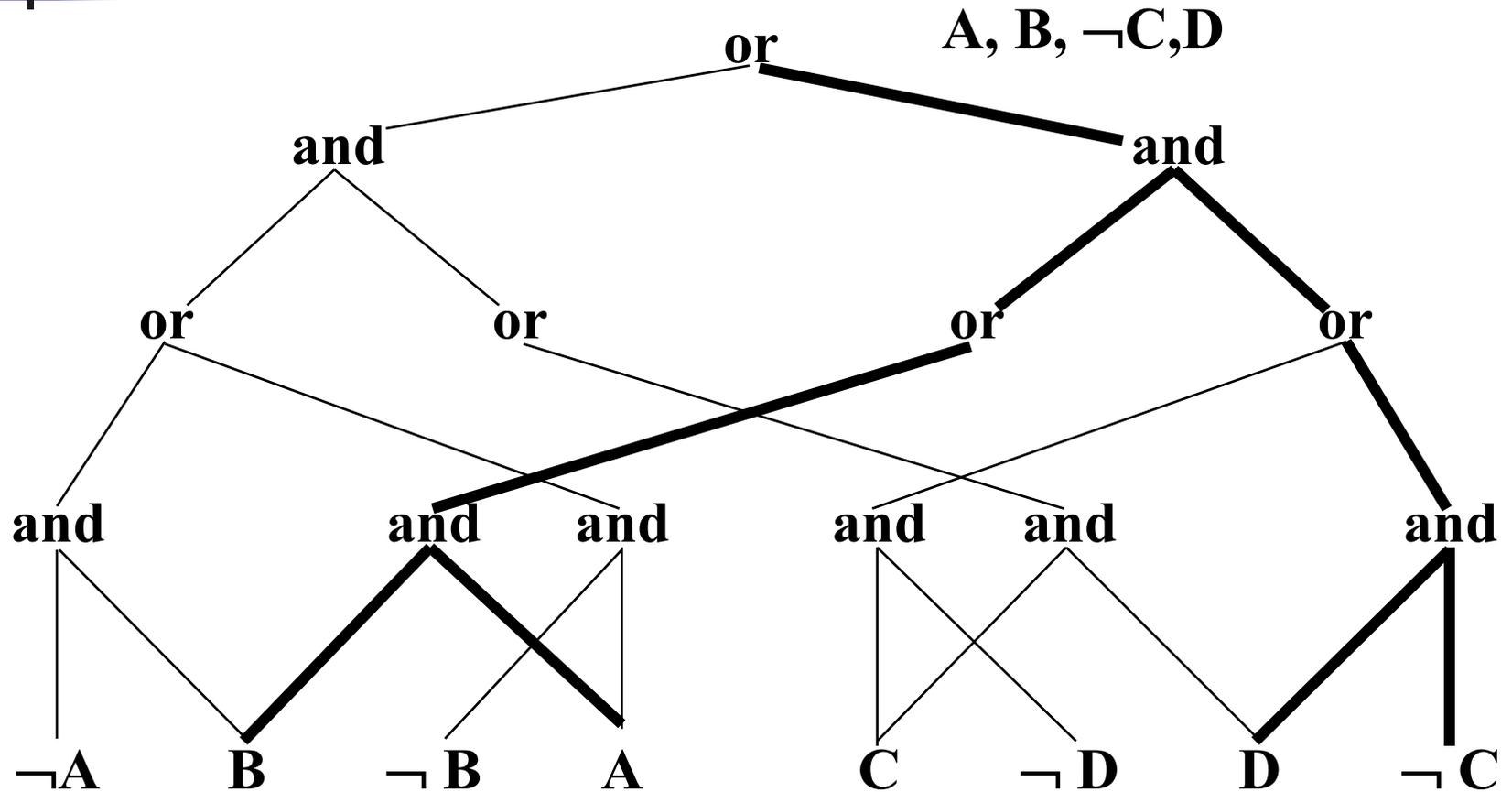
Minimizing

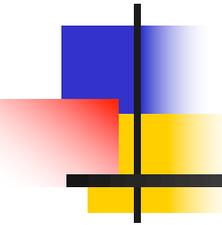


Minimizing



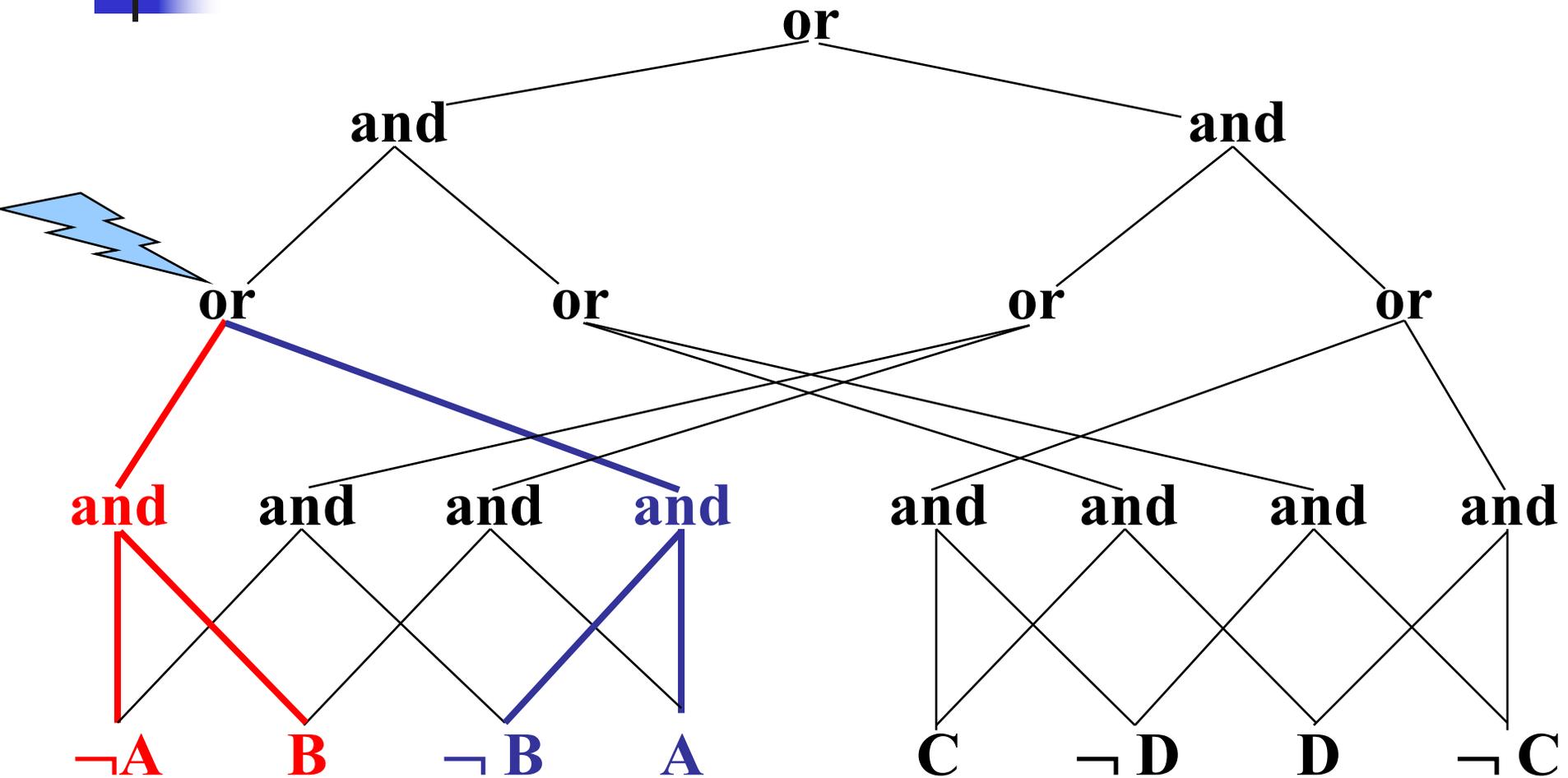
Minimizing



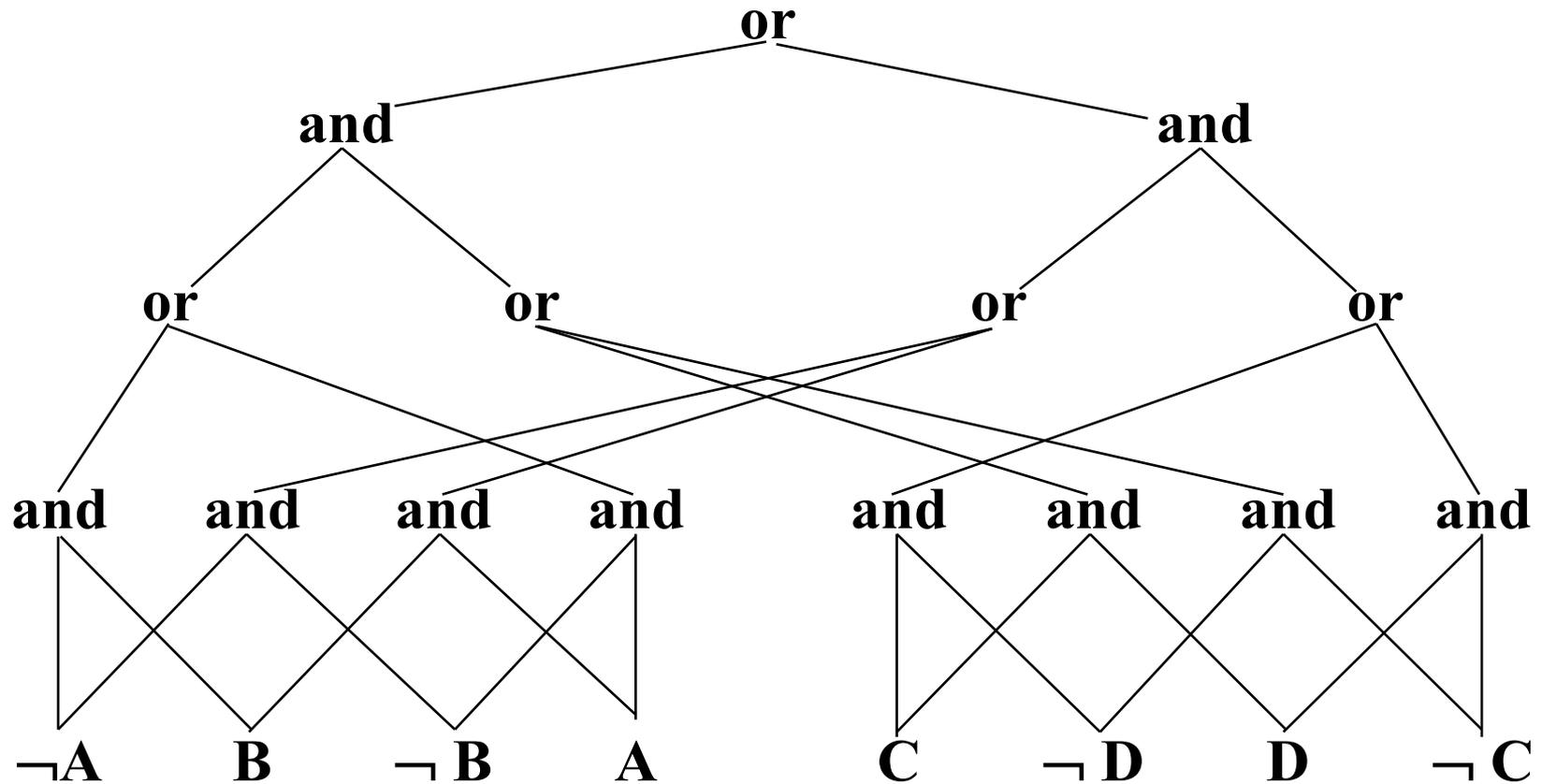


Determinism

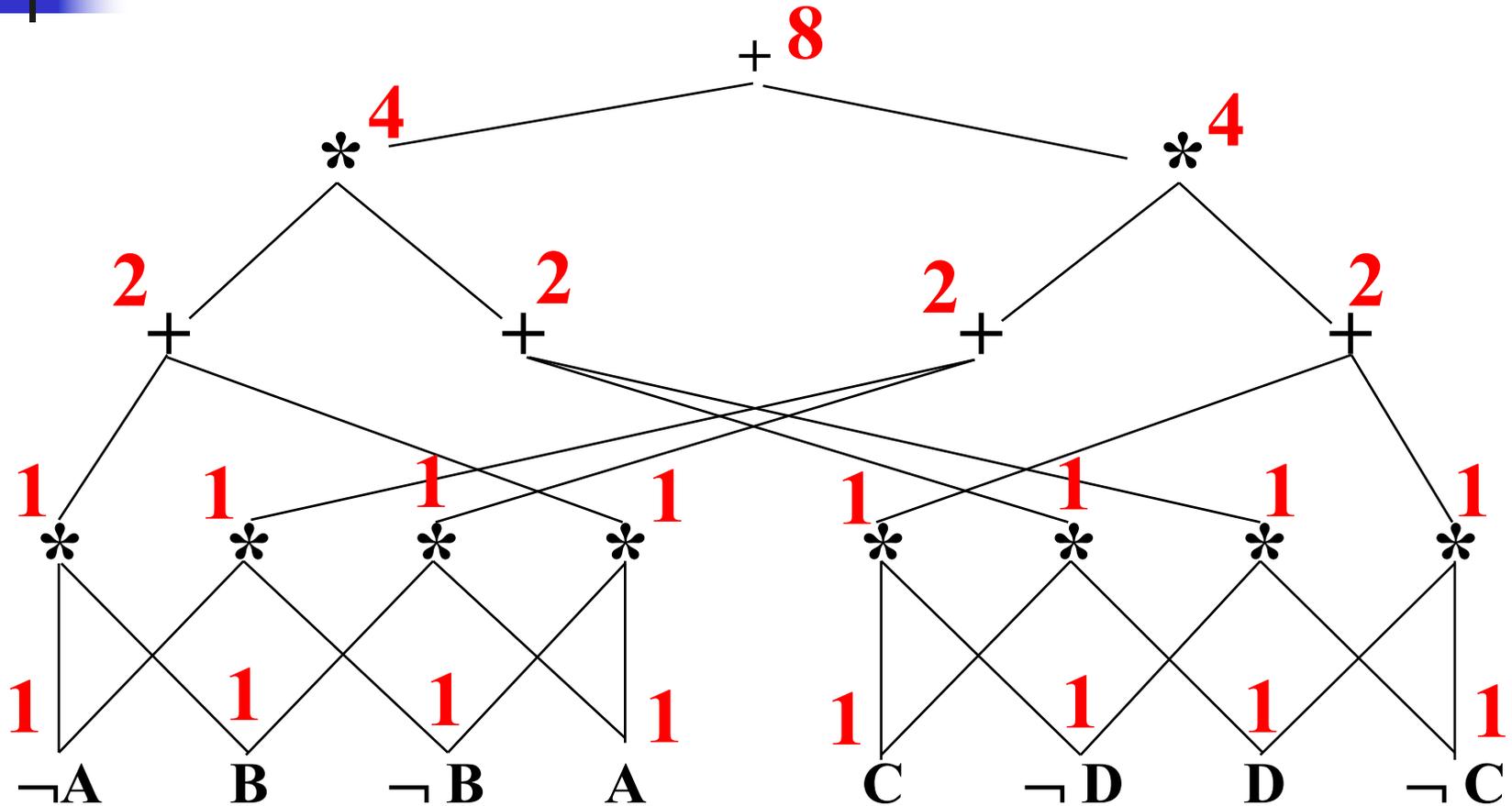
Determinism

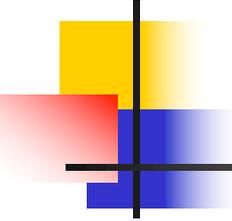


Counting Models



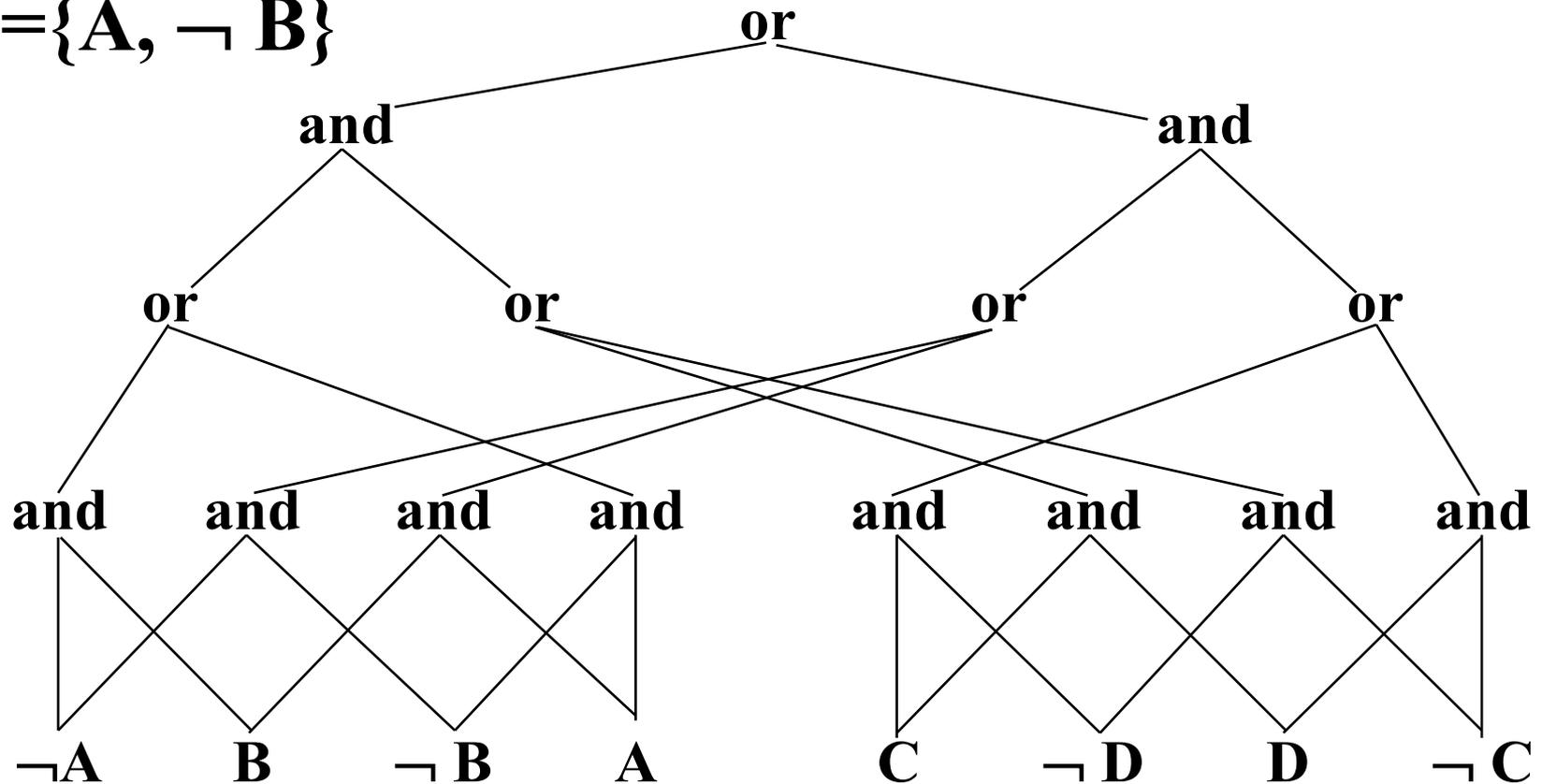
Counting Graph





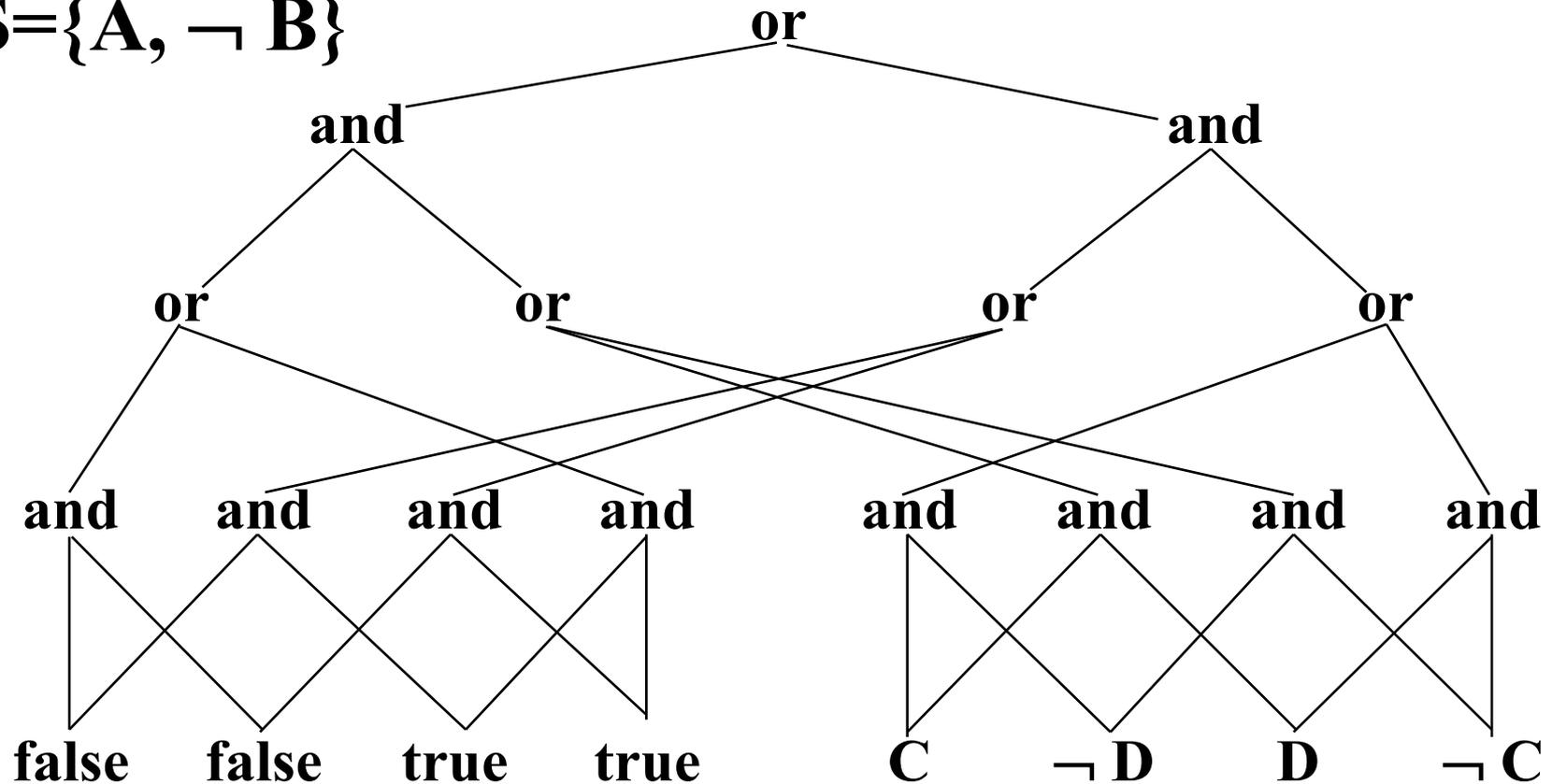
Counting Models

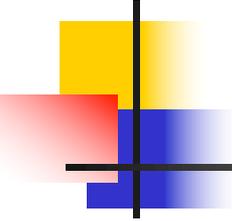
$S = \{A, \neg B\}$



Counting Models

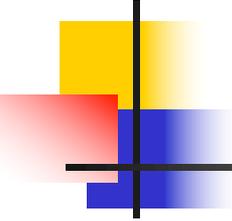
$S = \{A, \neg B\}$





Determinism

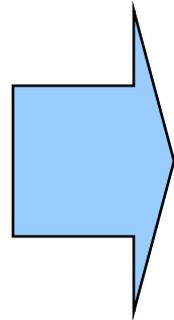
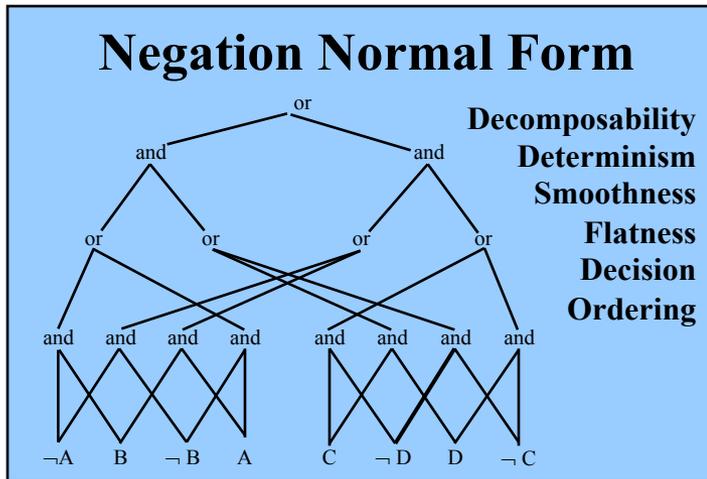
Query	d-DNNF
CO: Consistency	Yes
VA: Validity	Yes
CE: Clausal entailment	Yes
SE: Sentential entailment	
IP: Implicant testing	Yes
EQ: Equivalence testing	?
MC: Model Counting	Yes
ME: Model enumeration	Yes



OBDDS and SDDs

Query	OBDD
CO: Consistency	Yes
VA: Validity	Yes
CE: Clausal entailment	Yes
SE: Sentential entailment	Yes
IP: Implicant testing	Yes
EQ: Equivalence testing	Yes
MC: Model Counting	Yes
ME: Model enumeration	Yes

A Knowledge Compilation MAP



Polytime Operations

Consistency (CO)
Validity (VA)
Clausal entailment (CE)
Sentential entailment (SE)
Implicant testing (IP)
Equivalence testing (EQ)
Model Counting (CT)
Model enumeration (ME)

Projection (exist. quantification)
Conditioning
Conjoin, Disjoin, Negate

Succinctness

UCLA Automated Reasoning Group



4:30

CACM Oct. 2018 - Human-Level Intelligence or Animal...

Association for Computing Ma...
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The Three Eras

Knowledge Representation and Reasoning
(models); logic
AIGS

Machine Learning
(models + functions); probability
Statistical Optimisation

Neural Networks
(functions); neural networks
Engineering

12:15

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