

# Cardinality Encodings for Graph Optimization Problems

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# Definitions

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given a *propositional* formula in **CNF**,

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decide whether or not it is **satisfiable**

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**Example**

$$\mathcal{F}_1 = (x_1) \wedge (\neg x_1 \vee \neg x_2)$$

given a *propositional* formula in **CNF**,  
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### Example

$\mathcal{F}_1 = (x_1) \wedge (\neg x_1 \vee \neg x_2)$  — satisfiable

given a *propositional* formula in **CNF**,  
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$$\mathcal{F}_2 = (x_1) \wedge (\neg x_1 \vee \neg x_2) \wedge (x_2)$$

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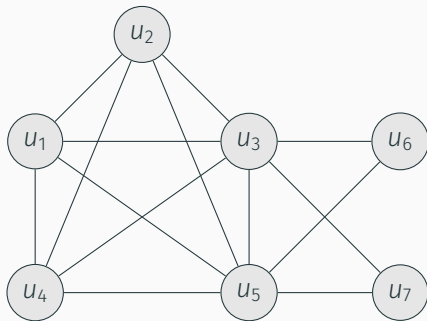
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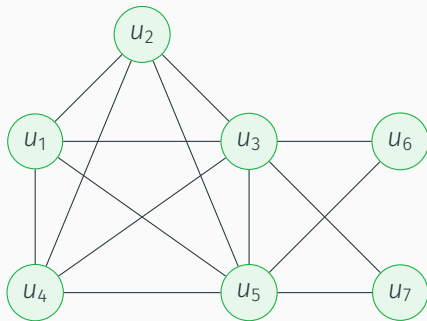
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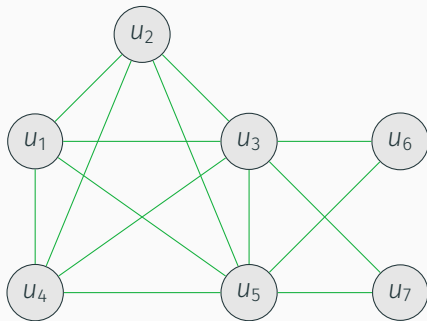
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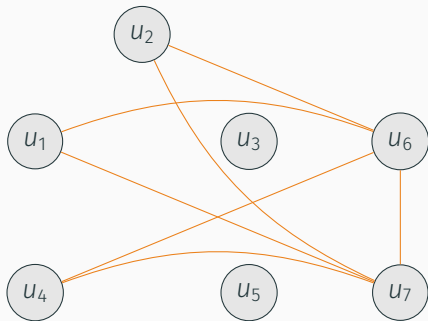
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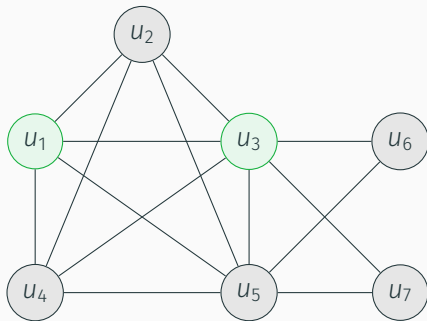
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graph  $G^C = (V, E^C)$  :



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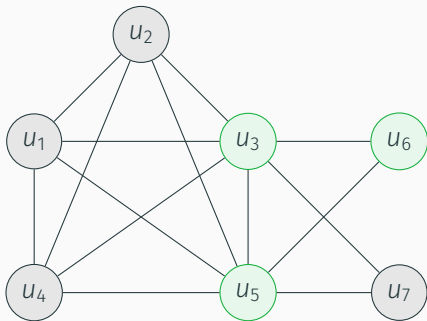
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clique of size 2

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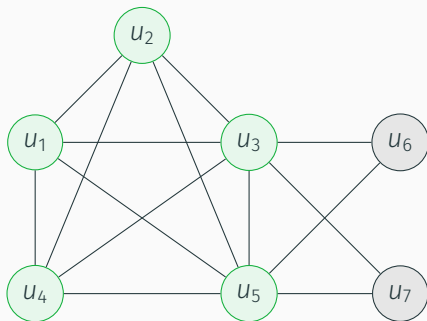
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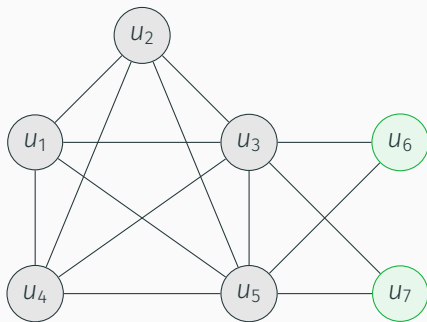
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maximum clique of size 5

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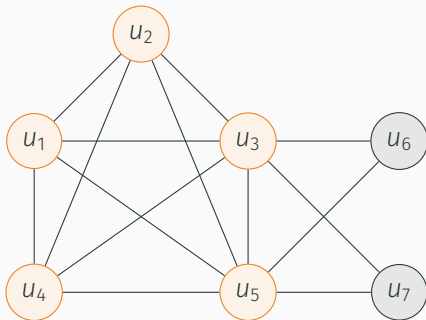


**independent set** of size 2



# Graph problems

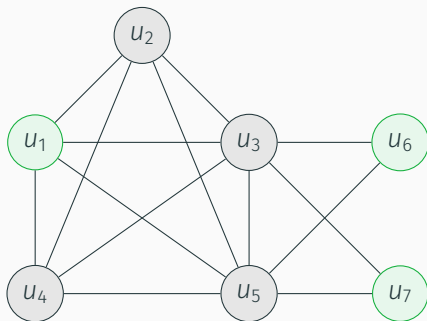
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**vertex cover** of size 5

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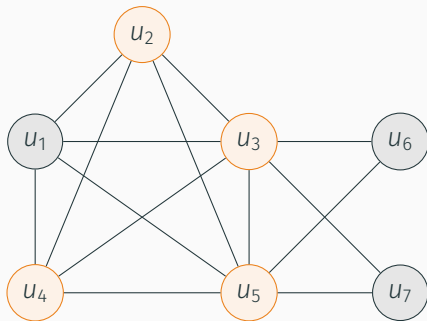
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# Graph problems

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minimum vertex cover of size 4

## Hidden pairwise encodings

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## Standard encoding of maximum clique

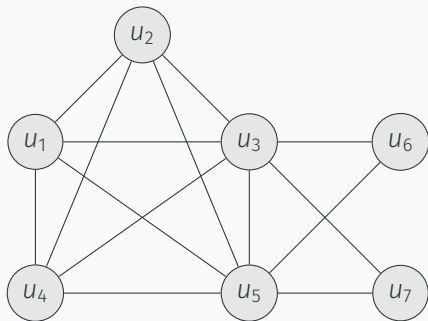
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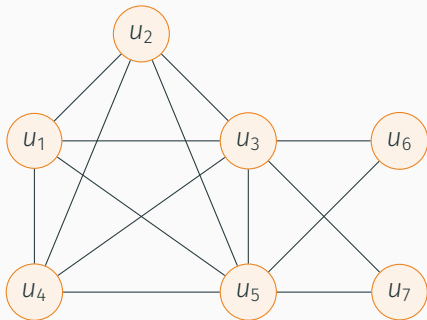
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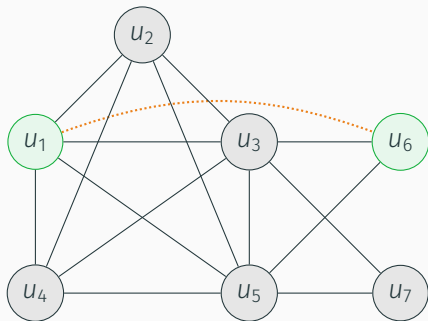


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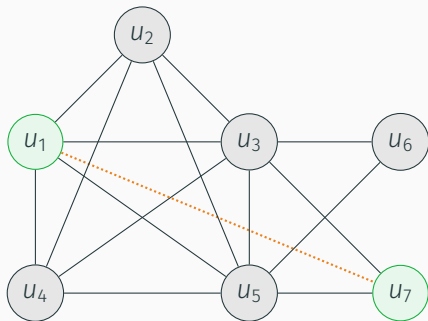


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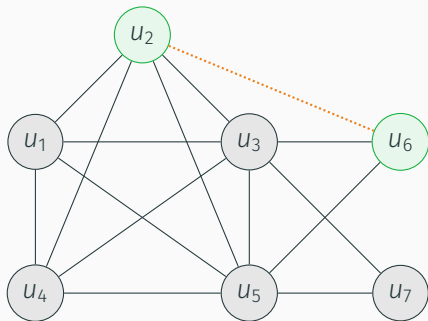


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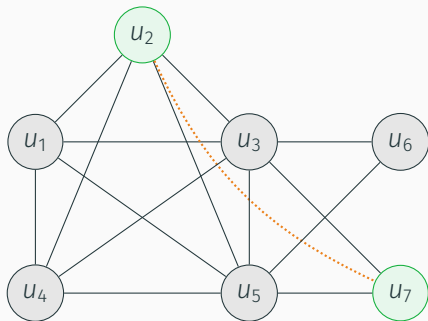


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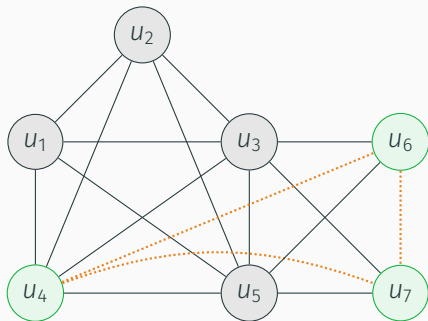


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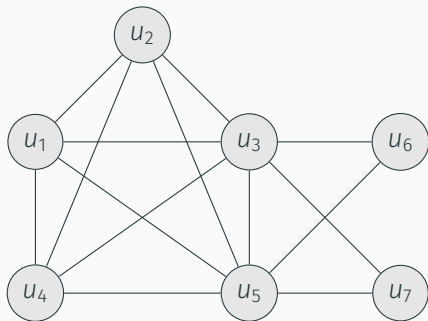


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solve  $\mathcal{F}$  with **MaxSAT**

for `ca-dblp-2012`<sup>1</sup>,

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**impossible** to solve and **hard** to *represent*

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$$G = (V, E)$$

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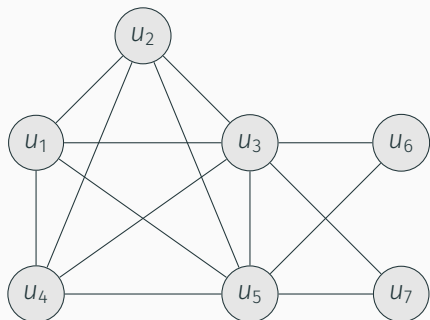
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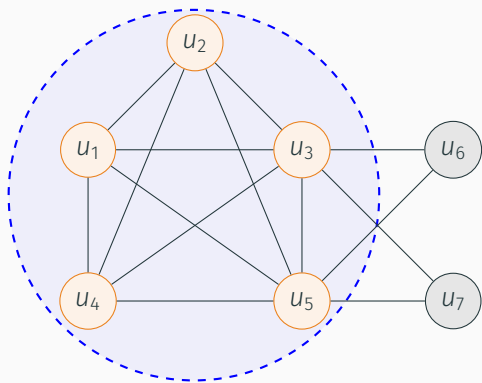
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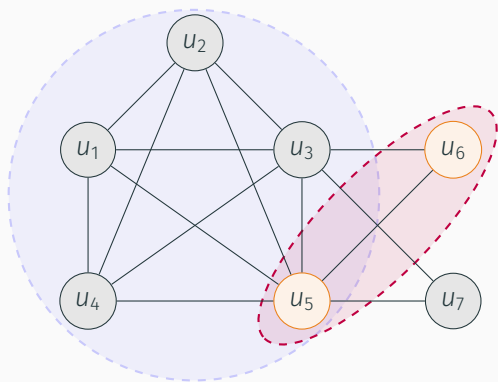
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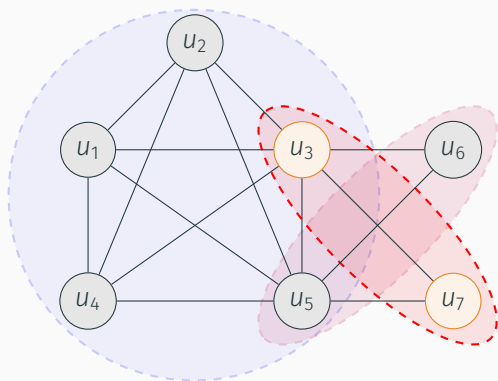
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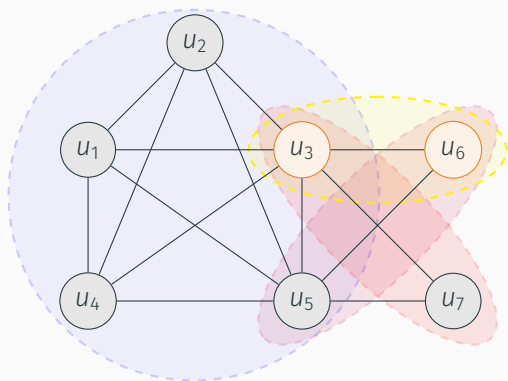
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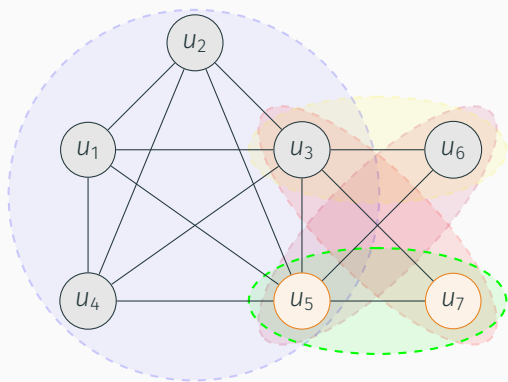
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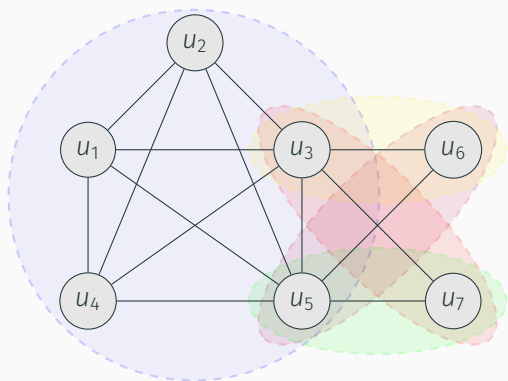
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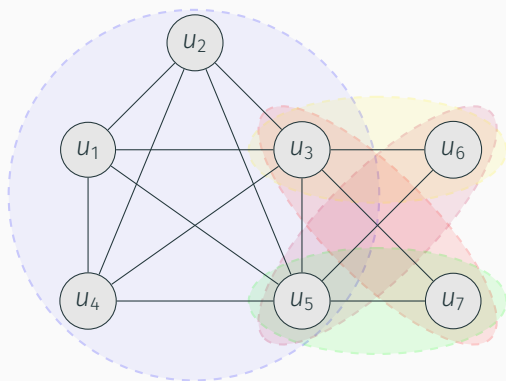
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# Solving MaxClique with SAT

---

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- **Competition:**
  1. Cliquer 1.21
  2. FMC
  3. IncMaxCLQ
  4. LMC



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- **Benchmarks:**

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- **Machine configuration:**

- Intel Xeon E5-2630 2.60GHz with 64GByte RAM
- running Ubuntu Linux
- 3600s timeout
- 10GByte memout



# Experimental evaluation

Instance	SATClq	Cliquer	FMC	IncMaxCLQ	LMC
comm-n1000	0.19	0.02	0.05	0.11	0.5
comm-n10000	—	0.99	—	12.33	0.93
ca-AstroPh	0	101.17	0.43	—	0.69
ca-citeseer	0	354.46	0.92	—	1.03
ca-coauthors-dblp	0	—	29.65	—	9.42
ca-CondMat	0	71	0.13	—	0.55
ca-dblp-2010	0	353.85	0.87	—	0.92
ca-dblp-2012	0	—	1.39	—	1.07
ca-HepPh	0	44.61	0.57	—	0.6
ca-HepTh	0	27.84	0.06	—	0.49
ca-MathSciNet	0	—	1.27	—	1.07
ia-email-EU	2.47	7.15	0.08	—	0.49
ia-reality-call	0	3.98	0.03	—	0.44
ia-retweet-pol	1.76	2.35	0.16	—	0.49
ia-wiki-Talk	—	60.48	4.21	—	0.73
rt-pol	1.7	2.39	0.19	—	0.49
rt_barackobama	0	0.46	0.45	6.63	0.46
soc-advogato	0.15	0.25	0.11	4.91	0.49
soc-epinions	—	101.67	0.21	—	0.82
soc-gplus	0.01	2.82	0.45	—	0.47
tech-as-caida2007	0.01	5.26	0.09	—	0.48
tech-internet-as	0.02	12.23	0.45	—	0.52
tech-pgp	3.05	0.71	0.07	—	0.45
tech-WHOIS	—	10.13	—	6.31	0.49
web-arabic-2005	0	151.31	2.43	—	1.57
web-baidu-baike-related	0.94	—	—	—	2.54
web-it-2004	0	—	25.32	—	4.87
web-NotreDame	0	—	3.76	—	1.37
web-sk-2005	0	97.44	0.34	—	0.64
p5sparse1	2.88	1031.15	—	12.17	0.48
p5sparse2+10clq20	1.9	24.42	—	—	0.54
p5sparse3+10clq20	3.62	150.15	—	—	0.58
p6sparse1	48.34	—	—	—	0.53
p6sparse2+10clq20	42.65	—	—	—	0.64
p6sparse3+10clq20	50.88	—	—	—	0.7
<b>Solved (out of 35)</b>	<b>31</b>	<b>26</b>	<b>26</b>	<b>6</b>	<b>35</b>

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ia-email-EU	2.47	7.15	0.08	—	0.49
ia-reality-call	0	3.98	0.03	—	0.44
ia-retweet-pol	1.76	2.35	0.16	—	0.49
ia-wiki-Talk	—	60.48	4.21	—	0.73
rt-pol	1.7	2.39	0.19	—	0.49
rt_barackobama	0	0.46	0.45	6.63	0.46
soc-advogato	0.15	0.25	0.11	4.91	0.49
soc-epinions	—	101.67	0.21	—	0.82
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Questions?